Parameterized Block-Based Non-Gaussian Statistical Gate Timing Analysis

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Abstract
As technology scales down, timing verification of digital integrated circuits becomes an increasingly challenging task due to the gate and wire variability. Therefore, statistical timing analysis (denoted by STA) is becoming unavoidable. This paper introduces a new framework for performing statistical gate timing analysis for non-Gaussian sources of variation in block-based STA. First, an approach is described to approximate a variational RC-π load by using a canonical first-order model. Next, an accurate variation-aware gate timing analysis based on statistical input transition, statistical gate timing library, and statistical RC-π load is presented. Finally, to achieve the aforementioned objective, a statistical effective capacitance calculation method is presented. Experimental results show an average error of 6% for gate delay and output transition time with respect to the Monte Carlo simulation with 10^4 samples while the runtime is nearly two orders of magnitude shorter.

1. Introduction
Process technology and environment-induced variability of gates and wires in VLSI circuits makes timing analysis of such circuits a challenging task [1]. More precisely, advanced analysis tools must be developed that are capable of verifying changes in the circuit timing which stem from various sources of variations [2]. In block-based statistical timing analysis (STA), every timing quantity of interest (e.g., delay and slew, arrival time and required arrival time) is represented as a function of global sources of variation (denoted by $\xi$) and independent random sources of variation (denoted by $\eta$) in the canonical first-order (denoted by CFO) form. The advantages of such a formulation are that a) it can capture all correlations and b) it can produce delay sensitivities due to changes in various environmental and process-related parameters [2]. Sources of variations have often been assumed to be Gaussian, which in turn simplifies the block-based STA. However, it has been recently reported that certain process parameters exhibit non-Gaussian probability distributions [3].

Block-based STA breaks its analysis into two parts: 1) variational interconnect timing analysis [4][5] and 2) variational gate timing analysis. Unfortunately, block-based STA is lacking in variation-aware gate timing analysis. The authors in [7] propose a modeling technique for gate delay variability considering multiple input switching. In [8], a model for calculating statistical gate delay variation caused by intra-chip and inter-chip variability is presented. Recent works do not provide an accurate means of analyzing the gate propagation delay and output slew as a function of variational input transition, variation-aware gate timing library, and variational gate load. In this paper a new framework is proposed for determining variational gate timing behavior. This is achieved by performing the following steps:

1. Given the variational resistive-capacitive load (where all resistances and capacitances are represented in the CFO form), an efficient and accurate algorithm is presented to calculate variation-aware RC-π load. To perform the analysis, we calculate the variation-aware admittance moments (cf. section 3), and as a result, the resistance and capacitances in the RC-π load can be written in the CFO form.

2. Based on the statistical RC-π load obtained in step 1, we calculate the variation-aware effective capacitance in the CFO form. In order to achieve the aforementioned goal, a new approach for effective capacitance calculation in static timing analysis (STA) is proposed (cf. section 4.1). This effective capacitance calculation method is used to calculate the variational effective capacitance considering non-Gaussian process and environmental sources of variation in the CFO form (cf. section 4.2).

3. Given the variational input transition time, statistical gate timing library, and variational effective capacitance ($C_{\eta}$) load in the CFO form, we calculate variational gate delay and output transition time in the CFO form (cf. sections 2.2.1).

We point out that although, in the remainder of this paper, we will mainly focus on the CFO random variables to represent process and environmental sources of variation as well as the performance quantities of interest; the work itself is not limited to the first-order approximation of these quantities. In fact, it is straightforward to extend the approach to more complex (e.g., second-order) forms regardless of considering Gaussian or non-Gaussian parameter variations.

The remainder of this paper is as follows. In section 2, we review the background of block-based STA. We also show how to convert a quantity, which itself is a function of global and independent sources of variation, into a canonical first-order (CFO) form. The variation-aware RC-π calculation is presented in section 3. Section 4 explains the statistical gate timing analysis for the variational input rise time, variation-aware gate timing library, and variational RC-π load. In this section a new statistical effective capacitance calculation will be proposed and used for gate timing analysis, which is the key contribution of this paper. Section 5 presents experimental results. Finally, conclusions are discussed in section 6. We use the notation shown in Table 1 throughout the paper.

<table>
<thead>
<tr>
<th>Notation</th>
<th>Description</th>
</tr>
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<tbody>
<tr>
<td>$A$</td>
<td>A deterministic variable (does not take into account any statistical variation)</td>
</tr>
<tr>
<td>$\nu^A$</td>
<td>An arbitrary (non-CFO) random variable, which is a function of $m$ global and $p$ independent random sources of variation</td>
</tr>
<tr>
<td>$\nu^{\nu^A}$</td>
<td>A CFO random variable, which is a function of $m$ global and $p$ independent random sources of variation $i.e., A = A_0 + \sum_{j=1}^{p} A_{\nu^j} \Delta x_j + \sum_{j=k+1}^{m} A_{\nu^{\nu^j}} \Delta S_j$</td>
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2. Background
As mentioned before, the sources of variation may exhibit non-Gaussian distributions. Therefore, in general, in addition to calculating the mean and variance of the electrical and timing parameters, we need to calculate the skewness of their distributions, i.e. using the first three moments of the parameters variations.
Definition: The degree of asymmetry of a distribution is called skewness (denoted by $\kappa$). A distribution, or data set, is symmetric if it looks the same to the left and right of the center point. The skewness for a normal distribution is zero. Negative values for the skewness indicate data that are skewed left whereas positive values for the skewness indicate data that are skewed right. By skewed left (right), we mean that the left (right) tail is heavier than the right (left) tail.

The skewness of a distribution is defined by $\kappa = \frac{\mu_3}{\sigma^3}$, where $\mu_3$ is the 3rd central moment and $\sigma^2$ is the variance (second central moment).

Lemma 1: Suppose $\Delta S_1,...,\Delta S_n$ are $n$ independent random variables with distribution $\Delta S_i \sim \text{Dist}(\mu=0, \sigma=1, \kappa)$. Then:

$$\sum_{i=1}^{n} a_i \Delta S_i = \sum_{i=1}^{n} a_i \Delta X_i + a_{m+1} \Delta S_i$$

where $a_0$ is the nominal value; $\Delta X_i$‘s represent the variation of $m$ global sources of variation, $X_i$ from their nominal values, $a_i$’s are the sensitivities to each of the global sources of variation, $\Delta S_i$ is the variation of independent random variable $S_i$ and $a_{m+1}$ is the sensitivity of the timing quantity to $S_i$. By scaling the sensitivity coefficients, we can assume that $\Delta X_i$ and $\Delta S_i$ have distributions with $\mu=0$ and $\sigma^2=1$ and skewness $\kappa$ denoted by $\text{Dist}(\mu=0, \sigma=1, \kappa)$.

Variation in the physical dimensions of the wire causes change in its resistance and capacitance, thereby, making the gate delay and slew as well as wire delay and slew to vary accordingly [9]. Therefore, we need to capture the effect of geometric variations on the electrical parameters. For instance, resistance and capacitance in the CFO form are calculated as follows:

$$r = r_0 + \sum_{i=1}^{m} c_i \Delta X_i + r_{m+1} \Delta S_i$$

$$c = c_0 + \sum_{i=1}^{m} c_i \Delta X_i + c_{m+1} \Delta S_i$$

where $r_0$ and $c_0$ represent nominal resistance and capacitance values, computed when the wire dimensions are at their nominal or typical values. The other parameters are as explained above.

Observation: Invariant Functional Form Property: This property states that: $y = f(x) \Rightarrow y = f'(\tilde{x})$, which follows from the fact that form of function $f$ is independent of its input type (deterministic or variational).

2.2 Converting a variational function into CFO form

It is important to represent timing and electrical quantities in the CFO form. This in turn enables one to propagate first order sensitivities to different sources of variation through timing graph [2][9]. In addition, it makes statistical computations efficient and practical and provides timing diagnostics at a very small cost in run time. The remaining question is how to convert a quantity of interest (which itself is a function of different CFO variables) into the CFO form.

The following subsection presents a method to answer the above question. We use an example to show the procedure. The problem we address is how to convert the gate output transition time into the CFO form. However, this method can be easily applied to any other quantity of interest.

2.2.1 Gate timing analysis for lumped capacitive load

Problem Statement 1: Given is a variational CMOS driver where its input rise time, $t_{in}$, is in the CFO form and drives an output capacitive load, also, in the CFO form. Note that the distribution characteristics of all global and independent sources of variation ($\mu=0$, $\sigma^2=1$, $\kappa$) are given. The objective is to calculate the output transition time, $t_r$, in the CFO form:

$$t_r = t_{r,0} + \sum_{i=1}^{p} t_{r,i} \Delta X_i + t_{r,m+1} \Delta S_i$$

i.e., calculate the nominal value ($t_{r,0}$) and the sensitivity coefficients ($t_{r,i}$ and $t_{r,m+1}$) as well as the skewness of distribution of $\Delta S_i$. The gate output transition time is a function of the input transition time, the logic gate characteristics (e.g., the $W/L$ ratio, threshold voltage of transistors, $V_{dd}$, and temperature), and the output load. In commercial ASIC cell libraries, it is possible to characterize various output transition times (e.g. 10%, 50%, and 90%) as a function of above variables; i.e.:

$$t_r = TF(t_{in}, c_j, z) \text{ where } z = \begin{cases} \frac{w}{l}, & v_i, v_{dd}, \text{Temp} \end{cases}$$

where $t_r$ is the output transition time and $TF$ is the corresponding output transition time function. $z$ captures the gate characteristics and environmental factors, $t_{in}$ is the input transition time, and $c_j$ is the output capacitive load. Based on the Invariant Functional Form Property, the form of function $TF$ is independent of its input type (deterministic or variational.) Hence, we extend the above equation to the variational case. In block-based statistical timing analysis tool, a first-order variational model is employed for all timing quantities such as the gate and wire delays, arrival times, required arrival times, slacks and slacks, i.e., all timing quantities are expressed in the CFO form as follows:

$$\Delta S_i \sim \text{Dist}(\mu=0, \sigma=1, \kappa)$$

By using Lemma 1:

$$\Delta S_i \sim \text{Dist} \left( \mu=0, \sigma^2=1, \kappa=\frac{\mu_3}{\sigma^3} \right)$$

where $\mu_3$ is the 3rd central moment and $\sigma^2$ is the variance (second central moment).

By scaling the sensitivity coefficients, we can assume that $\Delta X_i$ and $\Delta S_i$ have distributions with $\mu=0$ and $\sigma^2=1$ and skewness $\kappa$. It is omitted for brevity.

Proof: It is omitted for brevity.
In Lemma 2, we present how to calculate addition, multiplication, and division of two CFO forms in a new CFO form.

**Lemma 2:** Suppose, \( a \) and \( b \) are two given CFO random variables as:

\[
\begin{align*}
\bar{a} &= a_0 + \sum_{i=1}^{\infty} a_i \Delta X_i + a_{\sigma_1} \Delta S_{\sigma_1} \\
\bar{b} &= b_0 + \sum_{j=1}^{\infty} b_j \Delta X_j + b_{\sigma_1} \Delta S_{\sigma_1}
\end{align*}
\]

Therefore, for addition, subtraction, multiplication, and division of \( a \) and \( b \), we have:

a) Addition and subtraction:

\[
\bar{c} = \bar{a} \pm \bar{b} = (a_0 \pm b_0) + \sum_{i=1}^{\infty} (a_i \pm b_i) \Delta X_i + \sqrt{(a_{\sigma_1})^2 + (b_{\sigma_1})^2} \Delta S_{\sigma_1}
\]

b) Multiplication:

\[
\bar{c} \equiv \bar{a} \times \bar{b} = a_0 \cdot b_0 + \sum_{i=1}^{\infty} \sum_{j=1}^{\infty} (a_i b_j + a_j b_i) \Delta X_i \Delta X_j + \sqrt{(a_{\sigma_1})^2 + (b_{\sigma_1})^2} \Delta S_{\sigma_1}
\]

c) Division:

\[
\bar{c} \equiv \frac{\bar{a}}{\bar{b}} = \frac{a_0}{b_0} + \sum_{i=1}^{\infty} \frac{a_i b_j - a_j b_i}{b_0^2} \Delta X_i \Delta X_j + \sqrt{(a_{\sigma_1})^2 + (b_{\sigma_1})^2} \Delta S_{\sigma_1}
\]

**Proof:** It is omitted for brevity.

### 3. RC-\( \pi \) Load Calculation in the CFO Form

In VDSM technologies, one cannot neglect the effect of interconnect resistance of the load on the gate delay and output transition time. In STA, an adequate approximation of an \( n^\text{th} \) order load seen by the gate (i.e., a load with \( n \) distributed capacitances to ground) is obtained by replacing the load by a second order RC-\( \pi \) model [10]. Equating the first, second, and third moments of the real load with the first, second, and third moments of the RC-\( \pi \) load, one can compute \( c_n, r_k, \) and \( c_j \) as [11]:

\[
\begin{align*}
\bar{c}_n &= \frac{Y_{1,n}}{Y_{3,\sigma_1}} Y_{3,n} \quad r_k = \frac{Y_{1,n}^2}{Y_{3,\sigma_1}^2} \quad c_j = \frac{Y_{2,n}^2}{Y_{3,\sigma_1}}
\end{align*}
\]

where \( Y_{1,n} \) is the \( n^\text{th} \) moment of the admittance of the real load. In \( \sigma \)TA, it is required to consider the effect of variability of the load on the gate timing analysis, as detailed below.

**Problem Statement II:** Given is an RC network representation of the load of a logic gate in a design as exemplified in Figure 1(a), where each \( r \) and \( c \) is in the CFO form. Note that the distribution characteristics of all global and independent sources of variation (\( \mu=0, \sigma^2=1, \kappa \)) are given. The objective is to calculate an equivalent variational RC-\( \pi \) load (i.e., \( c_n, r_k, \) and \( c_j \) of Figure 1(b) are in the CFO form), while its admittance matches the admittance of the real load in the frequency range of interest.

\( c_n, r_k, \) and \( c_j \) are functions of the admittance moments as seen from Eqn. (3). Hence, by calculating the variational admittance moments, we can calculate the CFO parameters of RC-\( \pi \) load (using the technique explained in section 2.2.) This can be done by differentiating the expressions in Eqn. (3) with respect to the sources of variation (cf. section 2.2.). However, as it will be shown next, a recursive operation is utilized to calculate the variational admittance moments and since in each recursion step, we have a complex (non-CFO) random variable which will feeds in the next step and this may increase the complexity of the calculations;

We represent the admittance moments in the CFO form throughout the recursion. This helps us by controlling the complexity of presenting the moments as the recursive function proceeds. Following shows how to calculate the input admittance moments of the real load in the CFO form. Consider the RCY segment shown in Figure 2. Assume that the admittances at nodes \( i \) and \( j \) are represented by infinite series using the admittance moments:

\[
Y_i(s) = s Y_{1,i} + s^2 Y_{2,i} + \ldots + s^4 Y_{k,i} + \ldots
\]

\[
Y_j(s) = s Y_{1,j} + s^2 Y_{2,j} + \ldots + s^4 Y_{k,j} + \ldots
\]

where \( Y_{k,i} \) denotes the coefficient of \( s^k \) in the admittance of node \( i \). Thus, in STA, the admittance at node \( i \) is recursively computed in terms of the admittance at node \( j \) [11]:

\[
Y_{1,i} = Y_{1,j} + c_i
\]

\[
Y_{k,i} = Y_{k,j} - r_k \sum_{m=1}^{k-1} c_{k-m} Y_{m,j} - c_j Y_{k-1,j}
\]

Using the Invariant Functional Form Property, we extend the above equation to the variational case. Assume the admittance moments of node \( j \) are given (written in the CFO form). Thus, by differentiating \( Y_{i,j} \) with respect to the sources of variations, the \( Y_{k,i} \) moments can be also represented in the CFO form (cf. section 2.2.)

\[
\begin{align*}
\bar{c}_n &= \frac{Y_{1,n}}{Y_{3,\sigma_1}} Y_{3,n} \quad \bar{c}_j = \frac{Y_{2,n}^2}{Y_{3,\sigma_1}}
\end{align*}
\]

Figure 2: an RCY segment model for recursive admittance moment calculation.

By using the above recursive operations, we easily compute the moments of \( Y_{m}=Y_{1,j} \) in the CFO form, and hence we calculate the values of \( c_n, r_k, \) and \( c_j \) in the CFO form using Eqn. (3).

### 4. Gate Timing Analysis for the RC-\( \pi \) Load in Block-Based \( \sigma \)TA

**Problem statement III:** Given is a variational CMOS driver, whose input rise time, \( t_{in} \), is in the CFO form and drives a variational RC-\( \pi \) load. The resistance and capacitances of this load are also in the CFO forms. The distribution characteristics of all global and independent sources of variation (\( \mu=0, \sigma^2=1, \kappa \)) are given.

The objective is to calculate the output transition time, \( t_{out} \), in the CFO form:

\[
t_{out} = t_{in} + \sum_{m=1}^{\infty} t_{m} \Delta X_{m} + t_{m+1} \Delta S_{\sigma_1}
\]

i.e., calculate the nominal value (\( t_{in} \)) and the sensitivity coefficients (\( t_{j} \) and \( t_{m+1} \)) as well as the skewness of distribution of \( \Delta S_{\sigma_1} \).

Section 2.2.1 solves the same problem where the gate drives a variational purely-capacitive load in the CFO form. (cf. Eqn. (1)) Therefore, if we substitute the RC-\( \pi \) load with its equivalent variational effective capacitance, \( c_{\pi_{\sigma_1}} \) in the CFO form, then the solution to problem statement I is an acceptable solution to problem statement III. Based on this reasoning, the following subsections propose a solution for calculating the effective capacitance in the CFO form. Section 4.1 presents a new effective capacitance.
calculation in static timing analysis. This approach is used in section 4.2 where statistical effective capacitance is calculated.

4.1 A new approach for effective capacitance calculation in static timing analysis

By definition, the effective capacitance is a pure capacitance that replaces an RC-π load and has the property that it gives the most accurate result from a timing model that is characterized with lumped capacitance. Typically, the effective capacitance stores the same amount of charge as the RC-π load until a certain point of the output voltage transition [11][12][13] (e.g., the 50% point of the output transition.) Figure 3(a) depicts a typical CMOS driver with its input waveform and RC-π load. The output voltage waveform may be modeled as a weighted linear sum of ramp and exponential waveforms as shown in Figure 3(b). We therefore assume that the actual $c_{eff}$ can be obtained as a weighted average of that obtained for the ramp output waveform and that obtained for the exponential output waveform.

In the following, we calculate $c_{eff}$ for ramp and exponential waveforms of the gate output voltage.

\[ c_{eff}^r (\theta) = G \left( t_i, \left( t_{in}, c_{eff}^r (\theta) \right), c_i, r_i, c_f \right) \]

\[ c_{eff}^e (\theta) = H \left( t_i, \left( t_{in}, c_{eff}^e (\theta) \right), c_i, r_i, c_f \right) \]

(5)

where $0 \leq \zeta \leq 1$ is the weighting factor for the linear combination of exponential and ramp waveforms. However, we have observed that when $\theta = 50\%$, then $\zeta = 0.5$ results in the minimum error between the iterative $c_{eff}$ equation in Eqn. (5) and the actual sign-off $c_{eff}$ value.

4.2 Calculating $c_{eff}$ in the CFO form

Suppose $t_{in}$, $c_i$, $r_i$, and $c_f$ in the CFO form are given as:

\[ t_n = t_{in} + \sum_{i=1}^{\infty} t_{in}AX_i + t_{in}n\Delta S_{nu} \]

(6)

\[ c_n = c_{in} + \sum_{i=1}^{\infty} c_{in}AX_i + c_{in}n\Delta S_{nu} \]

(7)

\[ r_n = r_{in} + \sum_{i=1}^{\infty} r_{in}AX_i + r_{in}n\Delta S_{nu} \]

(8)

\[ c_f = c_{in} + \sum_{i=1}^{\infty} c_{in}AX_i + c_{in}n\Delta S_{nu} \]

(9)

$\Delta S_n = \text{Dist} (\mu = 0, \sigma^2 = 1, k_i)$

$\Delta S_n = \text{Dist} (\mu = 0, \sigma^2 = 1, k_i)$

(10)

The effective capacitance for this problem generally becomes a complex random variable, i.e. $c_{eff}^\prime$. Therefore, we approximate it with its CFO form and the objective becomes to calculate the coefficients of $c_{eff}$ in the CFO form as well as the skewness of $\Delta S_{eff}$ as:

\[ c_{eff} = c_{eff}^r + \sum_{n=1}^{\infty} c_{eff}^e nAX_i + c_{eff}^n\Delta S_{eff} \]

(11)

Such that $E \left[ c_{eff} - F \left( t_i, c_{eff}, c_i, r_i, c_f \right) \right]^2$ is minimized

where $F$ is given in Eqn. (5). Theorem 2 presents the solution for calculating these unknown values.

Theorem 2: For a variational circuit, where $t_{in}$, $c_i$, $r_i$ and $c_f$ in the CFO form are written as in Eqs. (6)-(10), the coefficients of $c_{eff}$ in the CFO form (Eqn. (11)), can be calculated as:

\[ c_{eff, i} = F \left( t_i, t_{in}, c_{eff, i}, \dot{c}_{eff, i}, \ddot{c}_{eff, i}, c_i, r_i, c_f \right) \]

(12)

\[ c_{eff, j} = \left( \frac{\partial F}{\partial t_{in}} \right)_{t_{in}} + \left( \frac{\partial F}{\partial c_{eff}} \right)_{t_{in}} t_{in} + \left( \frac{\partial F}{\partial c_i} \right)_{t_{in}} c_i + \left( \frac{\partial F}{\partial r_i} \right)_{t_{in}} r_i + \left( \frac{\partial F}{\partial c_f} \right)_{t_{in}} c_f \]

(13)

\[ c_{eff, n} = \sqrt{c_{eff, n, 1}^2 + c_{eff, n, 2}^2} \]

(14)

Proof: It is omitted for brevity.

Now, based on the assumption made above, an iterative equation for actual $c_{eff}$ calculation for any $\theta$% point of the output transition time may be written as:

\[ c_{eff} = c_{eff}^r + \sum_{n=1}^{\infty} c_{eff}^e nAX_i + c_{eff}^n\Delta S_{eff} \]
\[ \Delta s_{ce} = \text{Dist} \left( \mu = 0, \sigma^2 = 1, \kappa = \frac{\sum (c_{eff,m+1})^3}{\left( \sum (c_{eff,m+1})^2 \right)^{3/2}} \right) \]

where:

\[ c_{eff,m+1} = \frac{\left( \frac{\partial F}{\partial \tau} \right)_{\text{nom},m+1}}{1 - \left( \frac{\partial F}{\partial \tau} \right)_{\text{eff},m+1}} \quad \text{and} \quad U = \left\{ t_{in}, c_{i}, t_{out}, c_{j} \right\} \]

\[ c_{eff,m+1} = \frac{\left( \frac{\partial F}{\partial \tau} \right)_{\text{nom},m+1}}{1 - \left( \frac{\partial F}{\partial \tau} \right)_{\text{eff},m+1}} \]

\[ c_{eff,m+1} = \frac{\left( \frac{\partial F}{\partial \tau} \right)_{\text{nom},m+1}}{1 - \left( \frac{\partial F}{\partial \tau} \right)_{\text{eff},m+1}} \]

\[ c_{eff,m+1} = \frac{\left( \frac{\partial F}{\partial \tau} \right)_{\text{nom},m+1}}{1 - \left( \frac{\partial F}{\partial \tau} \right)_{\text{eff},m+1}} \]

\[ \text{Proof:} \] It is omitted for brevity.

Eqs. (12) is the iterative \( c_{eff} \) calculation under the nominal conditions of the circuit. Hence, \( c_{eff,0} \) can be evaluated by using the conventional effective capacitance calculation [12][13].

\( t_{in}, c_{i}, t_{out}, c_{j} \), are given (cf. Eqs. (6)-(9)).

1. Updating the gate library look-up table and utilizing the additional data during gTA: The revised tables now provide not only the timing quantity for each combination of \( t_{in} \) and \( c_{j} \) but also the derivatives of the timing quantity \( t_{in} \) with respect to \( t_{in} \) and \( c_{j} \) for each combination of \( t_{in} \) and \( c_{j} \).

2. Using the existing gate library look-up table, but performing additional calculations during gTA: To approximate calculate \( (\partial F/\partial t_{in})_{\text{nom}} \) we read \( t_{in} \) (from the gate library) for \( <t_{in}\text{0} : c_{j}\text{0} > \) and \( <t_{in}\text{0} : c_{j}\text{0} > \). Next, we calculate \( \Delta t_{in} \delta \) as the approximation. \( (\partial F/\partial c_{j})_{\text{nom}} \) can be similarly calculated.

Using any of the above solutions, Eqs. (13) and (14) become closed form expressions, which can be evaluated in constant time. Note that we calculate \( (\partial F/\partial t_{in})_{\text{nom}} \), \( (\partial F/\partial c_{j})_{\text{nom}} \), \( (\partial F/\partial \tau)_{\text{nom}} \), and \( (\partial F/\partial c_{j})_{\text{nom}} \) only once and in constant time. Therefore, complexity of our method is dominated by the iterative effective capacitance calculation under the nominal conditions.

5. Experimental Results

Our experiments use 90nm CMOS process parameters to model gates and interconnect parasitics. We assumed two different configurations for the experimental setup. The first one consists of two inverters connected in series whereas the second one is a CMOS inverter followed by a 2-input NAND gate. For both configurations, we apply a ramp input to the first inverter while its nominal value is chosen from the set \( (t_{in})_{\text{nom}} = \{10\text{ps, 80ps, 150ps, 220ps, 300ps}\} \). For the first configuration, size of the first inverter is fixed at \( W_{p}/W_{n} = 30/15 \mu m \) whereas size of the second inverter is chosen to be one of \( W_{p}/W_{n} = [20/10, 50/25, 70/35, 100/50] \mu m \). For the second configuration, size of the first inverter is again fixed at \( W_{p}/W_{n} = 30/15 \mu m \) whereas this time the size of the succeeding 2-input NAND gate is chosen to be one of \( W_{p}/W_{n} = \{40/40, 50/50, 100/100\} \mu m \).

To characterize the timing behavior of the gate, a look-up table based library is employed which represents the gate delay and output transition time as a function of input rise time, output capacitive load, \( V_{dd} \), and temperature. We apply different loading scenarios for the second-stage gate as explained in the following subsections, i.e., pure capacitive load, and general RC load. We have also considered four different global sources of variation (\( V_{dd} \), temperature, Metal layer 1 width, and ILD) and one independent random sources of variation for each electrical parameter (i.e., \( r \) and \( c \)) and timing parameter (for instance \( t_{in} \) in the circuit). The sensitivity of each data to the sources of variation is chosen randomly, while the total \( \sigma \) variation for each data is chosen to be 10% and 15% of their nominal value.

We also assumed that the sources of variation are skewed with different skewness values as explained in each subsection. Mean, variance, and skewness of effective capacitance, the gate 50% propagation delay, and 10%-90% output transition time are calculated using the approaches presented in this paper.

To compare the results, we ran Monte Carlo simulation with \( 10^4 \) samples on each test scenario and derived mean, variance, and skewness of the effective capacitance, gate 50% propagation delay, and 10%-90% output transition time. Average percentage errors for the mean, variance, and skewness of effective capacitance, the gate 50% propagation delay, and 10%-90% output transition time between the obtained results from the Monte Carlo and the calculated results based on using statistical gate timing analysis approach are reported.

A. Purely Capacitive Load

The load in this section is considered to be purely capacitive. Its nominal value is chosen to be \( (C)_{\text{nom}} = \{400, 500, 800, 1400\}/F \). The scaled distribution of the sources of variation is considered to have a skewness of 0.4, 0.6, and 0.8. We performed our experiments on both circuit configurations explained above. The results for the first configuration (where the second gate is an inverter) are presented in Table 2 (the skewness of the given data is 0.4) and Table 3 (for the skewness of 0.8). The results for the second configuration are provided in Table 4 (for the skewness of 0.6). Experimental results indicate an average error of about 3% for two different \( \sigma \) values, i.e., 10% and 15%. As we increase the \( \sigma \) value, i.e., the total \( \sigma \) variation for each data, e.g., \( \sigma \) variation of \( t_{in} \) and \( c_{j} \) from 10% to 15%, the error in calculated mean, variance, and skewness of the delay and slew increase, but slightly. The sources of error can be mainly classified into two groups: 1) the inaccuracy of the gate library table lookup and 2) the linear first order approximation of the timing and electrical parameters with respect to the sources of variation. Note that, the runtime of the proposed algorithm in average is 89 times faster than the Monte Carlo based approach.

Table 2: Average error for the inverter driving pure capacitive load (Skewness=0.4)

<table>
<thead>
<tr>
<th>( \sigma ) Value</th>
<th>Delay</th>
<th>Slew</th>
<th>Delay</th>
<th>Slew</th>
</tr>
</thead>
<tbody>
<tr>
<td>Mean</td>
<td>1.5%</td>
<td>1.7%</td>
<td>2.2%</td>
<td>2.3%</td>
</tr>
<tr>
<td>Variance</td>
<td>1.2%</td>
<td>1.3%</td>
<td>1.8%</td>
<td>1.9%</td>
</tr>
<tr>
<td>Skewness</td>
<td>1.0%</td>
<td>1.1%</td>
<td>1.4%</td>
<td>1.3%</td>
</tr>
</tbody>
</table>
In this paper we presented a framework to handle the variation-aware gate timing analysis in block-based STA considering non-Gaussian sources of variation. First, we proposed an approach to calculate variational RC-π load, which can be utilized in place of the actual variational RC load for the gate timing analysis purposes. Next, we presented a new approach for calculating effective capacitance in STA. We used this technique to calculate the statistical $c_{eff}$ in the CFO form, and thereby, calculated the gate delay and output slew in the that form. Experimental results show an average error of 6% with respect to Monte Carlo with $10^4$ samples simulation.

7. References