

Embedded System Design

Modeling, Synthesis, Verification

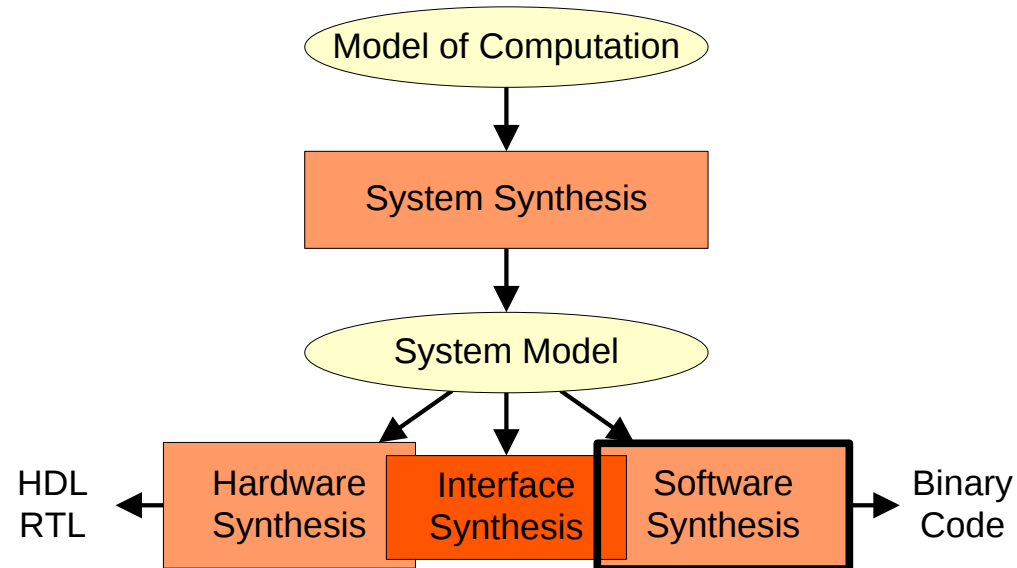
Daniel D. Gajski, Samar Abdi, Andreas Gerstlauer, Gunar Schirner



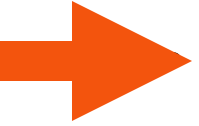
Chapter 5: Software Synthesis

SW Synthesis within Overall Design Flow

- **MoC**
 - Capture application
- **System Synthesis**
 - Define system wide decisions
- **System model**
 - Estimate / analyze system wide decisions
- **Component synthesis**
 - Generate component implementation
 - Software Synthesis
Produce SW binary for exec on processor



Outline



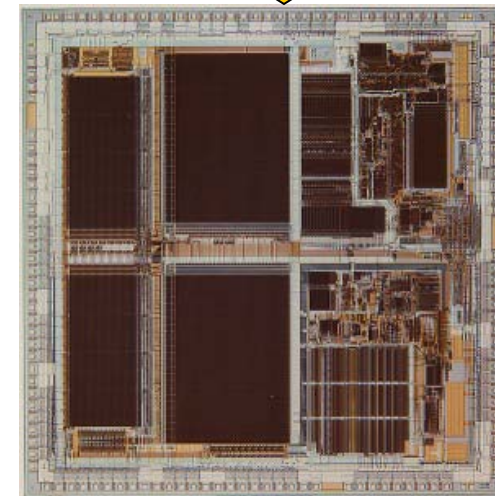
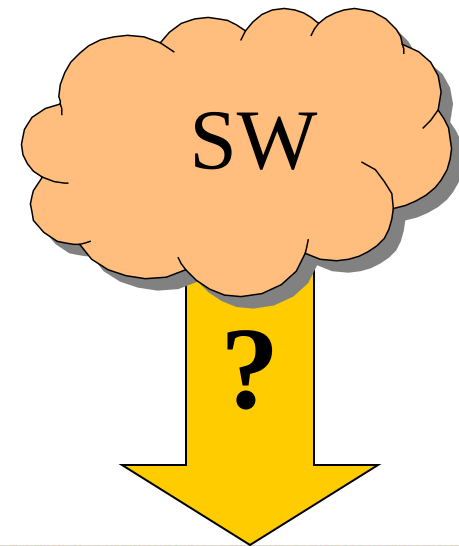
Introduction

- Preliminaries
- Software Synthesis Overview
- **Code Generation**
- **Hardware-dependent Software**
 - Multi-Task Synthesis
 - Internal Communication
 - External Communication
 - Startup Code Generation
 - Binary Image Generation
- **Execution**
- **Summary**



Motivation

- **Increasing complexity of Multi-Processor System-on-Chip (MPSoCs)**
 - Feature demands
 - Production capabilities + Implementation freedom
- **Increasing software content**
 - Flexible solution
 - Addresses complexity
- **How to create SW for MPSoC?**
 - Avoid break in ESL flow:
 - Synthesize SW from abstract models

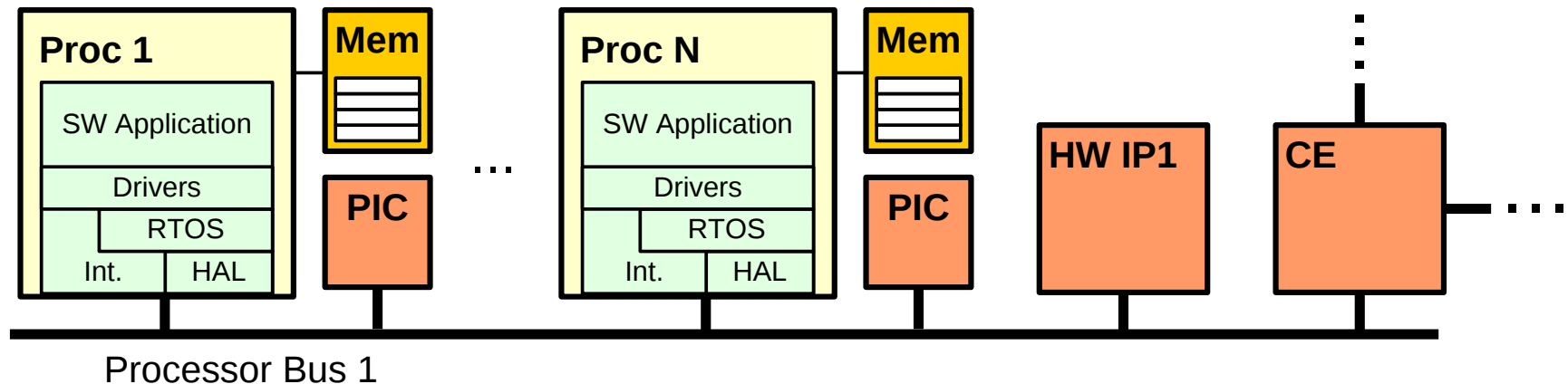


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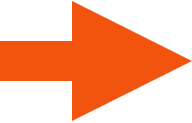


Goals

- **Generate SW binaries for MPSoC from abstract specification**
 - Eliminate tedious, error prone SW coding
 - Rapid exploration
 - High-level application development
 - Support wide range of system sizes
 - with RTOS / without RTOS (i.e. interrupt-driven)



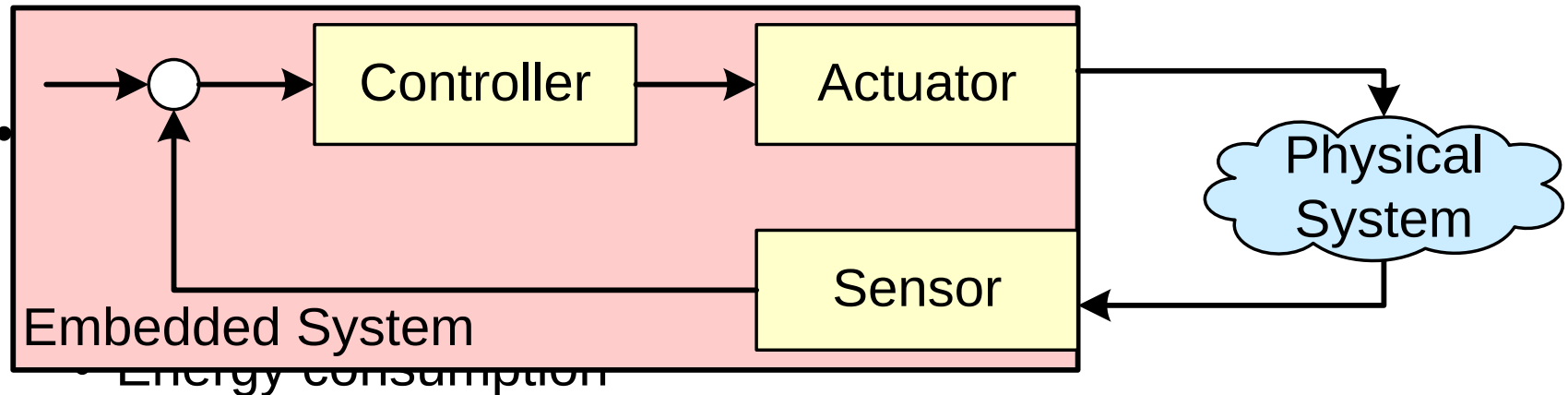
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Embedded Software Challenges

- **SW tightly coupled**
 - Underlying HW (accelerator)
 - Part of a physical control loop
- **Time constrained**
 - Correctness depends on time frame
 - Too late functional correct answer still wrong
- **Concurrent**



Example Target Languages

- **Assembly**
 - Basically a symbolic representation of machine code
 - Fine grained control: Registers, Instructions
- **C**
 - General purpose programming language, including level features
 - Bit operations, direct memory management
 - Low run-time overhead
 - Coding standards for portability
 - E.g. MISRA (automotive)
- **C++**
 - Backward compatible to C
 - Facilitates object oriented programming
- **Java**
 - Portability!
 - Hides memory management, pointer arithmetic
 - Interpreted (JVM), ahead of time and Just in Time (JIT) compiled
 - Java accelerators (e.g. ARM Jazelle), Java processors
 - Real-time extension: Real-time Specification for Java (RTSJ)



Real-Time Operating System (RTOS)

- **Similar to general purpose operating system**
 - SW layer above HAL
 - Provides services for
 - Concurrent execution
 - Communication
 - Synchronization
 - Resource management
- **Real-Time Operating System**
 - **Facilitates** constructing real-time systems
 - Predictable response time
 - Hard real-time
 - Catastrophic failure on deadline miss
 - Soft real-time
 - Missing some deadlines tolerable



Real-Time Operating System (RTOS)

- Characteristics of scheduling policies
 - Preemptive / Non-preemptive
 - Static / Dynamic
 - Off-line / On-line
- Example scheduling policies
 - Priority-based
 - Earliest Deadline First (EDF)
 - Rate Monotonic (RM)



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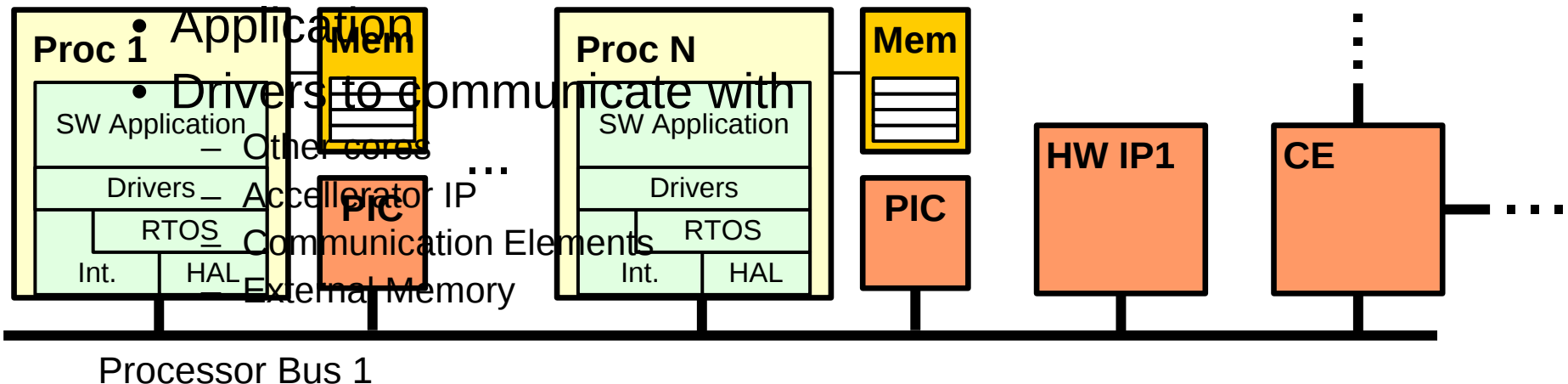


Software Synthesis Overview

- **Multi-core target platform**

- Processors
 - PIC, Timer, MEM
- HW IP
- Communication Element
 - E.g. Transducer
- Bus hierarchy

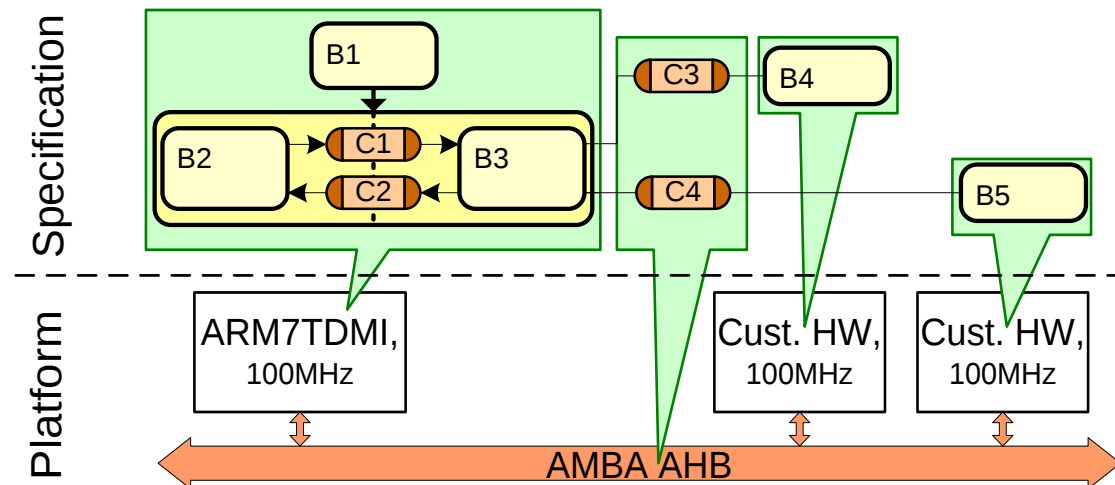
- **Generate for each core**



System Design Flow Overview

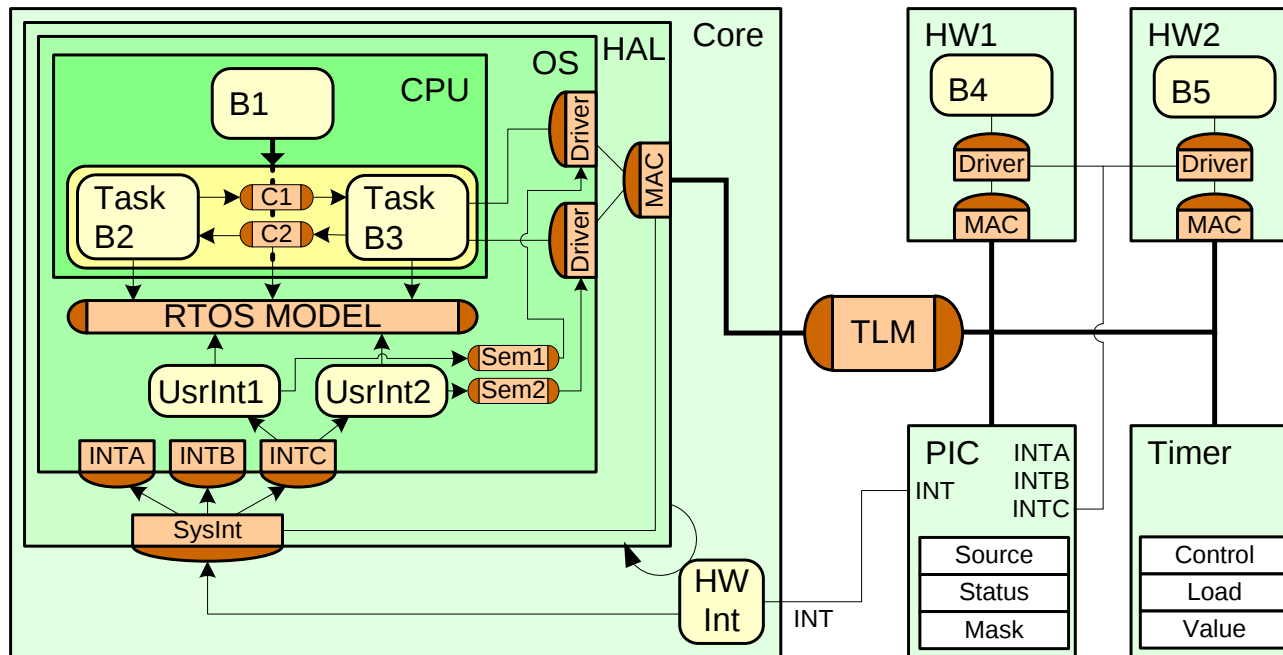
Input specification (review)

- **Capture Application in MoC, then SLDL**
 - Computation
 - Organize code in behaviors (processes)
 - Communication
 - Point-to-point channels, feature-rich
 - » Synchronous / Asynchronous
 - » Blocking / Non-Blocking
 - » Synchronization only (e.g. semaphore, mutex, barrier)
- **Architecture decisions:**
 - Processor(s)
 - HW component(s)
 - Busses
 - Mapping
 - ...



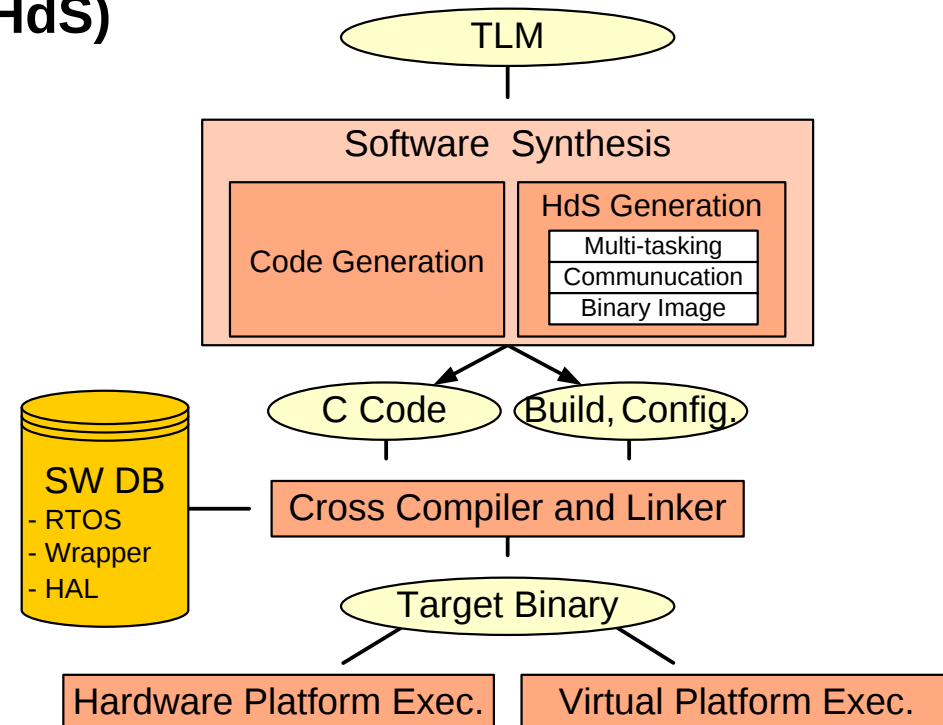
System Design Flow Overview

- **System synthesis output: system TLM**
 - Functional timed abstract model
 - Reflects system-wide decisions



Software Synthesis Overview

- **Code Generation**
 - Generate application code inside each task
 - Resolve behavioral hierarchy into flat C code
- **Hardware-dependent Software (HdS) Generation**
 - Multi-task Generation
 - Tasks -> RTOS
 - Communication Synthesis
 - Drivers
 - Binary Image Generation
 - Cross Compile and Link
- **Binary executable on**
 - Actual Hardware Platform
 - Virtual Platform with integrated ISS



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Code Generation

- **Rational:**
 - TLM in SLDL (e.g. SystemC, SpecC)
 - SLDL system level optimized, not target
 - Generic for any architecture
 - » Hierarchical composition of concurrent modules
 - » Connectivity between modules
 - » Communication encapsulation
 - » Hardware concepts
 - Inefficient direct compilation to target
 - Large simulation kernel on target
 - Need efficient (footprint, memory, speed) target implementation
 - Translate system model to target language
 - SystemC -> ANSI-C



Code Generation

- **Generate sequential code executing in a task**
 - Remove SLDL specific concepts
 - Resolve hierarchy and connectivity into flat C code
 - Set of C *structs*: Behaviors / channels
 - » Member variables + Port connectivity
 - Global functions + context pointer (*pThis)

Rules for C code generation

1. Module resolution

Behaviors and channels are converted into C *struct*

2. Hierarchy translation

Child behaviors and channels are instantiated as C *struct* members inside the parent C *struct*

3. Variable translation

Variables defined inside a behavior or channel are converted into data members of the corresponding C *struct*

4. Port resolution

Ports of behavior or channel are converted into data members of the corresponding C *struct*

5. Method globalization

Functions inside a behavior or channel are converted into global functions

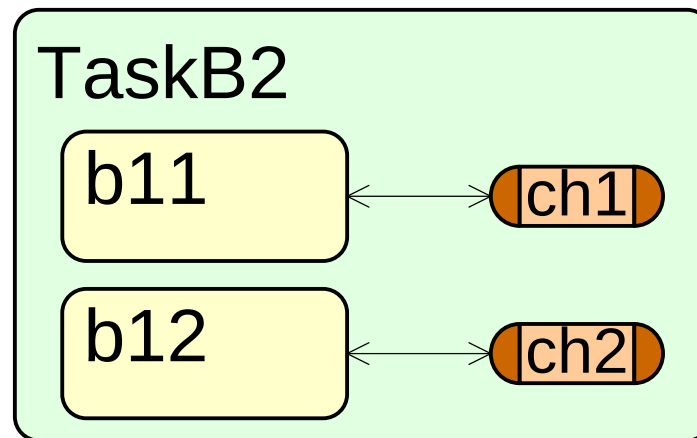
6. Global context creation

A static *struct* instantiation for each PE is added to allocate/initialize the data used by SW



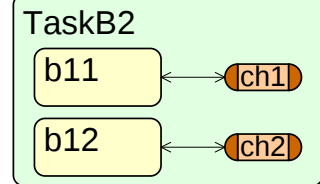
Code Generation

- **Example**
 - Sequential executing b11, b12
 - Instances of channel class B1
 - Port bound to ch1, ch2
 - Instances of channel class CH1



Code Generation

1) Module resolution



SystemC task specification

```
1 SC_MODULE(B1){
2     int A;
3     sc_port<iChannel> myCh;
4     SC_CTOR(B1){}
5     void main(void) {
6         A = 1;
7         myCh->chCall(A*2);
8     }
9 };
10
11 SC_MODULE(TaskB2){
12     CH1 ch11, ch12; //channel inst
13     B1 b11, b12;    //module inst
14     SC_CTOR(TaskB2):
15         ch11("ch11"), ch12("ch12"),
16         b11("b11"), b12("b12") {
17         b11.myCh(ch11); // connect ch11
18         b12.myCh(ch12); // connect ch12
19     }
20     void main(void) {
21         b11.main();
22         b12.main();
23     }
24 };
```

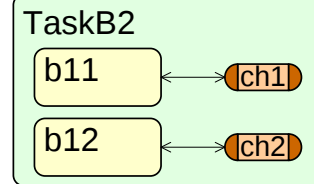
ANSI-C task code

```
1 struct B1 {
2     struct CH1 *myCh; /* port iChannel*/
3     int A;
4 };
5 struct TaskB2 {
6     struct B1 b11, b12;
7     struct CH1 ch11, ch12;
8 };
9 void B1_main(struct B1 *This) {
10     (This->A) = 1;
11     CH1_chCall(This->myCh, (This->A)*2);
12 }
13 void TaskB2_main(struct TaskB2 *This){
14     B1_main(&(This->b11));
15     B1_main(&(This->b12));
16 }
17 struct TaskB2 taskB2= {
18     {&(taskB2.ch11),0/*A*/}/*b11*/,
19     {&(taskB2.ch12),0/*A*/}/*b12*/,
20     {} /*ch11*/, {} /*ch12*/
21 };
22 void TaskB2() {
23     TaskB2_main( &task1);
24 }
```



Code Generation

2) Hierarchy translation



SystemC task specification

```
1 SC_MODULE(B1){
2     int A;
3     sc_port<iChannel> myCh;
4     SC_CTOR(B1){}
5     void main(void) {
6         A = 1;
7         myCh->chCall(A*2);
8     }
9 };
10
11 SC_MODULE(TaskB2){
12     CH1 ch11, ch12; //channel inst
13     B1 b11, b12;    //module inst
14     SC_CTOR(TaskB2):
15         ch11("ch11"), ch12("ch12"),
16         b11("b11"), b12("b12") {
17         b11.myCh(ch11); // connect ch11
18         b12.myCh(ch12); // connect ch12
19     }
20     void main(void) {
21         b11.main();
22         b12.main();
23     }
24 };
```

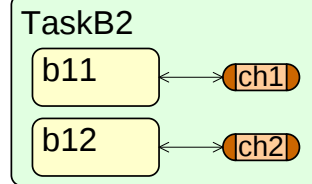
ANSI-C task code

```
1 struct B1 {
2     struct CH1 *myCh; /* port iChannel*/
3     int A;
4 };
5 struct TaskB2 {
6     struct B1 b11, b12;
7     struct CH1 ch11, ch12;
8 };
9 void B1_main(struct B1 *This) {
10     (This->A) = 1;
11     CH1_chCall(This->myCh, (This->A)*2);
12 }
13 void TaskB2_main(struct TaskB2 *This){
14     B1_main(&(This->b11));
15     B1_main(&(This->b12));
16 }
17 struct TaskB2 taskB2= {
18     {&(taskB2.ch11),0/*A*/}/*b11*/,
19     {&(taskB2.ch12),0/*A*/}/*b12*/,
20     {} /*ch11*/, {} /*ch12*/
21 };
22 void TaskB2() {
23     TaskB2_main( &task1);
24 }
```



Code Generation

3) Variable translation



SystemC task specification

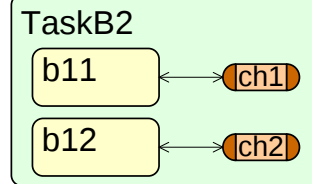
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1 SC_MODULE(B1){
2   int A;
3   sc_port<iChannel> myCh;
4   SC_CTOR(B1){}
5   void main(void) {
6     A = 1;
7     myCh->chCall(A*2);
8   }
9 };
10
11 SC_MODULE(TaskB2){
12   CH1 ch11, ch12; //channel inst
13   B1 b11, b12;    //module inst
14   SC_CTOR(TaskB2):
15     ch11("ch11"), ch12("ch12"),
16     b11("b11"), b12("b12") {
17     b11.myCh(ch11); // connect ch11
18     b12.myCh(ch12); // connect ch12
19   }
20   void main(void) {
21     b11.main();
22     b12.main();
23   }
24 };
```

ANSI-C task code

```
1 struct B1 {
2   struct CH1 *myCh; /* port iChannel*/
3   int A;
4 };
5 struct TaskB2 {
6   struct B1 b11, b12;
7   struct CH1 ch11, ch12;
8 };
9 void B1_main(struct B1 *This) {
10   (This->A) = 1;
11   CH1_chCall(This->myCh, (This->A)*2);
12 }
13 void TaskB2_main(struct TaskB2 *This){
14   B1_main(&(This->b11));
15   B1_main(&(This->b12));
16 }
17 struct TaskB2 taskB2= {
18   {&(taskB2.ch11),0/*A*/}/*b11*/,
19   {&(taskB2.ch12),0/*A*/}/*b12*/,
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21 };
22 void TaskB2() {
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24 }
```



4) Port resolution



SystemC task specification

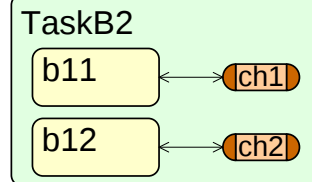
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2   int A;
3   sc_port<iChannel> myCh;
4   SC_CTOR(B1){}
5   void main(void) {
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8   }
9 };
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11 SC_MODULE(TaskB2){
12   CH1 ch11, ch12; //channel inst
13   B1 b11, b12;    //module inst
14   SC_CTOR(TaskB2):
15       ch11("ch11"), ch12("ch12"),
16       b11("b11"), b12("b12") {
17       b11.myCh(ch11); // connect ch11
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ANSI-C task code

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1 struct B1 {
2   struct CH1 *myCh; /* port iChannel*/
3   int A;
4 };
5 struct TaskB2 {
6   struct B1 b11, b12;
7   struct CH1 ch11, ch12;
8 };
9 void B1_main(struct B1 *This) {
10   (This->A) = 1;
11   CH1_chCall(This->myCh, (This->A)*2);
12 }
13 void TaskB2_main(struct TaskB2 *This){
14   B1_main(&(This->b11));
15   B1_main(&(This->b12));
16 }
17 struct TaskB2 taskB2= {
18   {&(taskB2.ch11),0/*A*/}/*b11*/,
19   {&(taskB2.ch12),0/*A*/}/*b12*/,
20   {} /*ch11*/, {} /*ch12*/
21 };
22 void TaskB2() {
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```



5) Method globalization



SystemC task specification

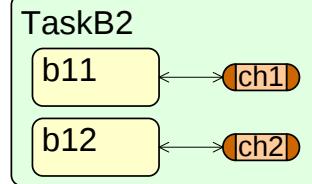
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4     SC_CTOR(B1){}
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11 SC_MODULE(TaskB2){
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ANSI-C task code

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3     int A;
4 };
5 struct TaskB2 {
6     struct B1 b11, b12;
7     struct CH1 ch11, ch12;
8 };
9 void B1_main(struct B1 *This) {
10     (This->A) = 1;
11     CH1_chCall(This->myCh, (This->A)*2);
12 }
13 void TaskB2_main(struct TaskB2 *This){
14     B1_main(&(This->b11));
15     B1_main(&(This->b12));
16 }
17 struct TaskB2 taskB2= {
18     {&(taskB2.ch11),0/*A*/}/*b11*/,
19     {&(taskB2.ch12),0/*A*/}/*b12*/,
20     {} /*ch11*/, {} /*ch12*/
21 };
22 void TaskB2() {
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```



6) Global context creation



SystemC task specification

```
1 SC_MODULE(B1){
2   int A;
3   sc_port<iChannel> myCh;
4   SC_CTOR(B1){}
5   void main(void) {
6     A = 1;
7     myCh->chCall(A*2);
8   }
9 };
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11 SC_MODULE(TaskB2){
12   CH1 ch11, ch12; //channel inst
13   B1 b11, b12;    //module inst
14   SC_CTOR(TaskB2):
15     ch11("ch11"), ch12("ch12"),
16     b11("b11"), b12("b12") {
17     b11.myCh(ch11); // connect ch11
18     b12.myCh(ch12); // connect ch12
19   }
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23   }
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```

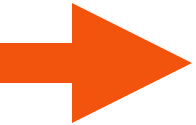
ANSI-C task code

```
1 struct B1 {
2   struct CH1 *myCh; /* port iChannel*/
3   int A;
4 };
5 struct TaskB2 {
6   struct B1 b11, b12;
7   struct CH1 ch11, ch12;
8 };
9 void B1_main(struct B1 *This) {
10   (This->A) = 1;
11   CH1_chCall(This->myCh, (This->A)*2);
12 }
13 void TaskB2_main(struct TaskB2 *This){
14   B1_main(&(This->b11));
15   B1_main(&(This->b12));
16 }
17 struct TaskB2 taskB2= {
18   {&(taskB2.ch11),0/*A*/}/*b11*/,
19   {&(taskB2.ch12),0/*A*/}/*b12*/,
20   {} /*ch11*/, {} /*ch12*/
21 };
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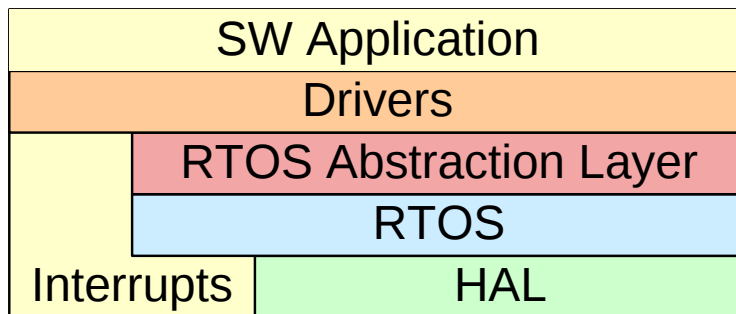


Multi-Task Synthesis

- **Variants**

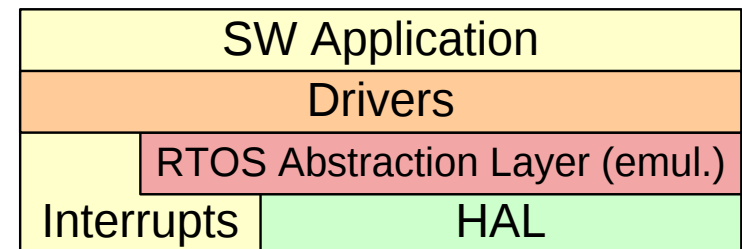
- **RTOS-based**

- General solution
 - Off-the-shelf RTOS
 - Flexible, well tested



- **Interrupt-based**

- Few reactive tasks
 - Resource constraints inhibit RTOS
 - Unavailability of RTOS
 - E.g. DSP with encoder / decoder application



RTOS-based Multi-Tasking

- **Based on off the shelf RTOS**
 - e.g. μ C/OS-II, eCos, vxWorks
- **Uses RTOS Abstraction Layer (RAL)**
 - Canonical interface
 - Limits interdependency RTOS / Synthesis
- **Multi-task synthesis**
 - Generate task management code
 - Adjust internal communication

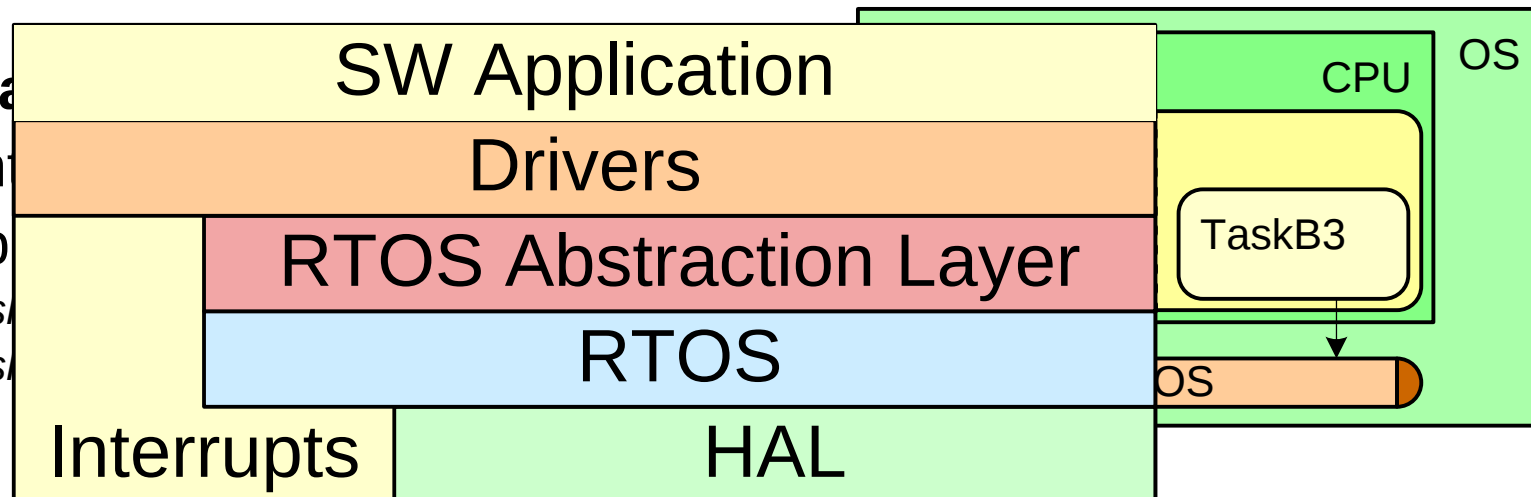
- **RTOS-based**

- Parent

- Two p

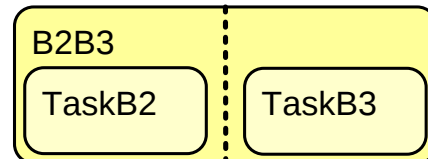
- Task

- Task



Multi-Task Synthesis

- RTOS-based multi-tasking example



```
1 SC_MODULE(B2B3) {
2   public:
3     sc_port<iRTOS> rtos;
4     TaskB2 taskB2;
5     TaskB3 taskB3;
6     SC_CTOR(B2B3):
7       taskB2("taskB2", 5, 4096),
8       taskB3("taskB3", 2, 4096) {
9         taskB2.rtos(rtos);
10        taskB3.rtos(rtos);
11      }
12      void main(void) {
13        taskB2.release();
14        taskB3.release();
15        taskB2.join();
16        taskB3.join();
17      }
18 };
```

SystemC

```
1 struct B2B3{
2   struct TaskB2 task_b2;
3   struct TaskB3 task_b3;};
4 void *TaskB2_main(void *arg){
5   struct TaskB2 *this=(struct TaskB2*)arg;
6   /* ...*/
7 }
8 void *TaskB3_main(void *arg){
9   struct TaskB3 *this=(struct TaskB3*)arg;
10  /* ... */
11 }
12 void *B2B3_main(void *arg){
13   struct B2B3 *this= (struct B2B3*)arg;
14   os_task_handle task_b2, task_b3;
15   task_b2 = taskCreate(TaskB2_main,
16                       &this->taskB2, 5, 4096);
17   task_b3 = taskCreate(TaskB3_main,
18                       &this->taskB3, 2, 4096);
19
20   taskJoin(task_b2);
21   taskJoin(task_b3);
22 }
```

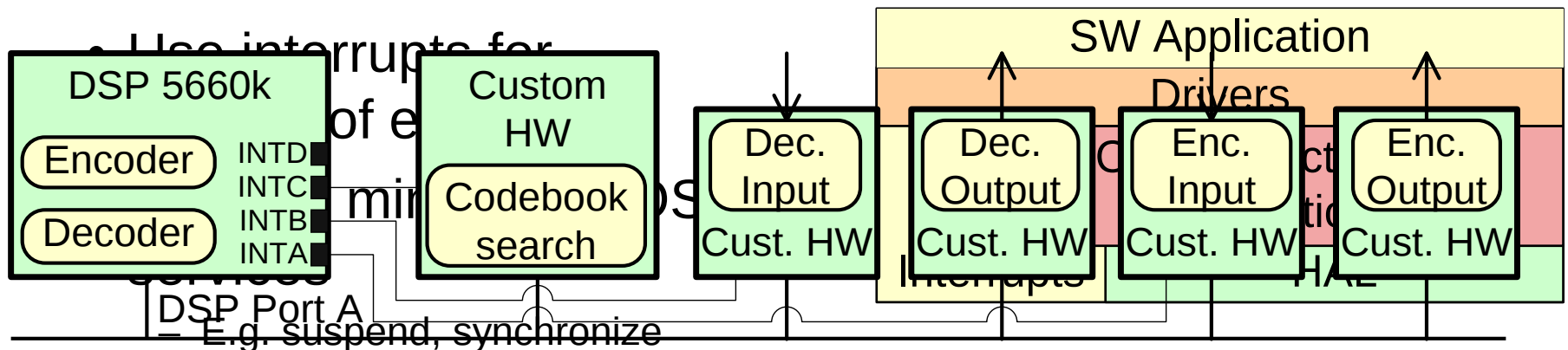
ANSI-C



Interrupt-based multi-tasking

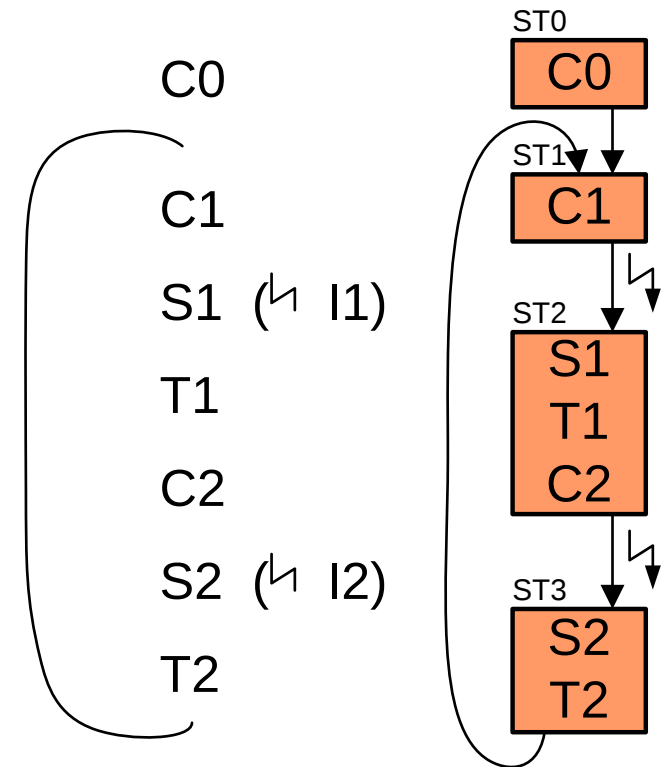
- **Alternative multi-tasking**
 - Few tasks
 - Resource constraints
 - Memory, footprint, performance
 - Unavailability of RTOS port

- **Overview**



Interrupt-based multi-tasking

- **Procedure overview**
 - Convert tasks to state machine
 - Execute state machine in interrupt handler
 - Execute lowest priority task as main()
- **Assume only interrupt synchronization**
- **Code composed of**
 - Computation C_n
 - Synchronization S_n ,
 - Data transfer T_n
 - Interrupt sych. I_n
- **Convert task code to state machine**
 - New state for each synchronization
 - E.g. ST2, ST3
 - New state for each conditional
 - E.g. Loop: ST0, ST1
- **Execute state machine in interrupt**



Multi-Task Synthesis

- **Interrupt-based multi-tasking example**

- **Example execution**

- Assume computation C1 [line 10] just finished

- New state ST2

- Attempt S1

- » Unavailable

- » ISR terminates

- Receive I1

- Release S1

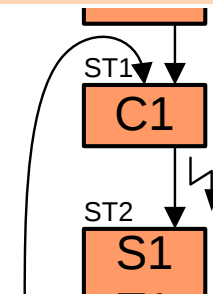
- executeTask0()

- Continue ST2

- Attempt S1

- » Available

- Continue

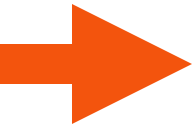


```
1  /* interrupt handler */
2  void intHandler_I1() {
3      release(S1);    /* set S1 ready */
4      executeTask0(); /* task state machine */
5  }
6  /* task state machine */
7  void executeTask0() {
8      do { switch(Task0.State) {
9          /* ... */
10         case ST1: C1(...);
11                 Task0.State = ST2;
12         case ST2: if(attempt(S1))
13                   T1_receive(...);
14                   else break;
15                   C2(...);
16                   Task0.State = ST3;
17         case ST3: /* ... */
18         } } while (Task0.State == ST1);
19 }
```



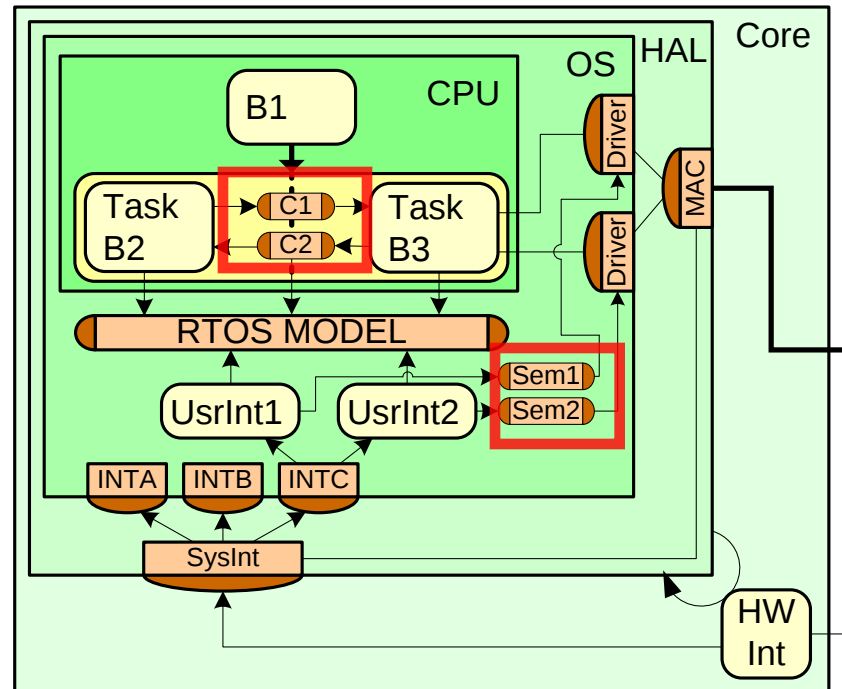
Outline

- **Introduction**
 - Preliminaries
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- **Code Generation**
- **Hardware-dependent Software**
 - Multi-Task Synthesis
 - Internal Communication
 - External Communication
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 - Binary Image Generation
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Internal Communication

- **Communication within processor**
 - Replace with target specific implementation
 - e.g. RTOS semaphore, event, msg. queue



Internal Communcation

- **Implementation example**

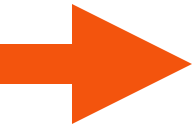
- Single handshake
 - One way synchronous without data
- Xilkernel (Xilinx proprietary RTOS) implementation
 - Semaphore-based

```
1  /** SHS OS-specific struct */
2  typedef struct {
3      sem_t req; /** os semaphore */
4  } tCh_shs;
5  void Shs__init(tCh_shs *pThis){
6      int retVal = sem_init(&pThis->req, 0, 0);
7      /* ... error handling */
8  }
9  void Shs_send(tCh_shs *pThis){
10     int retVal = sem_post(&pThis->req);
11     /* ... error handling */
12 }
13 void Shs_receive(tCh_shs *pThis){
14     int retVal = sem_wait(&pThis->req);
15     /* ... error handling */
16 }
```



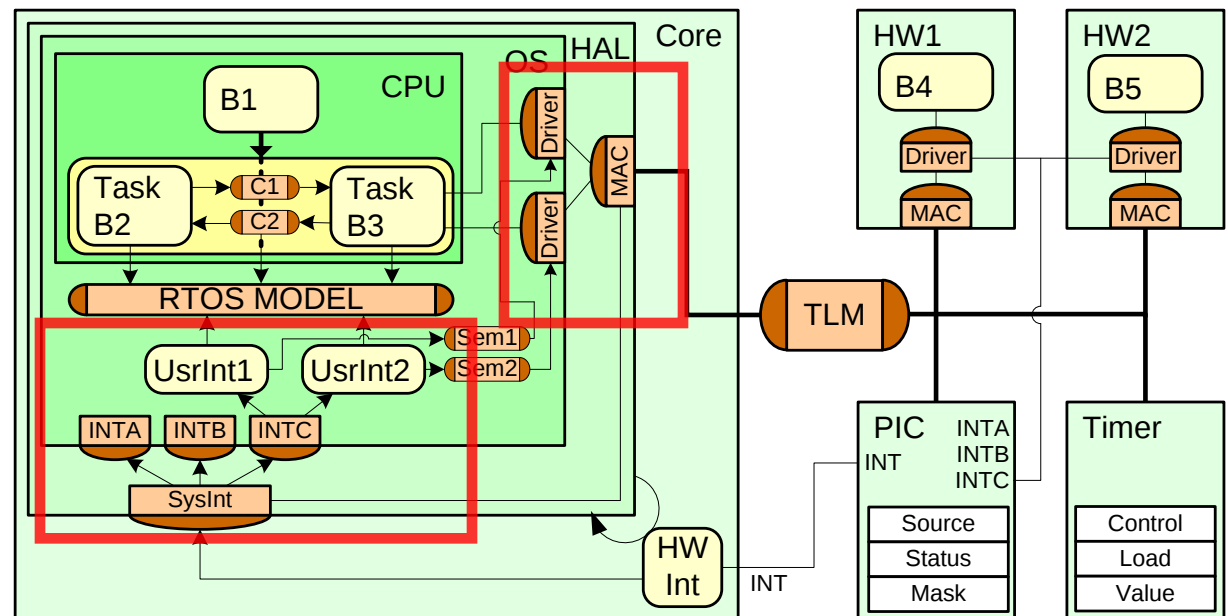
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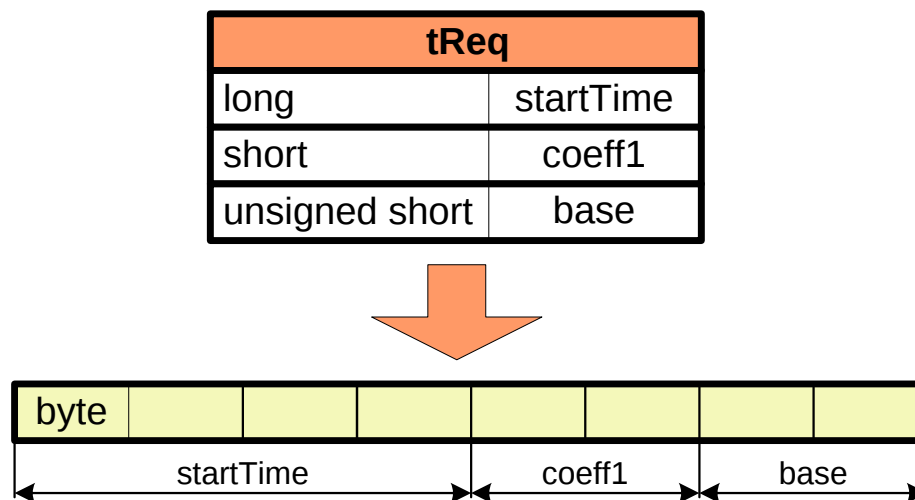
External Communication

- **Communication outside of processor**
 - E.g. with HW Accelerator
- **ISO / OSI layered to support heterogeneous architectures**
 - Data formatting
 - Packetization
 - Synchronization
 - MAC



External Communication

- **Data formatting (marshalling)**
 - Problem: different memory representation for same data
 - Endianness (byte order)
 - Packing rules
 - Bit widths for data types
 - Convert typed data into flat untyped byte stream
 - Interpretable by everybody



External Communication

- **Data formatting example:**
 - Marshalling code uses
 - `hton<datatype>`
 - `ntoh<datatype>`

User type definition

```
1 typedef struct stReq {
2     long      startTime;
3     short     coeff1;
4     unsigned short base;
5 } tReq;
```

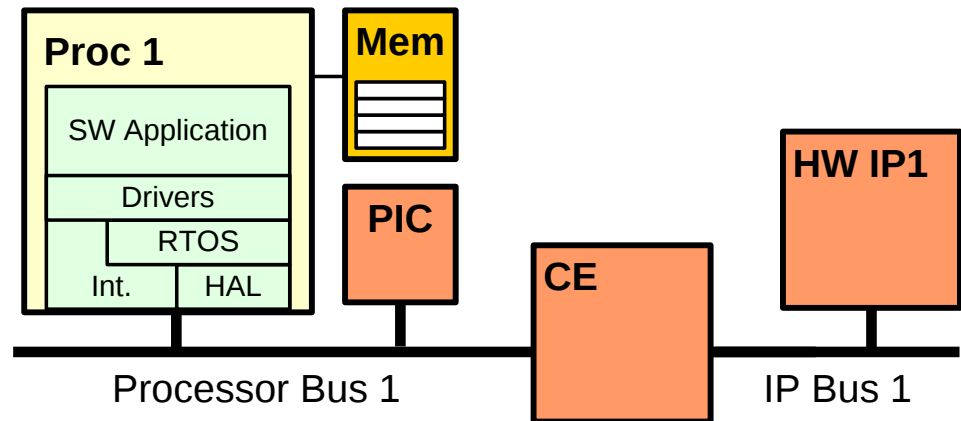
Marshalling code

```
1 void myCh_send(/* ...*/ *This,
2               struct tReq *pD){
3     unsigned char *pB = This->buf;
4     htonl(pB, pD->startTime);
5     pB += 4;
6     htons(pB, pD->coeff1);
7     pB += 2;
8     htons(pB, pD->base);
9     pB += 2;
10    DLink0_trans_send(/*...*/This->buf, 8);
11 }
```



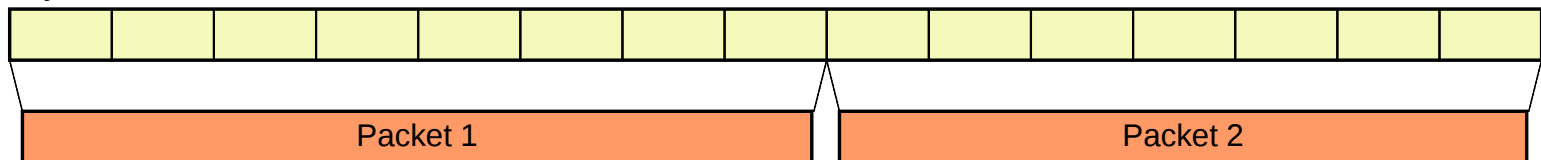
External Communication

- Packetization



- Break user defined length stream into packets
- Limit storage requirements during transport (CEs)

Byte stream



External Communication

- **Packetization example code**
 - Transmits input *pMsg*, of length *len* in packets of up to *CONFIG_PACKET_SIZE*

```
1 DLink0_trans_send(void *pMsg, unsigned int len){
2     unsigned char *pPos = pMsg;
3     while(len) {
4         unsigned long pktLen;
5         /* length is minimum of max size and len */
6         pktLen = min(len, CONFIG_PACKET_SIZE);
7
8         DLink0_net_send(pPos, pktLen); /* transfer */
9
10        len -= pktLen; /* decr. transferred len */
11        pPos += pktLen; /* advance pointer */
12    }
13 }
```

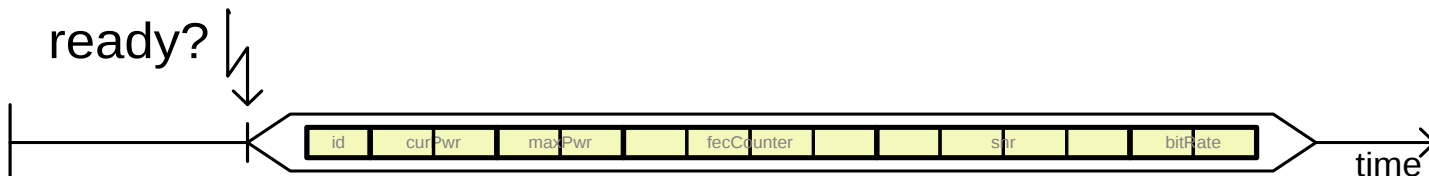


External Communication

- **Synchronization**

- Ensure slave is ready before master initiates transaction

- Ready to receive data / Data ready



- Master / slave bus: Separate event
 - Node-based bus: Synchronization packet

- Processor synchronization options:

- Interrupt
 - » Separate interrupt line connected to interrupt controller
 - » Low latency
 - » Interrupt overhead
 - » Shared interrupts
 - Polling
 - » Master periodically checks slave using data connection
 - » Polling period
 - Hybrid solutions possible



External Communication

- **Synchronization by interrupt**

- 1) Low level interrupt handler

- Preempts current task

- 2) System interrupt handler

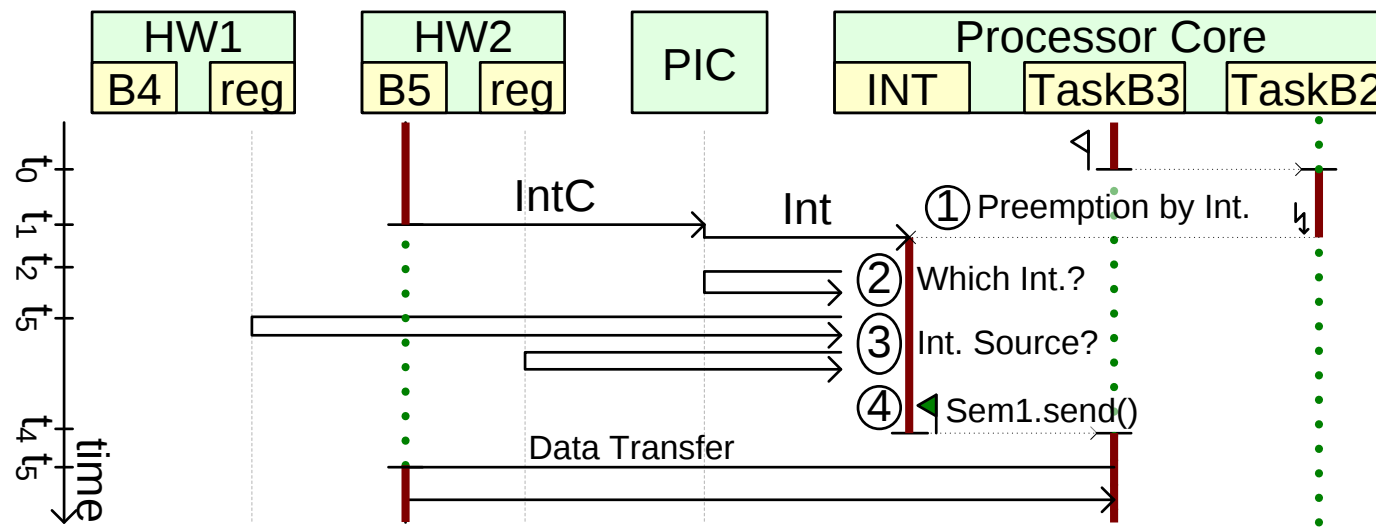
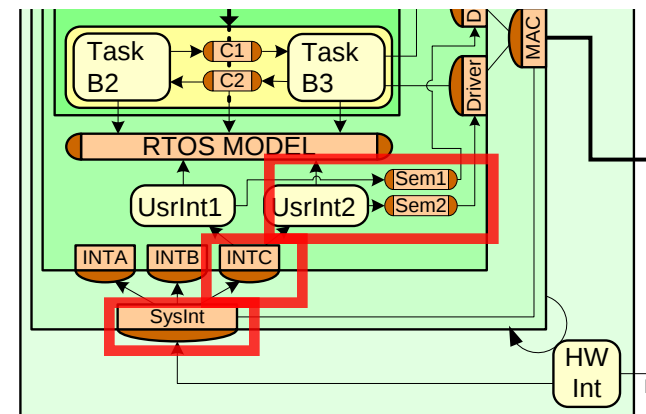
- Checks PIC

- 3) User-specific interrupt handler

- Handles shared interrupts

- 4) Semaphore

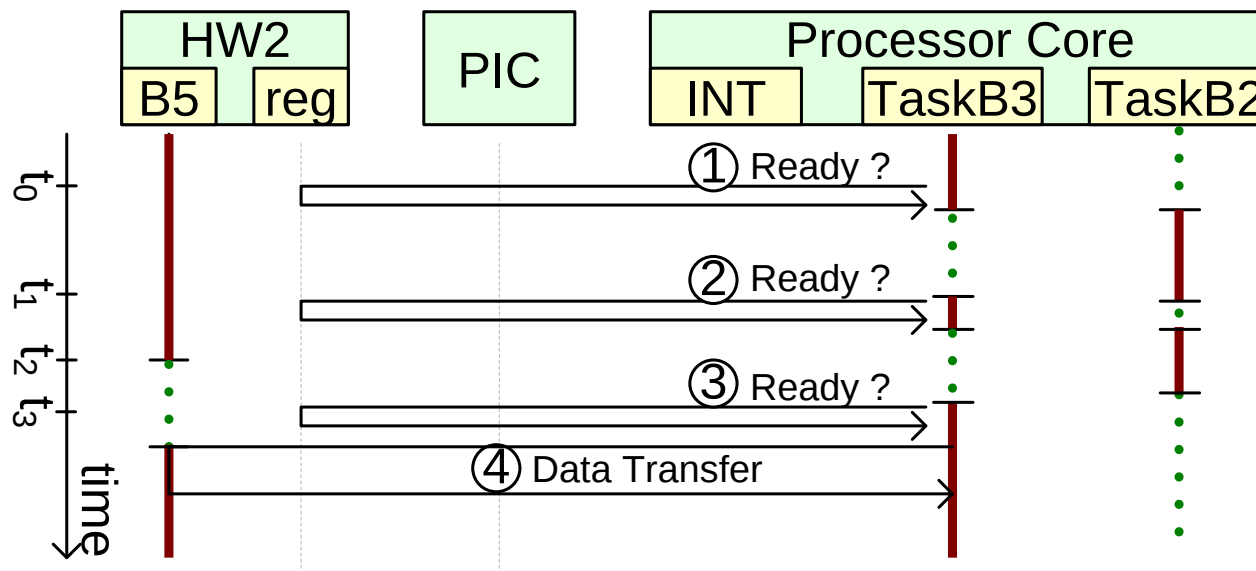
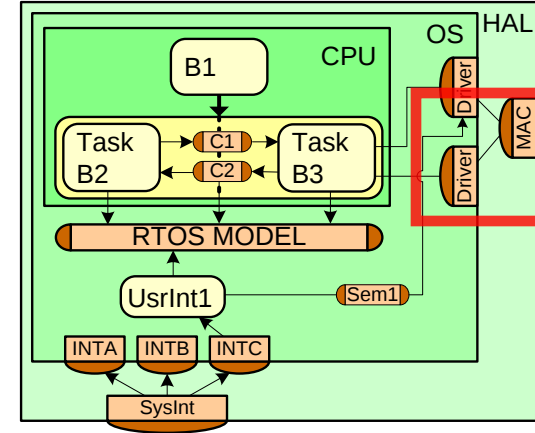
- Releases task



External Synchronization

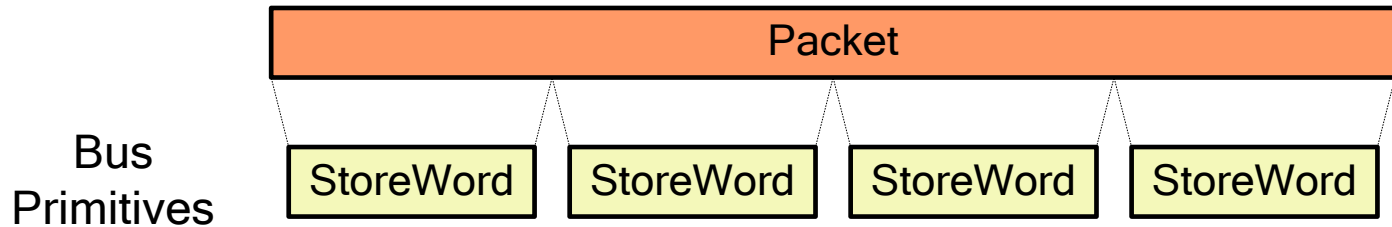
- **Polling example**

- 1) Expect message, check ready?
 - Suspend for polling period
- 2) Check ready?
 - Suspend for polling period
- 3) Check ready?
 - Finish synchronization
- 4) Data transfer



External Communication

- **Media Access Control (MAC)**
 - Provides access to bus medium
 - Low level driver
 - Break packet into bus transactions



- Simple drivers use processor memory interface
 - Memory mapped bus interface
- More complex communicate / synchronize with protocol controller
 - I2C, CAN, FlexRay



External Communication

- **Media Access Control (MAC)**
 - Example for memory mapped bus access
 - Left: Cast integer *addr* to pointer, take value of
 - Right: Cast void **pD*, to int pointer, get value
 - Assign value to “value of” results in bus transaction
 - Repeat for word, short, byte

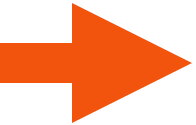
```
1 void masterWrite(unsigned int addr, void *pData, unsigned int len) {
2     unsigned char *pD = (unsigned char*)pData;
3     while (len >= 4 ) {
4         *((unsigned int*)addr) = *((unsigned int*)pD);
5         len -= 4; pD += 4;
6     }
7     if (len >= 2 ) {/* remaning short */
8         *((unsigned short*)addr) = *((unsigned short*)pD);
9         len -= 2; pD += 2;
10    }
11    if (len >= 1) {/* the last byte */
12        *((unsigned char*)addr) = *((unsigned char*)pD);
13        len -= 1; pD += 1;
14    }
15 }
```

Note: code assumes 32bit processor.



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Startup Code

- **Initialize system and release tasks**

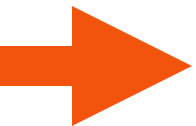
- Global data structure
- Board Support Package (BSP)
- Operating System
- Create synchronization channels
- Register interrupts
- Create and release user tasks

```
1 void main(void) {  
2     PE_Struct_Init(&PE0);  
3     BSP_init();  
4     OSInit();  
5  
6     c_os_handshake_init(&PE0->sem1);  
7     c_os_handshake_init(&PE0->sem2);  
8     BSP_UserIrqRegister(INT1, Int1Handler, /*...*/);  
9     BSP_UserIrqRegister(INT2, Int2Handler, /*...*/);  
10  
11     taskCreate(task_b2b3, NULL,  
12                B2B3_main, &this->task_b2b3);  
13  
14     OSStart();  
15 }
```



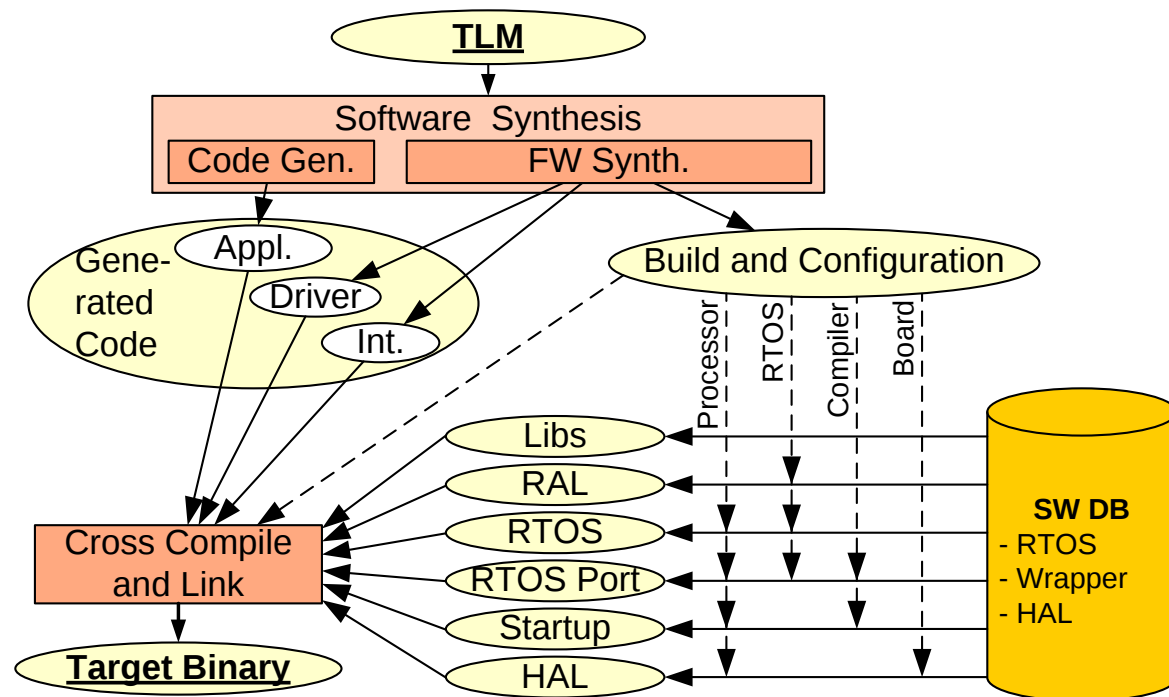
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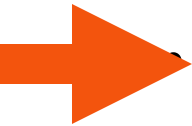
Binary Image Generation

- **Generate binary for each processor**
 - Generate build and configuration files
 - Select software database components
 - Configure RTOS
 - Cross compile and link
- Significant effort in DB design
 - Minimize components
 - Analyze dependencies
 - Goal: flexible composition



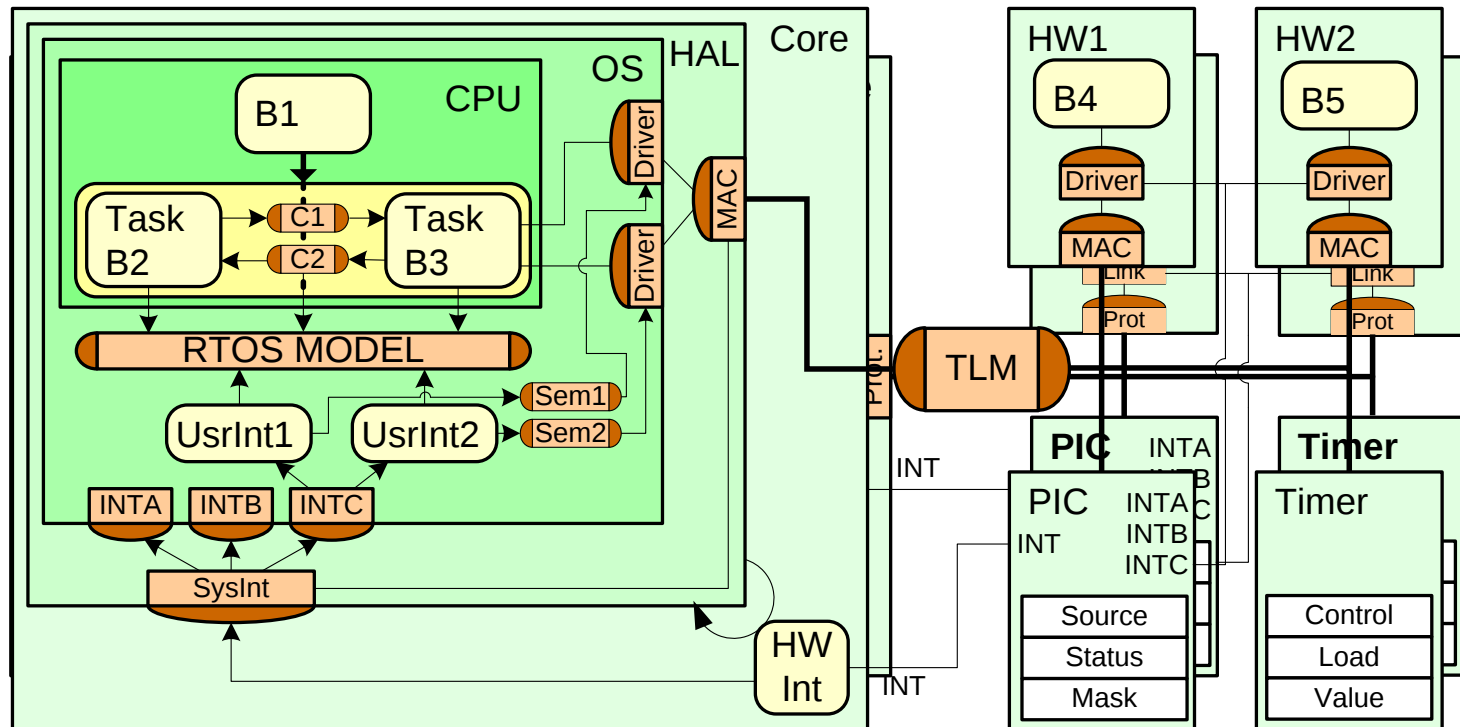
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Execution

- **Validate binary image**
 - Target processor on prototype
 - ISS-based virtual platform



Summary

- **Embedded software generation from system model**
 - Including
 - Communication synthesis (external / internal)
 - Multi-task synthesis
 - Binary image creation
- **Integrated into ESL flow**
 - Seamless solution
- **Complete: from abstract model to implementation!**
 - Completes ESL flow for software
 - Eliminates tedious and error prone manual HdS development
 - Significant productivity gain
 - Enables rapid design space exploration

