# N-Detection Under Transparent-Scan 

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#### Abstract

We study the quality of test sequences under a test application scheme called transparent-scan as $n$-detection test sequences. We obtain transparent-scan sequences from combinational test sets. We show that for the same number of clock cycles required to apply a compact single-detection combinational test set, a transparent-scan sequence detects faults more times than the combinational test set. We note that a transparent-scan sequence based on a combinational test set contains unspecified values. We consider several procedures for specifying the unspecified values of the transparent-scan sequence, and study their effects. We also study the extension of a transparent-scan test sequence into an $n$-detection test sequence that detects every target fault at least $n$ times.


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## 1. Introduction

Test sets that contain several different tests to detect each target fault, also referred to as $n$-detection test sets, were shown to be effective in detecting unmodeled faults and defects [1]-[9]. In this work we study the quality of test sequences under a test application scheme called transparent-scan as $n$-detection test sequences.

The transparent-scan approach was proposed in [10], and was shown to achieve higher levels of compaction than other test compaction approaches for scan circuits. It was also found to be necessary in [11] for testing critical paths in a microprocessor that uses partial scan. Under the transparent-scan approach, the distinction that typically exists between scan operations and functional clock cycles is eliminated. This is achieved by using the scan select and scan chain inputs as regular primary inputs of the circuit, and using the scan chain outputs as regular

[^0]primary outputs of the circuit during test generation and test compaction. The result is a synchronous sequential circuit with the extra controllability and observability provided by scan.

A test under the transparent-scan approach is a sequence $T$ of primary input vectors, which assign values to the original primary inputs of the circuit, to the scan select input, and to the scan chain inputs. An output vector includes values corresponding to the original primary outputs of the circuit, and to the scan chain outputs. This unified view of the original primary inputs (outputs) and the scan inputs (outputs) provides complete flexibility in interleaving scan operations and functional clock cycles, and results in high levels of compaction.

The test translation method was proposed in [10] for obtaining an initial test sequence under the transparent-scan approach. Under this method, a combinational test set $C$ (a test set $C$ for the combinational logic of the circuit) is translated into a test sequence $T$ under the transparent-scan approach by translating a scan operation of $C$ into a sequence of primary input vectors where the scan select is high, and translating a functional clock cycle of $C$ into a primary input vector where the scan select is low. In [10], the translated test sequence $T$ is compacted by a static test compaction procedure for synchronous sequential circuits. In this work, we will focus on the quality of $T$ (before compaction) as an $n$-detection test sequence. Several properties of a transparent-scan test sequence $T$, which is translated from a combinational test set $C$, are important.

The number of clock cycles required for applying $T$ (before compaction) is equal to the number of clock cycles required for applying $C$. Nevertheless, $T$ provides more options for detecting faults, as we demonstrate later.

The test sequence $T$ obtained after translation of a combinational test set $C$ is incompletely specified. Even as an incompletely specified test sequence, $T$ may detect faults more times than $C$. By specifying the unspecified values of $T$, the qualify of $T$ as an $n$-detection test sequence can be improved even further.

The paper is organized as follows. In Section 2 we review the test translation method from [10] for obtaining a transparent-scan test sequence. In Section 3 we provide definitions related to the quality of a test set as an $n$ detection test set. We then consider the quality of transparent-scan test sequences as $n$-detection test sequences. We describe three procedures for specifying the unspecified values of a transparent-scan sequence so as to improve its quality. Experimental results are presented in Section 4. In Section 5 we study the possibil-
ity of extending a transparent-scan sequence into an $n$ detection test sequence that detects every target fault at least $n$ times.

## 2. Transparent-scan

We demonstrate the test translation method under the transparent-scan approach by considering ISCAS-89 benchmark circuit $s 27$. The original circuit (without scan) has four primary inputs $a_{0}, a_{1}, a_{2}, a_{3}$, and three state variables. A combinational test set $C$ for the circuit is shown in Table 1. Each combinational test vector $c_{i} \in C$ is divided into a subvector $c_{i I}$ that defines the values of the original primary inputs of the circuit, and a subvector $c_{i S}$ that defines the values of the state variables. We denote the scan select input by $s_{\text {sel }}$, and the scan chain input by $s_{i n p}$. We assume that the circuit has a single scan chain, and that the scan chain is shifted from left to right. An input vector of a transparent-scan sequence consists of the values of the original primary inputs of the circuit, followed by the values of $s_{\text {sel }}$ and $s_{i n p}$ (in this order).

Table 1: Combinational test set $C$ for $s 27$

| $i$ | $c_{i I}$ | $c_{i S}$ |  |
| :---: | :---: | :---: | :---: |
| 0 | 0000 |  | 11 |
| 1 | 1001 |  | 10 |
| 2 | 0100 |  | 10 |
| 3 | 0111 |  | 01 |
| 4 | 1101 |  | 11 |
| 5 | 1010 |  | 00 |
|  | 0 | 1 | 2 |
| sinp | 0 | 1 | 1 |

Figure 1: Single scan chain
To apply the combinational test vector $c_{0}$ of $s 27$, we need the following primary input vectors. We first need to hold $s_{\text {sel }}=1$ for three clock cycles while scanning in the state subvector $c_{0 S}=011$. Due to the assumption that scan chains are shifted to the right, this requires us to set $s_{i n p}=1$ in the first clock cycle, $s_{i n p}=1$ in the second clock cycle, and $s_{i n p}=0$ in the third clock cycle, as demonstrated by Figure 1. We obtain the input vectors $a_{0} a_{1} a_{2} a_{3} s_{\text {sel }} s_{\text {inp }}=\operatorname{xxxx} 11, \mathrm{xxxx} 11$ and xxxx 10 , where an $x$ stands for an unspecified value. Note that we list the values of the four original primary inputs first, followed by $s_{\text {sel }}$ and $s_{\text {inp }}$. We then need to hold $s_{\text {sel }}=0$ while applying $c_{0 I}=0000$ to the original primary inputs of the circuit. This is done by the input vector $a_{0} a_{1} a_{2} a_{3} s_{\text {sel }} s_{\text {inp }}=00000 \mathrm{x}$. Finally, we need to hold $s_{s e l}=1$ for three clock cycles in order to scan out the final state. The resulting transparent-scan sequence is $T=$ xxxx 11 xxxx 11 xxxx 1000000 x xxx 1 x xxxx 1 x xxxx 1 x .

While scanning out the final state of $c_{0}$, we scan-in the state $c_{1 S}=010$ of $c_{1}$. We then hold $s_{s e l}=0$ while applying $c_{1 I}=1001$ to the original primary inputs of the circuit. Continuing in the same way to apply the remaining test vectors in Table 1, we obtain the transparent-scan sequence shown in Table 2 under column $T_{x}$.

It is important to note that the test sequence $T_{x}$ is incompletely specified. Unspecified values are obtained since the combinational test set $C$ does not specify pri-

Table 2: Transparent-scan test sequence for $s 27$

| $u$ | $T_{x}$ | $T_{r n d}$ | CI_ind | $T_{c p i}$ |
| :---: | :---: | :---: | :---: | :---: |
| 0 | xxxx11 | 001111 | 2 | 010011 |
| 1 | xxxx11 | 110111 | 4 | 110111 |
| 2 | xxxx10 | 001110 | 3 | 011110 |
| 3 | 00000 x | 000000 | - | 000000 |
| 4 | xxxx10 | 011110 | 1 | 100110 |
| 5 | xxxx11 | 110011 | 1 | 100111 |
| 6 | xxxx10 | 011110 | 5 | 101010 |
| 7 | $10010 x$ | 100100 | - | 100100 |
| 8 | xxxx10 | 110010 | 0 | 000010 |
| 9 | xxxx11 | 010011 | 3 | 011111 |
| 10 | xxxx11 | 110111 | 0 | 000011 |
| 11 | $01000 x$ | 010001 | - | 010001 |
| 12 | xxxx11 | 000111 | 0 | 000011 |
| 13 | xxxx10 | 110110 | 1 | 100110 |
| 14 | xxxx10 | 100110 | 5 | 101010 |
| 15 | $01110 x$ | 011100 | - | 011101 |
| 16 | xxxx11 | 010011 | 0 | 000011 |
| 17 | xxxx11 | 011111 | 0 | 000011 |
| 18 | xxxx10 | 111010 | 5 | 101010 |
| 19 | $11010 x$ | 110100 | - | 110101 |
| 20 | xxxx10 | 101110 | 5 | 101010 |
| 21 | xxxx10 | 000010 | 3 | 011110 |
| 22 | xxxx10 | 000010 | 3 | 011110 |
| 23 | $10100 x$ | 101001 | - | 101001 |
| 24 | xxxx1x | 101011 | 3 | 011110 |
| 25 | xxxx1x | 001111 | 0 | 000011 |
| 26 | xxxx1x | 010111 | 2 | 010010 |

mary input vectors during scan operations, and it does not specify the value of $s_{i n p}$ during a functional clock cycle. Even with the incompletely specified values, $T$ detects all the faults detected by the combinational test set $C$.

It is also important to note that the length of the test sequence $T$ is equal to the number of clock cycles required for applying the combinational test set $C$ from which $T$ is obtained.

## 3. $N$-detections

In this section we discuss the quality of a transparent-scan sequence as an $n$-detection test set. We first provide a definition of this quality for a combinational test set and for a transparent-scan sequence. We then discuss the use of the unspecified values in $T_{x}$ to improve the quality of the test sequence.

### 3.1. Quality definitions

The quality of an arbitrary combinational test set $C$ as an $n$-detection test set can be defined as follows.

We perform $n$-detection fault simulation of $C$. Under this process, a fault is dropped only after it is detected $n$ times, by $n$ test vectors in $C$. We denote by $n_{\text {det }}(f)$ the number of times a fault $f$ is detected at the end of this fault simulation process. We denote by $F$ the set of faults that are detected at least once by $C$. For a test set $C_{i}$ we use the notation $n_{\text {det }}^{i}(f)$ for the number of times $f$ is detected by $C_{i}$, and $F^{l}$ for the set of faults detected at least once by $C_{i}$.

A simple measure of the quality of $C$ as an $n$ detection test set is obtained by using the total (or the average) number of detections under $C$. The total number of detections is equal to $\sum\left\{n_{d e t}(f): f \in F\right\}$. The average number of detections is equal to $\sum\left\{n_{d e t}(f): f \in F\right\} /|F|$.

However, the total (or average) number of detections does not take into account the following situation. Consider a circuit with three faults $f_{1}, f_{2}$ and $f_{3}$. Let $n=5$. Suppose that a test set $C_{1}$ results in $n_{d e t}^{1}\left(f_{1}\right)=1$, $n_{\text {det }}^{1}\left(f_{2}\right)=2$, and $n_{\text {det }}^{1}\left(f_{3}\right)=5$. Suppose that a test set $C_{2}$ results in $n_{\text {det }}^{2}\left(f_{1}\right)=1, n_{\text {det }}^{2}\left(f_{2}\right)=3$, and $n_{\text {det }}^{2}\left(f_{3}\right)=4$. The total number of detections in both cases is eight. However, defects associated with $f_{2}$, which is detected twice by $C_{1}$, are more likely to be detected by $C_{2}$, where $f_{2}$ is detected three times. The reduction in the number of detections for $f_{3}$ from five to four is less likely to affect the overall quality of $C_{2}$. The resulting metric considers the numbers of detections of the faults as follows.

For $m=1,2, \cdots, n$, we denote by $F_{m}$ the subset of faults in $F$ that are detected $m$ times by $C$, i.e., $F_{m}=\left\{f \in F: n_{\text {det }}(f)=m\right\}$. We denote the size of $F_{m}$ by $\left|F_{m}\right|$. For a test set $C_{i}$ we use the notation $F_{m}^{i}=\left\{f \in F^{i}: n_{d e t}^{i}(f)=m\right\}$. In comparing two test sets $C_{1}$ and $C_{2}$ with subsets of faults $F_{1}^{1}, F_{2}^{1}, \cdots, F_{n}^{1}$ and $F_{1}^{2}, F_{2}^{2}, \cdots, F_{n}^{2}$, respectively, such that $F^{2}=F^{2}$, we consider the first value of $m$ for which $\left|F_{m}^{1}\right| \neq\left|F_{m}^{2}\right|$. We say that $C_{1}$ is of higher quality if $\left|F_{m}^{1}\right| \stackrel{m}{<}\left|F_{m}^{2}\right|$. We say that $C_{2}$ is of higher quality if $\left|F_{m}^{2}\right|<\left|F_{m}^{1}\right|$.

For example, for $C_{1}$ and $C_{2}$ above we have $F_{1}^{1}=$ $\left\{f_{1}\right\}$ and $F_{1}^{2}=\left\{f_{1}\right\} ; F_{2}^{1}=\left\{f_{2}\right\}$ and $F_{2}^{2}=\phi$. The first value of $m$ for which $\left|F_{m}^{\mathrm{T}}\right| \neq\left|F_{m}^{2}\right|$ is $m=2$, and we have $\left|F_{m}^{2}\right|<\left|F_{m}^{1}\right|$. We conclude that $C_{2}$ is of higher quality.

To define the quality of a test sequence $T$ under the transparent-scan approach, we perform $n$-detection fault simulation of $T$. During this process, $n_{d e t}(f)$ is incremented by one for every time unit where $f$ is detected by $T$. A fault $f$ is dropped from further simulation when $n_{\text {det }}(f)$ reaches $n$. This is consistent with the first definition of an $n$-detection test sequence given in [5]. Using the numbers of detections $n_{d e t}(f)$, we can define the total and average number of detections, as well as the metric based on the subsets $F_{m}$ in the same way as for a combinational test set.

### 3.2. The quality of transparent-scan sequences

Several effects may cause a test sequence $T$, which is translated from a combinational test set $C$, to have a higher quality than $C$ as an $n$-detection test set. (1) Under $T$, there is no distinction between scan clock cycles (where $s_{s e l}=1$ ) and functional clock cycles (where $s_{s e l}=0$ ). The test sequence $T$ is simulated during both types of clock cycles, and fault detections are captured on the original primary outputs during both types of clock cycles. For a combinational test set $C$, fault detections are captured on the original primary outputs only during functional clock cycles. As a result, $T$ may have higher numbers of fault detections than $C$. (2) We have the option of specifying the unspecified values of $T$ so as to improve the numbers of detections of faults. It is important to note that the number of clock cycles required for applying $C$ and $T$ (translated from $C$ ) are the same. Thus, we obtain improved $n$-detection quality for the same test application time. Next, we consider each one of the effects listed above.

We first compare $C$ with the transparent-scan sequence $T$ before any unspecified values in $T$ are specified. We consider the example of $s 27$. In row comb of Table 3 we show the results of 6-detection fault simulation for the combinational test set $C$ of $s 27$ from Table 1. We show the average number of detections under column ave. We then show the size of $F_{m}$ under $C$, for $1 \leq m \leq 6$.

Table 3: Quality of tests for $s 27$

| type | ave | $\mathrm{m}=1$ | $\mathrm{~m}=2$ | $\mathrm{~m}=3$ | $\mathrm{~m}=4$ | $\mathrm{~m}=5$ | $\mathrm{~m}=6$ |
| :--- | :--- | ---: | ---: | ---: | ---: | ---: | ---: |
| comb | 2.00 | 15 | 7 | 5 | 5 | 0 | 0 |
| x | 3.25 | 5 | 9 | 5 | 4 | 4 | 5 |
| rnd | 4.28 | 3 | 3 | 5 | 5 | 3 | 13 |
| cpi | 4.38 | 3 | 2 | 5 | 6 | 2 | 14 |

In row $x$ of Table 3 we show the results of 6 detection fault simulation for the transparent-scan test sequence $T_{x}$ obtained by translating $C$. This sequence is shown in Table 2 under column $T_{x}$. It can be seen that the average number of detections is higher for $T_{x}$ than for $C$. More important, $C$ has 15 faults that are detected only once. Under $T_{x}$, only five faults are detected once, and the remaining faults are detected higher numbers of times. This is a result of the fact that $T_{x}$ provides more time units where faults may be detected.

We can specify the unspecified values of $T_{x}$ randomly to obtain the test sequence $T_{\text {rnd }}$ shown in Table 2 under column $T_{\text {rnd }}$. The results of 6-detection fault simulation for $T_{r n d}$ are shown in row rnd of Table 3. It can be seen that specification of the unspecified values in $T_{x}$ improves the quality of the transparent-scan sequence.

Next, we consider a different way of specifying the unspecified values of $T_{x}$. We concentrate on the unspecified values of the original primary inputs of the circuit, which are obtained when $s_{\text {sel }}=1$. We will specify the unspecified values on $s_{i n p}$ when $s_{\text {sel }}=0$ randomly. We note that each test vector $c_{i} \in C$ includes a subvector $c_{i I}$ for the original primary inputs of the circuit. This subvector may be useful in detecting faults at any time unit of $T_{x}$. We define $C_{I}=\left\{c_{i I}: c_{i} \in C\right\}$. We use the primary input subvectors in $C_{I}$ as follows. For every time unit $u$ of $T_{x}$ where $s_{\text {sel }}=1$ and the original primary inputs are unspecified, we randomly select an index $i$ of a subvector $c_{i I} \in C_{I}$. We then assign $c_{i I}$ to the primary inputs at time unit $u$ of $T_{x}$. We also set $C I \_i n d(u)=i$ to indicate that $c_{i I}$ is used at time unit $u$ of the sequence. This will be useful later. We denote the resulting test sequence by $T_{c p i}$ (cpi stands for combinational primary input vectors).

A test sequence $T_{c p i}$ based on the test sequence $T_{x}$ of $s 27$ is shown under column $T_{c p i}$ of Table 2. Before every vector of $T_{c p i}$ we show the index of the combinational test vector whose primary input subvector is used in $T_{\text {cpi }}, C I \_i n d(u)$.

We show the results of 6-detection fault simulation of $T_{c p i}$ for $s 27$ in row cpi of Table 3. It can be seen that the use of primary input vectors based on $C$ improves the quality of the transparent-scan test sequence relative to the use of random values. In the following subsection we describe a procedure for further improving the quality of the test sequence $T_{c p i}$.

### 3.3. Modifying $T_{c p i}$ to improve its quality

We first demonstrate the modification procedure of $T_{c p i}$ by considering the test sequence $T_{c p i}$ of $s 27$ in Table 2.

Based on Table 1, the candidate primary input vectors for $s 27$ are $C_{I}=\left\{c_{i I}: 0 \leq i \leq 5\right\}$. The time units where primary input vectors from $C_{I}$ are embedded in $T_{c p i}$ are $U=\{0,1,2,4,5,6, \cdots, 24,25,26\}$. The index of the primary input vector included in $T_{c p i}$ at time unit $u$, CI_ind ( $u$ ), is shown in Table 2 for $0 \leq u \leq 26$. We consider the time units in $U$ one at a time in a random order. For every time unit $u$, we attempt to replace $c_{C I_{-} i n d(u) I}$ with a different subvector $c_{i I}$. We select the index $i$ randomly. We accept the replacement only if the $n$-detection quality of $T_{c p i}$ is improved.

When $u=2$ is considered, we attempt to replace $c_{C I_{-} i n d(2)}=c_{3 I}=0111$ with $c_{1 I}=1001 \quad(u=2$ and $i=1$ are selected randomly). We find that the number of detections of the fault $f_{7}$ will go down from four to three. We therefore keep $c_{3 I}$ at time unit $u=2$.

When $u=14$ is considered, we attempt to replace $c_{C I \text { ind (14) }}=c_{5 I}=1010$ with $c_{3 I}=0111(u=14$ and $i=3$ are selected randomly). We find that the number of detections of the fault $f_{6}$ changes from six to five; the number of detections of $f_{14}$ changes from five to six; the number of detections of $f_{19}$ changes from three to four; the number of detections of $f_{21}$ changes from six to five; and the number of detections of $f_{31}$ changes from five to six. The effect on the metric based on the sets $F_{m}$ is shown in Table 4. The sizes of the sets $F_{m}$ for $T_{c p i}$ are shown in row $c p i$, and the sizes of the sets $F_{m}$ for the modified test sequence are shown in row mod. Overall, the quality of the modified test sequence is better, and we accept the change. Note that we accept the change even though the number of detections went down for some faults with higher numbers of detections. The reduction in the number of faults with three detections from five to four determines that the modified test sequence is better.

Table 4: Effect of modifying $T_{c p i}$

| type | $\mathrm{m}=1$ | $\mathrm{~m}=2$ | $\mathrm{~m}=3$ | $\mathrm{~m}=4$ | $\mathrm{~m}=5$ | $\mathrm{~m}=6$ |
| :--- | ---: | ---: | ---: | ---: | ---: | ---: |
| cpi | 3 | 2 | 5 | 6 | 2 | 14 |
| mod | 3 | 2 | 4 | 7 | 2 | 14 |

After we consider all the time units in $U$ once we consider them again until two consecutive passes over $U$ do not improve the quality of the test sequence.

To reduce the simulation effort required for computing the quality metric of a modified test sequence, we only consider faults that are detected up to $n_{0}$ times by $T_{c p i}$, for a constant $n_{0}<n$. Thus, we compute the sets $F_{m}$ over a subset of the faults that have $1 \leq n_{\operatorname{det}}(f) \leq n_{0}$ before the modification, and we use these sets to determine whether a modification of $T_{c p i}$ is acceptable. By using $1 \leq m \leq n_{0}$, we ignore the effect of the modification on faults that are detected more than $n_{0}$ times. This is justified since their effect on the overall quality is lower. Even if they are detected fewer times after a modification, an increase in the number of detections of faults detected fewer times is advantageous. We accept a modification only if such an improvement occurs.

After we accept a modification, we compute the numbers of detections of all the faults under the modified test sequence to ensure that the next modification is done using accurate information.

The procedure for modifying $T_{c p i}$ is given next as Procedure 1. We denote by $T_{\text {mod }}$ the test sequence that results from applying Procedure 1 to $T_{c p i}$
Procedure 1: Modifying $T_{c p i}$
(1) Let $C=\left\{c_{0}, c_{1}, \cdots, c_{k-1}\right\}$ be a combinational test set. Let $T_{c p i}=t_{0} t_{1} \cdots t_{L-1}$ be the transparent-scan sequence with original primary input vectors out of $C_{I}=\left\{c_{0 I}, c_{1 I}, \cdots, c_{(k-1) I}\right\}$. If $t_{u}$ contains a subvector $c_{i I} \in C_{I}$, CI_ind $(u)=i$; otherwise, CI_ind $(u)=-$. Set $T_{\text {mod }}=T_{c p i}$.
(2) Perform $n$-detection fault simulation of $T_{\text {mod }}$. Store for every fault $f$ the number of times $n_{d e t}(f)$ the fault is detected by $T_{\text {mod }}$.
(3) Include in $U$ every time unit $u$ such that $C I \_i n d(u) \geq 0$.
(4) Select a time unit $u \in U$ randomly. Remove $u$ from $U$.
(5) Select a subvector $c_{i I} \in C_{I}$ randomly such that CI_ind $(u) \neq i$.
(6) Replace $c_{C I_{-} i n d(u) I}$ in $t_{u}$ with $c_{i I}$.
(7) Let $\hat{F}=\left\{f: n_{\text {det }}(f) \leq n_{0}\right\}$. Simulate $T_{\text {mod }}$ under every fault $f \in \hat{F}$. Considering only the faults in $\hat{F}$, if the quality of $T_{m o d}$ is improved by the change in $t_{u}$ :
(a) Set $C I \_i n d(u)=i$.
(b) Perform $n$-detection fault simulation of $T_{\text {mod }}$. Store for every fault $f$ the number of times $n_{d e t}(f)$ the fault is detected by $T_{\text {mod }}$.
Otherwise, replace $c_{i I}$ in $t_{u}$ with $c_{C I-i n d(u) I}$.
(8) If $U \neq \phi$, go to Step 4.
(9) If $T_{m o d}$ was not modified in two consecutive passes over all the time units in $U$, stop.
(10) Go to Step 3.

## 4. Experimental results

We consider ISCAS-89 and ITC-99 benchmark circuits with $n=10$. For the combinational test set $C$ we use a compact single-detection test set. We consider four transparent-scan sequences for every circuit. (1) The incompletely-specified test sequence $T_{x}$ obtained after translation of $C$. (2) The test sequence $T_{r n d}$ obtained after specifying the unspecified values of $T_{x}$ randomly. (3) The test sequence $T_{c p i}$ obtained after using original primary input vectors in $C$ to specify unspecified values of $T_{x}$. (4) The test sequence $T_{m o d}$ obtained by applying Procedure 1 to $T_{c p i}$. We use $n_{0}=2$ in Procedure 1 .

We always simulate $C$ using a combinational fault simulation procedure. We also perform limited simulation of $T_{x}$, where fault effects are observed on the original primary outputs only when $s_{\text {sel }}=0$, and on the scan output(s) only when $s_{\text {sel }}=1$. This simulation process imitates combinational fault simulation while eliminating the following effects. (1) There are slight variations in fault lists when scan is inserted into a circuit. (2) Combinational fault simulation counts a fault as detected once even if the fault is propagated to several flip-flops and results in a fault effect at several scan clock cycles.

We consider circuits with $1,2,4, \cdots$ scan chains. We use the smallest number of scan chains that results in a transparent-scan sequence of length 5000 or less. We denote the circuit circ that has $S$ scan chains by circ. $S$.

Table 5: Numbers of detections

| circuit | type | ave | $\mathrm{m}=1$ | $\mathrm{m}=2$ | $\mathrm{m}=3$ | $\mathrm{m}=4$ | $\mathrm{m}=5$ | $\mathrm{m}=10$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| s510.1 | comb | 5.45 | 113 | 76 | 51 | 44 | 27 | 178 |
|  | $\lim$ | 5.70 | 104 | 74 | 46 | 41 | 29 | 199 |
|  | x | 7.72 | 47 | 46 | 30 | 16 | 18 | 362 |
| s526.1 | comb | 5.45 | 92 | 112 | 48 | 30 | 28 | 189 |
|  | $\lim$ | 5.68 | 89 | 100 | 45 | 32 | 28 | 204 |
|  | x | 9.24 | 0 | 4 | 2 | 5 | 2 | 431 |
| s641.1 | comb | 4.97 | 99 | 77 | 42 | 25 | 31 | 109 |
|  | $\lim$ | 6.04 | 67 | 70 | 31 | 22 | 32 | 180 |
|  | x | 6.13 | 67 | 70 | 31 | 22 | 30 | 193 |
|  | rnd | 8.52 | 20 | 16 | 22 | 7 | 13 | 341 |
|  | cpi | 9.36 | 3 | 6 | 10 | 8 | 9 | 406 |
|  | mod | 9.36 | 1 | 3 | 14 | 10 | 8 | 405 |
| s820.1 | comb | 3.54 | 336 | 180 | 75 | 27 | 28 | 136 |
|  | $\lim$ | 3.71 | 324 | 166 | 72 | 29 | 43 | 144 |
|  | x | 5.07 | 162 | 142 | 82 | 84 | 51 | 235 |
|  | rnd | 6.00 | 93 | 120 | 76 | 80 | 65 | 314 |
|  | cpi | 6.02 | 85 | 125 | 80 | 80 | 54 | 309 |
|  | mod | 6.11 | 67 | 116 | 100 | 88 | 47 | 312 |
| s953.1 | comb | 5.58 | 205 | 159 | 102 | 75 | 37 | 387 |
|  | $\lim$ | 6.38 | 135 | 142 | 81 | 81 | 37 | 484 |
|  | x | 9.52 | 5 | 16 | 9 | 16 | 5 | 969 |
| s1196.1 | comb | 6.11 | 239 | 139 | 72 | 58 | 62 | 522 |
|  | x | 6.65 | 141 | 172 | 74 | 58 | 51 | 607 |
|  | rnd | 8.60 | 25 | 60 | 46 | 37 | 39 | 919 |
|  | cpi | 9.59 | 8 | 8 | 13 | 13 | 12 | 1131 |
|  | mod | 9.61 | 0 | 8 | 19 | 14 | 12 | 1134 |
| s1423.1 | comb | 5.36 | 291 | 164 | 165 | 102 | 107 | 388 |
|  | x | 8.51 | 8 | 61 | 126 | 34 | 19 | 1057 |
|  | rnd | 8.62 | 5 | 47 | 127 | 43 | 11 | 1096 |
|  | cpi | 8.66 | 5 | 48 | 104 | 58 | 19 | 1104 |
|  | mod | 8.66 | 5 | 46 | 106 | 58 | 19 | 1104 |
| s5378.4 | comb | 7.66 | 359 | 356 | 203 | 197 | 167 | 2837 |
|  | X | 9.17 | 72 | 129 | 104 | 76 | 60 | 3899 |
|  | rnd | 9.50 | 48 | 74 | 72 | 51 | 26 | 4182 |
|  | cpi | 9.52 | 49 | 74 | 73 | 43 | 22 | 4199 |
| $\overline{\text { s } 9234.8 ~}$ | comb | 6.11 | 1058 | 846 | 501 | 351 | 296 | 2753 |
|  | X | 7.22 | 545 | 672 | 348 | 361 | 250 | 3592 |
| $\overline{\text { s15850.16 }}$ | comb | 6.89 | 1596 | 1171 | 648 | 445 | 342 | $\overline{6160}$ |
|  | x | 8.41 | 621 | 570 | 360 | 309 | 246 | 8320 |
|  | rnd | 8.47 | 611 | 543 | 360 | 307 | 251 | 8568 |
|  | cpi | 8.48 | 607 | 547 | 362 | 293 | 251 | 8565 |
| s35932.8 | comb | 3.27 | 9025 | 8227 | 5070 | 3531 | 2536 | 145 |
|  | x | 4.66 | 8522 | 7121 | 3791 | 2066 | 1046 | 9120 |
| $\overline{\mathrm{b} 04.1}$ | comb | 7.67 | 63 | 74 | 77 | 87 | 78 | 770 |
|  | x | 9.71 | 0 | 11 | 9 | 5 | 11 | 1253 |
| b09.1 | comb | 5.88 | 40 | 48 | 51 | 37 | 37 | 122 |
|  | x | 8.10 | 7 | 31 | 17 | 20 | 11 | 262 |
| b10.1 | comb | 6.04 | 83 | 60 | 45 | 33 | 27 | 207 |
|  | x | 8.36 | 19 | 12 | 17 | 35 | 19 | 348 |
| b11.1 | comb | 6.71 | 119 | 106 | 78 | 75 | 59 | 500 |
|  | x | 9.34 | 0 | 12 | 29 | 21 | 20 | 934 |
| b14.16 | comb | 7.93 | 264 | 359 | 465 | 541 | 413 | 4982 |
|  | x | 9.13 | 129 | 161 | 165 | 226 | 121 | 6778 |

The results are given in Table 5. In row comb we show the results of 10 -detection fault simulation for the combinational test set $C$. In row lim we show for some of the circuits the results of limited 10-detection fault simulation for the transparent-scan sequence $T_{x}$ as described above. In rows $x$, rnd, cpi and mod we show the results of full 10-detection fault simulation for the transparentscan sequences $T_{x}, T_{r n d}, T_{c p i}$ and $T_{m o d}$, respectively. In every case we show the average number of detections under column ave, and the size of $F_{m}$ under column $m=\hat{m}$, for $1 \leq \hat{m} \leq 5$ and $\hat{m}=10$. We omit the results for
$T_{r n d}, T_{c p i}$ and $T_{m o d}$ if the quality of the sequence does not change relative to $T_{x}$. The following points can be seen from Table 5.

The quality of $T_{x}$ as an $n$-detection test set is always better than the quality of $C$. This is due to the fact that more input vectors are available for detecting faults under $T_{x}$, even though the number of clock cycles for applying $T_{x}$ is the same as that of $C$, and $T_{x}$ is incompletely specified.

Specifying $T_{x}$ improves the quality of the test sequence in many cases. In the cases where the quality is not improved, fault detection has only a weak dependence on the original primary input vectors. This typically happens when the circuit has a small number of primary inputs. In the cases where the quality is improved, $T_{c p i}$ is in most cases better than $T_{r n d}$ as an $n$-detection test sequence. Procedure 1 is able to improve the quality of the test sequence even further in these cases.

## 5. Generating $n$-detection sequences

The transparent-scan sequences $T_{x}, T_{r n d}, T_{c p i}$ and $T_{\text {mod }}$ detect faults more times than $C$. However, they do not reach $n$ detections for each target fault, even for small values of $n>1$. In this section we complete a transparent-scan sequence into an $n$-detection sequence. We use $T_{r n d}$ as the initial test sequence. Other test sequences can be used in a similar way.

We perform $n$-detection fault simulation to find the number of times $n_{d e t}(f)$ every fault $f$ is detected by $T_{r n d}$. Since $C$ is a single-detection test set, $T_{r n d}$ is also guaranteed to be a single-detection test sequence, i.e., $n_{\text {det }}(f) \geq 1$ for every target fault $f$. For $d=2,3, \cdots, n$, we extend $T_{r n d}$ into a $d$-detection test sequence by targeting every fault $f$ such that $n_{d e t}(f)<d$. When a fault $f$ is targeted, we attempt to generate a test subsequence $T(f)$ that detects $f$ starting from the final states reached under $T_{r n d}$. We search for $T(f)$ by performing sequential test generation. If the fault is activated on a flip-flop but not detected, we use an activation sequence followed by a scan-out operation of the appropriate length to detect it. If a test subsequence $T(f)$ that detects $f$ is found, then $T_{r n d} T(f)$ detects $f$ one additional time compared to $T_{r n d}$. We replace $T_{r n d}$ by $T_{r n d} T(f)$ and perform $n$-detection fault simulation in order to update the numbers of detections of all the faults. We repeat this process until all the faults are detected $d$ times. We then consider the next value of $d$.

To reduce the $n$-detection fault simulation effort during test generation, we only simulate a fault $f$ with $n_{d e t}(f)<d$ during $d$-detection test generation, since only such a fault may have to be targeted. We perform a complete $n$-detection fault simulation pass only after the $d$ detection test sequence is obtained. In addition, we drop a fault from consideration once it is detected $n$ times.

Table 6: $N$-detection test generation

| circuit |  | $\mathrm{d}=1$ | $\mathrm{~d}=2$ | $\mathrm{~d}=3$ | $\mathrm{~d}=4$ | $\mathrm{~d}=5$ | $\mathrm{~d}=10$ |
| :--- | :--- | ---: | ---: | ---: | ---: | ---: | ---: |
| s510.1 | len | 384 | 537 | 639 | 685 | 722 | 869 |
|  | rtio | 1.00 | 1.40 | 1.66 | 1.78 | 1.88 | 2.26 |
|  | ndet<d | 0 | 0 | 0 | 0 | 0 | 0 |
| s526.1 | len | 1121 | 1121 | 1140 | 1193 | 1241 | 2948 |
|  | rtio | 1.00 | 1.00 | 1.02 | 1.06 | 1.11 | 2.63 |
|  | ndet<d | 0 | 0 | 0 | 0 | 0 | 10 |
| s641.1 | len | 459 | 768 | 1255 | 1764 | 2168 | 4288 |
|  | rtio | 1.00 | 1.67 | 2.73 | 3.84 | 4.72 | 9.34 |
|  | ndet<d | 0 | 0 | 0 | 1 | 1 | 1 |
| s820.1 | len | 569 | 850 | 1163 | 1363 | 1568 | 2509 |
|  | rtio | 1.00 | 1.49 | 2.04 | 2.40 | 2.76 | 4.41 |
|  | ndet<d | 0 | 7 | 1 | 3 | 5 | 7 |
| s953.1 | len | 2309 | 2327 | 2394 | 2445 | 2518 | 3193 |
|  | rtio | 1.00 | 1.01 | 1.04 | 1.06 | 1.09 | 1.38 |
|  | ndet<d | 0 | 0 | 0 | 0 | 0 | 0 |
| s1196.1 | len | 2640 | 2770 | 3288 | 3681 | 4063 | 6423 |
|  | rtio | 1.00 | 1.05 | 1.25 | 1.39 | 1.54 | 2.43 |
|  | ndet<d | 0 | 4 | 7 | 5 | 5 | 9 |
| s1423.1 | len | 2024 | 2169 | 2742 | 390 | 4767 | 8462 |
|  | rtio | 1.00 | 1.07 | 1.35 | 1.93 | 2.36 | 4.18 |
|  | ndet<d | 0 | 0 | 4 | 8 | 13 | 18 |
| b04.1 | len | 4220 | 4220 | 4468 | 4545 | 4678 | 5465 |
|  | rtio | 1.00 | 1.00 | 1.06 | 1.08 | 1.11 | 1.30 |
|  | ndet<d | 0 | 0 | 0 | 0 | 0 | 0 |
| b09.1 | len | 956 | 1007 | 1476 | 1985 | 2633 | 5985 |
|  | rtio | 1.00 | 1.05 | 1.54 | 2.08 | 2.75 | 6.26 |
|  | ndet<d | 0 | 0 | 0 | 0 | 0 | 2 |
| b10.1 | len | 1043 | 1400 | 1626 | 1792 | 1962 | 2801 |
|  | rtio | 1.00 | 1.34 | 1.56 | 1.72 | 1.88 | 2.69 |
|  | ndet<d | 0 | 0 | 0 | 0 | 0 | 0 |
| b11.1 | len | 2417 | 2417 | 2633 | 2902 | 3055 | 4044 |
|  | rtio | 1.00 | 1.00 | 1.09 | 1.20 | 1.26 | 1.67 |
|  | ndet<d | 0 | 0 | 0 | 0 | 0 | 1 |
|  |  |  |  |  |  |  |  |

Results of this process using $n=10$ are shown in Table 6. For $1 \leq d \leq 5$ and for $d=10$ we show the results of $d$-detection test generation as follows. In row len we show the length $L_{d}$ of $T_{\text {rnd }}$ after $d$-detection test generation. In row rtio we show the test length ratio $L_{d} / L_{1}$, where $L_{1}$ is the length of the initial test sequence $T_{r n d}$. In row ndet $<d$ we show the number of faults with $n_{\text {det }}(f)<d$. Ideally, this number should be zero if every fault is detected at least $d$ times after $d$-detection test generation. The following points can be seen from Table 6.

A small number of faults have $n_{d e t}(f)<d$ at the end of the $d$-detection test generation process. These faults can be handled by a more complete test generation procedure. Due to the incompleteness of the procedure we use, and due to accidental detections, the number of faults with $n_{d e t}(f)<d$ may not change monotonically with $d$.

For many circuits, the growth of the test length with $d$ is slow relative to the approximately linear growth observed for compact combinational test sets. This is due to the fact that additional clock cycles are available for detecting faults compared to a combinational test set. In addition, the sequential test generation procedure is able to interleave scan clock cycles with functional clock cycles in order to derive short test subsequences.

## 6. Concluding remarks

We studied the quality of test sequences under transparent-scan as $n$-detection test sequences. The test sequences we considered were obtained from combina-
tional test sets by a process called test translation. A transparent-scan sequence $T$ requires the same number of clock cycles as the combinational test set $C$ from which it is translated, and it is incompletely specified. We showed that under these conditions, a transparent-scan sequence detects faults more times than a combinational test set. This is due to the fact that under transparent-scan, fault detections are captured on all the outputs during scan clock cycles as well as during functional clock cycles.

We considered three procedures for specifying the unspecified values of the transparent-scan sequence. In many cases, the procedure that used randomly selected primary input subvectors out of $C$ increased the numbers of times faults are detected relative to random specification. Other methods of filling unspecified values can be used as well, and are guaranteed to have a quality at least equal to that of the incompletely specified sequence. We also considered the extension of a transparent-scan test sequence into an $n$-detection test sequence that detects every target fault at least $n$ times.

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