MOTION ESTIMATION BY QUADTREE PRUNING AND MERGING

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ABSTRACT

In this paper we propose a rate-distortion optimized motion estimation algorithm that is built upon a quadtree structure. Each node of the quadtree represents a block in the current frame together with its motion vector, and the block size decreases from the root to the leaves. In the first step, the quadtree is pruned according to a ratedistortion criterion in order to obtain blocks of variable sizes. A further rate rebate can be achieved by merging those leaf nodes of the quadtree that can be efficiently represented by the same motion vector. The proposed merging scheme provides a reduction of up to 50% of the rate spent for the motion model with respect to the case that performs pruning only.

1. INTRODUCTION

Accurate motion modeling plays an essential role in any video coding architecture, especially to efficiently represent video contents at low and very low bit-rates. In fact, most of the coding gain achieved by the latest state-of-the-art standard, H.264/AVC [1], comes from improved motion estimation and signalling tools, including blocks of variable sizes, quarter-pel motion accuracy, and a more efficient entropy coding of motion information.

Motion estimation algorithms used in video coding application cannot neglect the amount of information needed to represent the motion model. In fact, at low bit-rates, most of the bit budget is usually allocated to describe the motion, while little remains for encoding prediction residuals. Therefore, the accuracy of the motion representation needs to be tuned according to the target bit-rate: at high bit-rates an accurate (thus costly) motion representation is usually desirable. When the bit-rate decreases the amount of bits allocated to motion is usually reduces, therefore achieving a coarser motion representation. Motion estimation algorithms used in state-of-theart video coding schemes always employ rate-distortion optimization criteria in order to obtain the best motion model representation, measured in terms of the energy of the prediction error, satisfying a given rate constraint [2][3].

In this paper we propose a motion estimation algorithm, specifically designed for video coding applications, that produces a region based motion representation. Here the goal is not to identify the motion models of independently moving regions but to provide a more compact representation of the motion model, without sacrificing its accuracy. For the sake of clarity, before illustrating the details of the proposed algorithm, Figure 1 shows the result of the motion estimation for a frame of the *Table Tennis* sequence. We notice that blocks characterized by the same motion model are grouped together in order to reduce the amount of bits allocated to motion.



Fig. 1. Table Tennis, SIF@30fps. Each region has a single motion vector assigned to it.

The work presented in this paper has been inspired from [4], where a quadtree-based coding scheme is used to efficiently encode images that can be modeled as piece-wise polynomials. Each node in the quadtree corresponds to an image block. A prune-merge scheme is used to segment the input image into regions. The novelty of the work in [4] lies in the fact that merging might involve leaf nodes in the quadtree that are not necessarily children of the same node. One or two polynomial models are used to approximate the image within each region. Both the pruning and the merging phase are rate-distortion optimized in order to achieve the best image approximation for a target rate. The idea of this work is to apply and adapt the prune-merging scheme to the problem of motion estimation.

A similar work has recently appeared in [5], where a merging scheme is applied to the variable block size representation provided by H.264/AVC. The goal of the present work is to study the benefits of merging blocks without reference to a specific coding architecture.

2. QUADTREE MOTION MODEL

The proposed motion estimation algorithm is based on a quadtree data structure, where each node represents a block of the current frame. At the top level of the quadtree¹ the block size is 16×16 . At the third (lowest) level, the block size is 4×4 .

First, motion estimation is performed for blocks of size 4×4 using an exhaustive search approach with a search window of $\pm W$ pixels which is chosen based on the spatial resolution of the sequence (QCIF: W = 16 pixels, CIF/SIF W = 32 pixels, 4CIF W = 48 pixels). In each node we store the distortion (SAD - Sum of Absolute Differences) associated with each of the $(2W+1)^2$ candidate motion

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¹the quadtree is traversed top-to-bottom going from the root to the leaves



Fig. 2. Table Tennis, SIF@30fps. a) Pruning with $\lambda = 0.5$. b) Pruning with $\lambda = 1.5$.

vectors (MVs). Without the need of performing further block comparisons, we infer the distortion associated to each candidate motion vector for blocks of size 8×8 . In fact, for a given candidate motion vector for a block of size 8×8 , the SAD is simply obtained as the sum of the previously stored SADs of the same candidate motion vector applied to its child 4×4 nodes. The same procedure allows to compute the distortions for 16 blocks.

3. PRUNING ALGORITHM

We assume that motion vectors are encoded in a similar way as in H.264/AVC, by encoding the prediction error between the motion vector and its predictor using Exp-Golomb codes. The motion vector predictor is obtained as the median of the motion vectors of its causal neighbors. Therefore, the rate needed to encode a motion vector is:

$$R_{MV} = 2|mv_x - mvp_x| + 1 + 2|mv_y - mvp_y| + 1$$
(1)

where (mv_x, mv_y) are the two components of the motion vector to be encoded and (mvp_x, mvp_y) is the motion vector predictor.

The goal of pruning is to find the optimal partitioning of each 16×16 block into sub-blocks. Using the quadtree representation, this is equivalent to pruning the quadtree in such a way that leaf nodes are not necessarily represented by 4×4 blocks. The pruning algorithm processes the quadtree according to a depth-first order. This means that each 16×16 block is considered one after the other and the pruning decisions for a 16×16 block are taken before considering the next 16×16 block.

The pruning algorithm can be summarized as follows. For each 16×16 block

- 1. Compute the motion vector predictor
- 2. For each candidate MV k, compute the Lagrangian cost as

$$J_{16\times 16}(k) = SAD_{16\times 16}(k) + \lambda(R_{MV}(k) + R_{16\times 16}^{nosplit})$$
(2)

3. Compute the lowest Lagrangian cost

$$J_{16\times 16}^* = \min_{k} J_{16\times 16}(k) \tag{3}$$

- 4. For each 8×8 child block i, i = 0, 1, 2, 3
 - (a) Compute the motion vector predictor

(b) For each candidate MV l, compute the Lagrangian cost as

$$J_{8\times8}^{i}(l) = SAD_{8\times8}^{i}(l) + \lambda(R_{MV}^{i}(l) + H_{8\times8}(split/nosplit))$$
(4)

(c) Compute the lowest Lagrangian cost

$$J_{8\times8}^{i*} = \min J_{8\times8}^{i}(l)$$
 (5)

5. Compute the total Lagrangian cost of the child blocks:

$$J_{8\times8}^{tot*} = \sum_{i=0}^{3} J_{8\times8}^{i*} + \lambda R_{16\times16}^{split}$$
(6)

6. If $J_{16\times 16}^{*} \leq J_{8\times 8}^{tot*}$ do not split the block. Assign to the 16×16 block the motion vector

$$MV = \arg\min J_{16\times 16}(k) \tag{7}$$

Go to next 16×16 block

- 7. Otherwise, for each 8×8 child block i, i = 0, 1, 2, 3
 - (a) Compute the motion vector predictor
 - (b) For each candidate MV *l*, compute the Lagrangian cost as

$$J_{8\times8}^{i}(l) = SAD_{8\times8}(l) + \lambda \left(R_{MV}^{i}(l) + R_{8\times8}^{nosplit}\right)$$
(8)

(c) Compute the lowest Lagrangian cost

$$J_{8\times8}^{i*} = \min J_{8\times8}^{i}(l)$$
(9)

- (d) For each 4×4 child block j, j = 0, 1, 2, 3
 - i. Compute the motion vector predictor
 - ii. For each candidate MV *m*, compute the Lagrangian cost as

$$J_{4\times4}^{j}(m) = SAD_{4\times4}^{j}(m) + \lambda R_{MV}^{j}(m) \quad (10)$$

iii. Compute the lowest Lagrangian cost

$$J_{4\times4}^{j*} = \min_{m} J_{4\times4}(m)$$
(11)

(e) Compute the total Lagrangian cost of the child blocks:

$$J_{4\times4}^{tot*} = \sum_{j=0}^{3} J_{4\times4}^{j*} + \lambda R_{8\times8}^{split}$$
(12)

(f) If $J_{8\times 8}^{i*} \leq J_{4\times 4}^{tot*}$ do not split the block. Assign to the 8×8 block *i* the motion vector

$$MV^{i} = \arg\min_{l} J^{i}_{8\times8}(l) \tag{13}$$

Go to next 8×8 block

(g) Otherwise, split the block into its four child 4×4 blocks. Assign to each of the 4×4 blocks *j* the motion vector

$$MV^{j} = \arg\min_{m} J^{j}_{4\times 4}(m) \tag{14}$$

where λ is the Lagrangian multiplier that can be adjusted based on the target bit-rate. $R_{16\times16}^{nosplit/split}$ is the number of bits used to communicate to the decoder the binary decision of splitting the 16×16 block. We use a context adaptive arithmetic coder to encode this decision, in such a way that the average number of bits is less than 1. The context used is the decision taken for the previously encoded 16×16 block. The term $H_{8\times8}(split/nosplit)$ is the estimated entropy relative to the binary decision of splitting each 8×8 block into 4×4 blocks. In fact, when the decision of splitting the 16×16 block is taken, there is no clue about the splitting decision taken at the next level. We notice that this term does not appear when taking the decision of splitting the 8×8 block as 4×4 is the smallest block size allowed by the proposed algorithm.

Intuitively, the proposed pruning algorithm tends to favor small block sizes when λ is small, as the Lagrangian cost function depends more on the distortion than on the rate needed to encode the motion vectors. By increasing λ , block sizes becomes larger, since the rate allocated to motion vector decreases. Figure 2 shows an example of the block sizes obtained with two different values of λ .

4. MERGING ALGORITHM

We notice by inspecting Figure 2 that several blocks represent the same moving object, therefore it is likely that they can be represented by the same motion vector. The idea underlying the proposed algorithm is to merge together leaf nodes of the quadtree representation, which is the output of the pruning phase. Merging might occur between blocks that are not necessarily children of the same node. A coding gain is obtained if the cost of signalling the merging decision is less than the cost of encoding the motion vector.

The proposed merging algorithm receives in input the quadtree representation produced by the pruning algorithm. As a preprocessing step, for each block a list of N target neighboring blocks is computed. This list contains only neighbors that are not smaller than the current block. For the problem at hand, this implies that each block can have up to four target neighbors it can be merged to $(0 \le N \le 4)$.

The merging algorithm analyzes each leaf node of the quadtree as shown in Figure 3. For each block, the algorithm evaluates if merging is advantageous in rate distortion sense. When a block is merged to one of its neighbors, the motion vector of the neighbor is assigned to it and the merging decision is signaled to the decoder.

The merging algorithm proceeds as follows:

For each block i (with motion vector m):

1. Populate the list of neighboring blocks



Fig. 3. Scanning order followed by the merging algorithm.

2. Compute the lagrangian cost when no merging occurs

$$J_{nomerge} = SAD^{i}(m) + \lambda (R^{i}_{MV}(m) + R_{nomerge})$$
(15)

3. For each neighboring block j, j = 0, ..., N-1 (with motion vector n_j)

(a) Compute the lagrangian cost when the two blocks merged

$$J_{merge}^{j} = SAD^{i}(n_{j}) + \lambda (R_{nomerge} + R_{target})$$
(16)

(b)

- 4. If $J_{nomerge} < \min_j J_{merge}^j$ do not merge the blocks
- 5. If $J_{nomerge} \ge \min_j J_{merge}^j$ merge block *i* with block $k = \arg \min_j J_{merge}^j$

 $SAD^{i}(m)$ is the distortion associated with block *i* when its motion vector *m* is used. $SAD^{i}(n)$ is the distortion associated with block *i* when its motion vector *n* of block *j* is used instead. $R_{nomerge/merge}$ is the cost of encoding the merging decision. We use a context adaptive arithmetic coder to efficiently encode this information. The context used is the merging decision taken for the previously encoded block. R_{target} is the rate spent to encode the index of the target block to merged to within the list identified at step 1 and it is equal to $\log N$ bits.

Figure 1 shows the result of the successive application of pruning followed by merging on the *Table Tennis* sequence. We can easily notice that large areas of constant motion (or zero motion) are clustered together, whereas smaller regions are used to describe more complex moving objects.

5. EXPERIMENTAL RESULTS

We carried out extensive experimental results on several test sequences. Figure 4 shows the rate-distortion curve for the *Foreman* (QCIF), *Table Tennis* (SIF), *Coastguard* and *Bus* sequences (CIF). We compare the performance of the merging algorithm with the case where only pruning is performed. The points of the curve are generated for different values of lambda, ranging from 0 to 2.5: at lower values of lambda, the cost of encoding the motion model increases but the distortion decreases. The rate refers to the total number of bits needed to encode the motion model (splitting decisions, merging decisions and motion vectors). As the merging algorithm is performed after pruning, the cost of encoding the splitting decisions is the same in both cases.



Fig. 4. Comparison between the quality of the estimated motion model after the pruning step and after the proposed pruning&merging algorithm. The quality is evaluated as the average PSNR of the displaced frame difference obtained applying either of the two motion models.

The proposed algorithm consistently outperforms the pruning algorithm at all bit-rates. As expected, the coding gain of the proposed algorithm is higher at low bit-rates (higher values of λ), as the Lagrangian cost function tends to weight more the cost associated with encoding the motion model. A gain of up to +0.6dB is reported for Bus. A smaller coding gain is observed for Foreman (up to +0.4dB), Table Tennis and Coastguard (up to +0.3dB). These figures should be interpreted carefully. In fact, the maximum coding gain is upper bounded by the lowest distortion that can be achieved when $\lambda = 0$, i.e when 4×4 blocks are used everywhere. In this case, no pruning nor merging is actually performed, therefore the two approaches converge to the same upper bound. Even with an arbitrary large bitbudget, a value of 34.4dB is obtained for Foreman, 31.6dB for Table Tennis, 29.9dB for Bus and 32.7dB for Coastguard. Therefore, it is more interesting to interpret these results from the bit-rate perspective. Figure 4 shows that a 40%-50% rate rebate is achieved at low bit-rates, and that the margin remains significant in the whole range of bit-rates relevant for video coding applications.

6. CONCLUSIONS

This paper proposes a two-step pruning/merging motion estimation algorithm based on a quadtree structure. Our experimental results show that there are benefits coming from merging together leaf nodes of the quadtree representation produced in output of the pruning phase. Ongoing research activities are focused on integrating the proposed algorithm into a complete video coding architecture.

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