EVALUATION OF THE INTERLEAVED SOURCE CODING (ISC) UNDER PACKET CORRELATION

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ABSTRACT

Network impairments such as delay and packet losses have severe impact on the presentation quality of many predictive video sources. Prior researches (e.g., [1]-[3][5]-[9][11]-[13]) have shown efforts to develop packet loss resilient coding methods to overcome such impairments for realtime streaming applications. *Interleaved Source Coding* (ISC) is one of the error resilient coding methods, which is based on an optimum interleaving of predictive video coded frames transmitted over a *single* erasure channel. ISC employs a Markov Decision Process (MDP) and a corresponding dynamic programming algorithm to identify the optimal interleaving pattern for a given channel model and a transmitting sequence. ISC has shown to significantly improve the overall quality of predictive video coded stream over a lossy channel without complex modifications to standard video coders [7][8]. In this paper, ISC is evaluated over channels with memory. In particular, we analyze the impact of *packet correlation* [15][16]of the popular Gilbert model on ISC-based packet video over a wide range of packet loss probabilities. Simulations have shown that ISC advances the traditional method as either the loss rate increases or the packet correlation decreases.

1. INTRODUCTION

With the expansion of the underlying infrastructure of the IP network, streaming video has become one of the most popular on-line realtime Internet applications. However, despite of the growth of internet infrastructure, such realtime streaming services often lack Quality-of-Service (QoS) guarantees due to the unreliable nature of IP networks. To improve playback quality of realtime streaming video under such network condition, coding techniques that are resilient to packet losses are in need. Coding techniques (e.g., [1]-[3][5]-[9][11]-[13].) are few examples of methods to be resilient to packet losses.

In this paper, we evaluate interleaved source coding (ISC) [7][8], a recently proposed packet loss resilient coding technique, over erasure channels with memory. ISC codes a single video sequence into two sub-sequences and transmits them over a single erasure channel to reduce the frequency and impact of the cascaded effect of packet losses and related propagation of errors of predictive coded video. Moreover, to avoid complex modifications to standard video coders, ISC uses frozen frame technique for the decoder failed frames from packet losses. The objective of ISC is to elect an optimum interleaving sub-sequence set with a Markov Decision Process (MDP) and a Dynamic Programming for a given erasure channel model such that the impact of packet losses is minimal, in other words, the number of loss impacted frames are minimum.

In this paper we evaluate the performance of ISC with various packet loss rates in conjunction with a *packet correlation* model [15][16]. The *packet correlation* model uses the average loss rate p and the packet correlation ρ to represent the state transition probabilities of the two state Markov model (a.k.a. Gilbert model).

The remainder of this paper is organized as follows: In Section 2, we briefly describe the ISC coding method following by the description of the *packet correlation* model in Section 4. In the last section, ISC is evaluated over *packet correlation* based lossy channel models with various packet loss rates and streams coded using the MPEG-4 video standard.

2. INTERLEAVED SOURCE CODING

In general, a predictive video coding partitions a single lengthy sequence into a number of shorter length Group Of Video object planes (GOVs) to limit the impact of possible errors or losses into individual GOVs. The *interleaved source coding* (ISC) [7][8] is a pre- and post-process of predictive source coders to reduce the impact of losses within a given GOV, hence improves the playback quality of predictive video over lossy packet networks. ISC separates a single video sequence into two sub-sequences and each sub-sequence is encoded using separate video encoders. Then the coded subsequences are merged into a single stream in the original-sequence frame order for transmission. Prior to the transmission of the merged ISC stream, information on the interleaving pattern employed by the ISC must be transmitted to the decoder. Once the merged stream is transmitted, with the interleaving pattern information, the stream is separated for the decoder and the separated streams are decoded independent to each other. For playback, the separated sub-sequences' frames are merged into the proper order.

$$\begin{gathered} s=\left\{0 \;\; 1 \;\; \cdots \;\; 2N-1\right\}=\bigcup_{j=1}^{2} s^{(j)} \\ \bigcap_{j=1}^{2} s^{(j)}=\varnothing, \qquad \forall j, size\left(s^{(j)}\right)=N \\ \sum\left\{s^{(j)}(2)-s^{(j)}(1),\cdots,s^{(j)}(N)-s^{(j)}(N-1)\right\}>N-1 \end{gathered} \tag{1}$$

ISC uses the above interleaving constraints when separating a single sequence into two sub-sequences, $s^{(j)}$, represented by an index set, $j=\{1,2\}$. Here, $(2N-1)$ represents the number of frames in the original non-interleaved sequence s where it could be the number of frames in a GOV.[1]

For an interleaving sub-sequence set $\mathbb{S}=\{s^{(1)},s^{(2)}\}$, the numbers in $s^{(j)}$ represent the frame locations in s and the coded stream's frame transmission order and this interleaving information must be transmitted (e.g., as meta data) prior to the coded stream transmission.

For a given GOV size, the size of the possible number of the interleaving set $\mathbb{K}$ could be quite large for any reasonable GOV size $(2N-1)$, hence over a lossy network channel, identifying an optimum interleaving set with the best quality decoded video could be very computationally expensive task. To resolve this problem, ISC employs a decision-based search algorithm to elect an optimal interleaving set for a given erasure-channel model and a video sequence.

The analysis and modeling of packet losses over the Internet (e.g., [14]) and wireless networks (e.g., [6]) have proven that these losses exhibit Markovian properties. A *Markov Reward Process* (MRP) (e.g., [4][10]) can be used to estimate the system's performance after n packet transmissions using the channel's Markov transition probabilities and some models for the *rewards* associated with each system state. For an erasure-channel model, the state space of two state Markov model is mapped to good (0) and bad (1) packet transmission states. In MRP, when the process reaches state i, the instant reward r_i assigned for each state are awarded to the process. Hence the aggregated reward $v(n-1)$ is a function of the number of transmitted packets which represents the performance of predictive sequence transmission after n packet transmissions over a lossy channel with the state transition matrix p. For a Gilbert channel model, $\{r_0,r_1\}=\{1,0\}$ are the instant reward constants and the process is awarded with *1* for a successful packet and *0* for a lost packet transmission, hence $v_i(n-1)$, represent the expected number of good packet transmissions with the initial process state i.

$$\begin{aligned} v(0) &= r=\{r_0,r_1\}^T, \quad v(1)=r+p\bullet v(0) \\ v(n-1) &= \left[v_0(n-1)\ \ v_1(n-1)\right]^T \\ &= r+p\bullet v(n-2)=\left(I+\sum_{m=1}^{n-1}p^m\right)\bullet r \end{aligned} \tag{2}$$

Associating Markov reward process with a series of actions and decision criteria gives a base model for a Markov Decision Process (MDP) [4][10]. For ISC, MDP is employed to find an interleaving set that is most suitable for a given decision criterium, maximize the number of frames (or associated packets) that can be decoded *correctly*. An aggregated MDP equation is obtained by incorporating an *interleaving set indicator* k, $k\in\mathbb{K}$ along with a set of *policies,* mappings from states to actions, and a set of *discount factors*, γ_a to equation.

$$v^{(k)}(n-1)=r_{a^{(k)}(n-1)}+\gamma_{a^{(k)}(n-1)}\times p\bullet v^{(k)}(n-2) \tag{3}$$

For predictive video coding, Coding, (C), or Skip, (S), action is taken for each state iteration. $\mathbb{S}^{(k)}=\{s^{(k,1)},s^{(k,2)}\}$ denotes an ISC sub-sequence set with respect to k, and to start the sub-sequence's reward computation from time instance 0, the frame numbers in $s^{(k,j)}$ are rewritten as following.

$$s^{*(k,j)}(n)=s^{(k,j)}(n)-s^{(k,j)}(0), \quad \text{for } 0\le n\le N-1 \tag{4}$$

In ISC, frames are coded (action C performed) at frame locations specified in $s^{*(k,j)}$. When the difference between two adjacent numbers in $s^{*(k,j)}$ exceeds 1, this indicates the presence of skipped frames and action S is performed for the frames in location m.

$$\begin{aligned} m &= \bigcup_{n=0}^{N-1}\left\{s^{*(k,j)}(n)+1,\cdots,s^{*(k,j)}(n+1)-1\right\} \\ &\forall n \ \mid\ s^{*(k,j)}(n+1)-s^{*(k,j)}(n)>1 \end{aligned} \tag{5}$$

In addition, ISC requires modification of the transition matrix p in association with actions. The state *1* for action C is considered as a trapping state since the decoder of predictive coding is forced to stop when a lost packet is detected. When stopped, ISC uses the last successfully decoded frame to replace the missing and effected frames and waits for a new GOV with successfully decoded *I*-frame to restart decoding process. Hence the transition probability for the policy $\{C,1\}$ to the next good state is set to *0*. Simultaneously, the discount factor for $\{C,1\}$ is set to *0*, since no further decoding is possible, the process is in the trapping state, hence, propagation of the aggregated reward is prohibited for the GOV. For all other policies, the aggregated rewards are propagated to the next state and the discount factors are set to *1*.

For the initial state, the instant reward is multiplied by a stationary probability π since it is assumed that the first packet in *I*-frames arrive to the process with the stationary probability since they do not have any temporal dependencies to the previously decoded frames. When frames are packetized, each frame is coded into different number of packets depending on the bitrate, the frame rate of the encoder, the packet size, the frame coding type, and the motion of the sequence. Hence, ISC incorporates this unpredictability of the variation of the number of packets per each coded frame with an average number of packets per frame η.

$$\eta=\left\lceil\frac{bitrate}{framerate\times packetsize}\right\rceil \tag{6}$$

$$v^{(k,j)}\left(s^{*(k,j)}(0)\right)=r_C(0)\times\left[\left(p_{00}\right)^{\eta-1}\ \ 1-\left(p_{00}\right)^{\eta-1}\right]^T\times\pi \tag{7}$$

$$\begin{aligned} &v^{(k,j)}\left(s^{*(k,j)}(n)\right)=r_C(n-1)+ \\ &\gamma_C\times\begin{bmatrix}\left(p_{00}\right)^{\eta-1} \\ 1-\left(p_{00}\right)^{\eta-1}\end{bmatrix}\times\begin{pmatrix}p_C\bullet p_S^{\left(s^{*(k,j)}(n)-s^{*(k,j)}(n-1)-1\right)} \\ \bullet v^{(k,j)}\left(s^{*(k,j)}(n-2)\right)\end{pmatrix} \text{ for } 1\le n\le N-1 \end{aligned} \tag{8}$$

Where the aggregated reward for a skipped frame is:

$$\begin{aligned} v^{(k,j)}(m) &= r_S(m)+\gamma_S\times p_S\bullet v^{(k,j)}(m-1) \\ &= p_S\bullet v^{(k,j)}(m-1) \end{aligned} \tag{9}$$

Since a frame is decoded if and only if all the packets in the coded frames are successfully transmitted, $\left[\left(p_{00}\right)^{\eta-1}\ \ 1-\left(p_{00}\right)^{\eta-1}\right]^T$ is multiplied to the aggregated reward. Hence the sum of aggregated rewards estimates the expected number of successfully decoded frames for each interleaving set k.

$$v^{(k)}=\sum_{j=1}^{2}\sum_{n=0}^{N-1}v^{(k,j)}\left(s^{*(k,j)}(n)\right) \tag{10}$$

[1] The same interleaving is applied to all GOVs in the sequence or a scene.

$$V^{(k)} = \bigcup_j V^{(k)}(s^{(k,j)})$$
$$\text{where } V^{(k)}(s^{(k,j)}(n)) = v^{(k,j)}\left(s^{*(k,j)}(n)\right), \quad 0 \le n \le N-1 \tag{11}$$

At last, ISC elects an optimal interleaving set k that satisfies our decision criterion, a set with the highest $v^{(k)}$.

$$\arg\max_k \left[v^{(k)}\right] \tag{12}$$

When the decoder encounters problem with the transmitted sequences, i.e., packet losses or errors in a transmitted packet, ISC replaces the decoder failed frames with the last successfully decoded frame, referred as the *replacement* frame, to continue the smooth video presentation until a successfully decoded frame arrives to restart the decoding process. In fact, the *distances* (in terms of number of frames) between the replacement frames and the replaced ones affects the smoothness of the flow and the overall playback quality of the sequence since the distances represent the correlations among the replacing frames. To employ the frame replacements action, ISC uses two different approaches. One is incorporating sequence specific correlation gain model to equation (10) and the other one is adopting a generalized correlation gain model to the same equation since measuring the sequence specific temporal correlation within a complete GOV may not be always feasible for realtime applications due to delay, complexity, and memory constraints.

For generalized correlation model ISC, the generalized correlation gain, $g^{(k)*}$, is computed with the following equations.

$$VI^{(k)} = \bigcup_j VI^{(k)}(s^{(k,j)})$$
$$\text{where } VI^{(k)}(s^{(k,j)}(0)) = v^{(k,j)}\left(s^{*(k,j)}(0)\right),$$
$$VI^{(k)}(s^{(k,j)}(n)) = v^{(k,j)}\left(s^{*(k,j)}(n)\right) - v^{(k,j)}\left(s^{*(k,j)}(n-1)\right), \tag{13}$$
$$\text{for } 1 \le n \le N-1$$

$$W^{(k)*} = \left(\left(D^{(k)*} \times VI^{(k)^T}\right) \div \sum_{y=1}^{\text{GOV SIZE}} D^{(k)*}_{x,y}, \forall x\right)^{\wedge} D^{(k)}$$
$$\text{where } D^{(k)*} = \begin{Bmatrix} 0, & \forall D^{(k)}_{x,y} \neq 0 \\ 1, & \text{otherwise} \end{Bmatrix} \tag{14}$$

$$G^{(k)*}_{x,y} = \begin{Bmatrix} W^{(k)*}_{x,y} \times V^{(k)*}\left(y - D^{(k)}_{x,y}\right), & \forall D^{(k)}_{x,y} \neq 0 \\ V^{(k)}(y), \quad \text{otherwise} \end{Bmatrix} \tag{15}$$

$$g^{(k)*} = \sum_{x,y=1}^{\text{GOV SIZE}} G^{(k)*}_{x,y}, \quad \forall x,y \le \text{GOV SIZE} \tag{16}$$

Here, $VI^{(k)}$ is the set of the reward increments at each sub-sequence's reward calculation iteration. With respect to distance matrix $D^{(k)}$ and $VI^{(k)}$, the weight matrix $W^{(k)*}$ is calculated where $\left(\left(D^{(k)*} \times VI^{(k)^T}\right) \div \sum_{y=1}^{\text{GOV SIZE}} D^{(k)*}_{x,y}, \forall x\right)$ is the average reward increment of the successfully decoded frames in case of a single error in a GOV. Multiplying (14) by the replacement frame's aggregated reward estimates the correlation-based aggregated reward of the replaced frame. Hence, the decrement is assumed to be exponential with respect to temporal distances from the replacement frames to the replaced ones.

The generalized ISC dynamic programming model elects an optimal interleaving set using the following equation.

$$\arg\max_k \left[v^{(k)} + g^{(k)*}\right] \tag{17}$$

3. GILBERT MODEL PARAMETER PAIRS USING PACKET CORRELATION

In the Gilbert channel model, the steady state probabilities in good state and bad state are represented as following.

$$\pi(0) = \frac{p_{10}}{p_{01} + p_{10}} \text{ and } \pi(1) = \frac{p_{01}}{p_{01} + p_{10}} \tag{18}$$

These values give a coarse measure of a given channel's packet transmission behavior. However, for statistical channel modeling, instead of the above probabilities, the transition probabilities p_{01} and p_{10} (or $p_{00} = 1 - p_{01}$ and $p_{11} = 1 - p_{10}$) could be used to characterize Gilbert channel model. Since it is difficult to properly model a Gilbert channel with arbitrary transition probabilities, p_{01} and p_{10}, a more meaningful pair of parameters is the average loss rate, p_1, and the *packet correlation*, ρ; this pair can provide a practical and useful insight of the channel while representing the state transition probabilities. The average loss rate and the correlation between two consecutive packets can be defined as follows [15][16]:

$$p_1 = \frac{p_{01}}{p_{01} + p_{10}}, \qquad \rho = p_{01} + p_{10} - 1 \tag{19}$$

Hence, the transition probabilities represented by p_1 and ρ are:

$$p_{00} = 1 - p_1(1-\rho), \qquad p_{01} = p_1(1-\rho)$$
$$p_{10} = (1-p_1)(1-\rho), \quad p_{11} = 1-(1-p_1)(1-\rho) \tag{20}$$

In addition, the steady state probabilities are directly related to the loss rate p_1; $\pi(0) = 1 - p_1$ and $\pi(1) = p_1$. Furthermore, the packet erasure correlation ρ provides an average measure of the correlation of two consecutive packets. In particular, when $\rho = 0$, $p_{01} + p_{10} = 1$, the loss process is memory-less, and the above probability measures reduce to the special case of a memory-less Binary Erasure Channel (BEC). In the sequel, we analyze the impact of the level of correlation among consecutive packets, as represented by ρ, on ISC-based packet video over a wide range of loss rate p_1 values.

4. SIMULATIONS AND RESULTS

For simulation, CIF sequences, *Akiyo*, *Foreman*, *Coastguard*, and *Mobile*, are encoded using MPEG-4 encoder. The un-interleaved GOV sizes of 10, 12, 14, 16, 18, and 20 were used and frame rate of *15fps*, bitrate of *250kbps* and *500kbps*, and packet size of *512Byte* were used when coding the sequences. 5%, 10%, and 15% packet loss rates were used and packet correlation value of 0.3, 0.6, and 0.9 were used to represent low, medium, and high correlation between the transmitted packets.

To limit the impact of a single packet loss to a single frame, no packets are shared among two consecutive coded frames. Furthermore, partial decoding is not allowed for the frames with errors and they are replaced with the last successfully decoded frames (frozen frames) for both ISC and traditional (non-ISC) cases. To obtain statistically viable results and to capture realistic network loss patterns, ten error traces were generated for each $p_1 - \rho$ pair. Each ISC pattern is fitted into these error traces and the PSNR values are averaged to provide statistically satisfying results for analysis.

As shown in Fig 1. and Fig 2., ISC shows improvements on most of the evaluation cases. It is clearly seen that ISC advances the traditional method as the channel loss rate increases or the packet correlation rate decreases. This is due to the fact that ISC reduces impact of packet losses to the GOV by isolating errors to

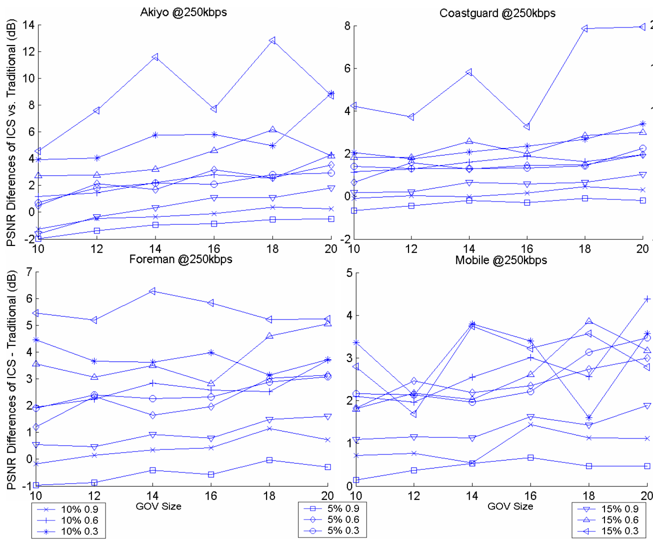

Fig 1. PSNR Differences between ISC and Non-ISC @*250kbps*

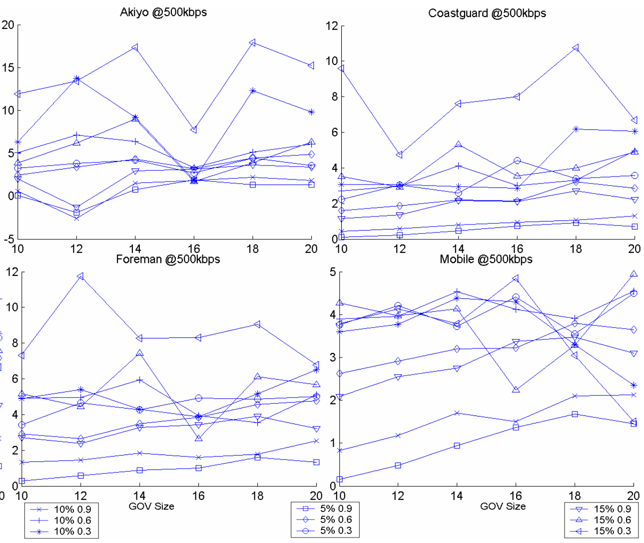

Fig 2. PSNR Differences between ISC and Non-ISC @*250kbps.*

one of the two sub-sequences and decreases frame replacement distances for decoder failed frames. However, with the increment of the packet correlation constants, the frequency of the long error bursts increases, and hence the chance that both sub-sequences are impacted by the long error bursts increases.

5. CONCLUSION

In this paper, we have evaluated the performance of an *interleaved source coding* (ISC) method for packet video over channels with memory. ISC is resilient to packet losses since it limits the errors due to packet losses to one of two sub-sequences (generated by ISC) and minimizes cascaded effects of packet losses over a *single* erasure-channel model and increases the number of successfully decoded frames and overall playback quality of the decoded video sequence. It is clearly shown that ISC advances traditional method for most of the cases; however, the performance highly depends on the channel's loss rate and the packet correlation increments. Since most of the packet losses over the best-effort network channel are caused by the buffer overflow of the routers on the path, hence the high frequency, short burst errors are more likely to be observed, therefore, adopting ISC to the realtime streaming services over the best-effort network would be beneficial.

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