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RISC Compiler and Simulator, Beta Release V0.3.0: Out-of-Order Parallel Simulatable SystemC Subset

Guantao Liu, Tim Schmidt, and Rainer Dömer

Technical Report CECS-16-06
September 30, 2016

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Abstract

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In recent years parallel SystemC simulators were proposed which run multiple threads in parallel based on synchronous Parallel Discrete Event Simulation (PDES) semantics. Synchronous PDES, however, limits parallel execution to threads that run at the same time and delta cycle.

In this report, we describe the advanced Recoding Infrastructure for SystemC (RISC) approach where a dedicated SystemC compiler and advanced parallel simulator implement Out-of-Order Parallel Discrete Event Simulation (OoO PDES) for SystemC. OoO PDES can execute threads in parallel and out-of-order (ahead of time) and thus achieves fastest simulation speed but nevertheless maintains the classic SystemC modeling semantics.

This report describes the RISC Compiler and Simulator and details the SystemC subset supported by the RISC Beta Release V0.3.0, as of September 30, 2016.

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1 Introduction

As an IEEE standard [1], the SystemC System Level Description Language (SLDL) is widely used for the specification, modeling, validation and evaluation of Electronic System Level (ESL) models. Under the Accellera Systems Initiative [2], the SystemC Language Working Group [3] maintains not only the official SystemC language definition, but also provides an open source proof-of-concept library [4] that can be used to simulate SystemC design models. However, implementing the classic scheme of Discrete Event Simulation (DES), this reference simulator runs sequentially and cannot utilize the parallel computing resources available on multi-core (or many-core) processor hosts. This severely limits the execution speed of SystemC simulation.

In order to provide faster simulation, Parallel Discrete Event Simulation (PDES) [5] has recently gained again significant attraction (examples include [6], [7], [8], [9], [10], and [11]). The PDES approach issues multiple threads (i.e. `SC_METHOD`, `SC_THREAD` and `SC_CTHREAD`) concurrently and runs them on the available processor cores in parallel. In turn, the simulation speed increases significantly.

Regular PDES is synchronous, however. That is, time advances globally and all threads execute in lock-step fashion. Here, the total order of time imposed by synchronous PDES still limits the opportunities for parallel execution. When a thread completes its evaluation phase, it has to wait until all other threads finish their evaluation

phases as well. Earlier completed threads must stop at the simulation cycle barrier and available processor cores are left idle until all runnable threads reach the cycle barrier.

In order to overcome this problem, we have developed a novel technique called Out-of-Order Parallel Discrete Event Simulation (OoO PDES) [12, 13, 14, 15]. By localizing the simulation time to individual threads and carefully handling events at different times, the simulation kernel can issue threads in parallel and ahead of time, following a partial order of time without loss of accuracy. Thus, Ooo PDES significantly reduces the idle time of available parallel processor cores and results in maximum simulation speed, while maintaining the traditional language and modeling semantics.

The OoO PDES technique was originally implemented based on the SpecC language [16, 17, 18, 19]. In this report, we document our efforts to apply OoO PDES to the SystemC SLDL [20, 21, 1] which is both the de-facto and official standard for ESL design today. In particular, we describe our Recoding Infrastructure for SystemC (RISC) [22] which consists of a dedicated SystemC compiler and corresponding out-of-order parallel simulator and implements OoO PDES for SystemC.

The remainder of this report is organized as follows: After a brief description of the simulator scheduling algorithms used for DES, PDES and OoO PDES in Section 2, we describe the RISC Compiler and Simulator proof-of-concept prototype in Section 3. In Section 4, we then list in detail the SystemC subset that is supported by the current RISC Beta Release V0.3.0 (2016-09-30)[1] and finally conclude this report in Section 5.

2 Out-of-Order Parallel Simulation

In this section, we briefly outline the scheduling algorithm used in out-of-order parallel simulation. We do this incrementally, starting from the traditional Discrete Event Simulation (DES) scheduler, then describe the synchronous Parallel DES (PDES) extension, and finally define the Out-of-Order PDES (OoO PDES) scheduling algorithm.

2.1 Notations

To formally describe the discrete event scheduling algorithms, the following notations are introduced.

1. Each SystemC thread (`SC_METHOD`, `SC_THREAD` and `SC_CTHREAD`) is assigned a localized time stamp (t_{th}, δ_{th}).

2. Each event (`sc_event`) is assigned a notification time stamp (t_e, δ_e), where $EVENTS = \cup EVENTS_{t,\delta}$.

3. Threads are grouped into different queues. Specifically,

   (a) $QUEUES = \{READY, RUN, WAIT, WAITTIME\}$.

   (b) $READY = \cup th_{t,\delta}$ where Thread th is ready to run at time (t, δ).

   (c) $RUN = \cup th_{t,\delta}$ where Thread th is running at time (t, δ).

   (d) $WAIT = \cup th_{t,\delta}$ where Thread th is waiting since time (t, δ).

   (e) $WAITTIME = \cup th_{t,0}$ where Thread th is waiting for simulation time advance to $(t, 0)$.

[1] An earlier version [23] of this technical report documents the prior Alpha Release V0.2.1 (2015-10-30).

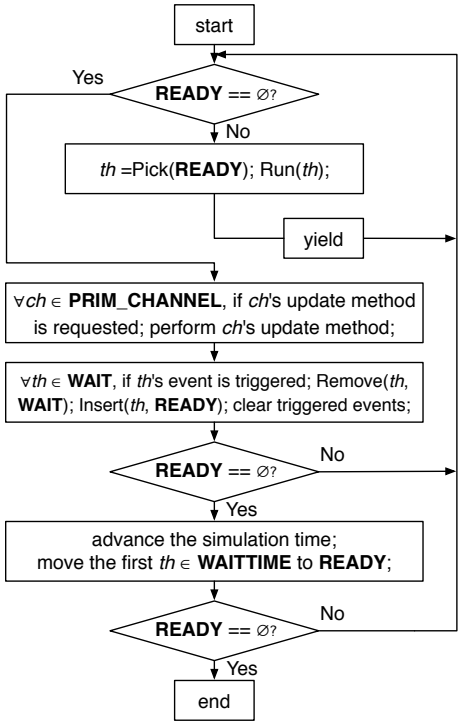

Figure 1: Traditional Discrete Event Simulation (DES) scheduler for SystemC.

2.2 Discrete Event Scheduler

The Accellera reference simulation library of SystemC [4] is based on DES. Figure 1 depicts such a traditional DES scheduling algorithm. In DES, a single thread is running at all times. When all threads in the *READY* and *RUN* queues complete their current delta cycle, the root thread resumes and performs the update and notification phase. Then threads are woken up and moved from the *WAIT* queue back into the *READY* queue. A new delta cycle begins.

If no threads are ready after the update and notification phase, the current time cycle finishes. The simulation kernel advances the simulation time and processes the earliest timed event from the *WAITTIME* queue. A new cycle begins for the updated simulated time.

Finally, when both the *WAITTIME* and *READY* queues are empty, the simulation terminates.

2.3 Parallel Discrete Event Scheduler

In comparison to DES, regular synchronous PDES issues multiple threads (`SC_METHOD`, `SC_THREAD` and `SC_CTHREAD`) concurrently in a delta cycle. These threads can then execute truly in parallel on the multiple available processor cores of the host.

Figure 2 shows the regular synchronous PDES scheduling algorithm. In the evaluation phase, as long as the *READY* queue is not empty and an idle core is available, the PDES scheduler will issue a new thread from the *READY* queue. If a thread finishes earlier than other threads in the same cycle, a new ready thread is assigned to the idle processor core, unless there is no thread available in the *READY* queue, in which case the core is keept idle until the next delta cycle.

It should be emphasized that synchronous PDES implies an absolute barrier at the end of each delta and time cycle. All threads need to wait at the barrier until all other runnable threads finish their current evaluation phase.

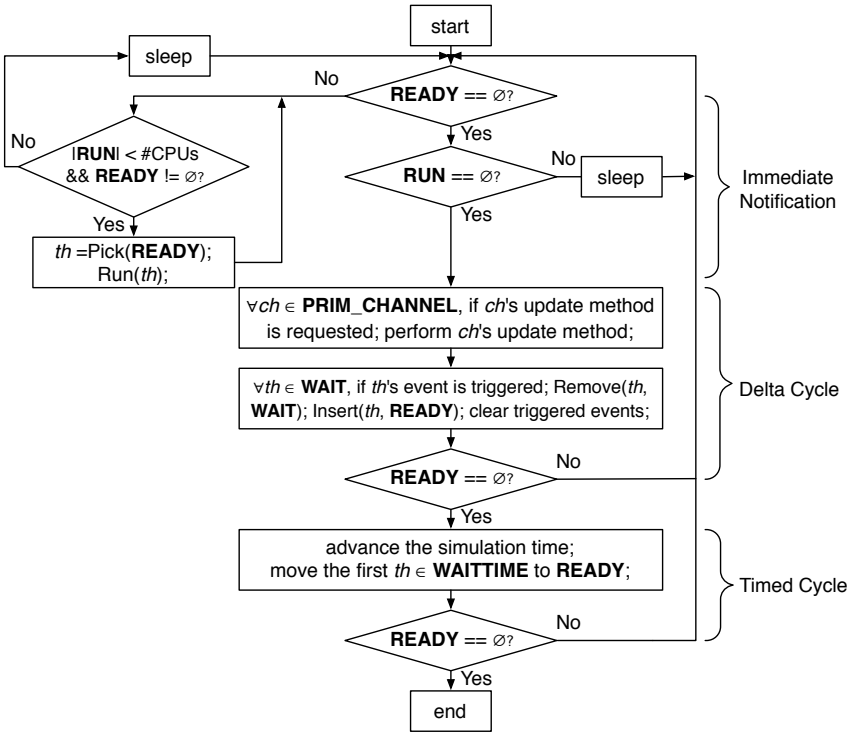

Figure 2: Synchronous Parallel Discrete Event Simulation (PDES) scheduler for SystemC.

Only then the synchronous PDES scheduler resumes and performs the update and notification phases, and finally advances to the next delta or time cycle.

For the SystemC language in particular, there is a very important aspect to consider when applying PDES. For semantics-compliant SystemC simulation, complex inter-dependency analysis over all variables in the system model is a prerequisite to parallel simulation [24].

The Standard SystemC Language Reference Manual (LRM) [1] clearly states that *"process instances execute without interruption"*. This requirement is also known as cooperative (or co-routine) multitasking which is assumed by the SystemC execution semantics. As detailed in [24], the particular problem of parallel simulation is specifically addressed in the SystemC LRM [1]:

> *"An implementation running on a machine that provides hardware support for concurrent processes may permit two or more processes to run concurrently, provided that the behavior appears identical to the co-routine semantics defined [...]. In other words, the implementation would be obliged to analyze any dependencies between processes and constrain their execution to match the co-routine semantics."*

We will describe the required dependency analysis in more detail below (in Section 3.2), as it is also needed for out-of-order PDES.

2.4 Out-of-Order Parallel Discrete Event Scheduler

In OoO PDES, we break the strict order of time (the synchronous barrier) by localizing time stamps to each thread. Figure 3 shows the out-of-order parallel DES scheduling algorithm. Since each thread has its own time stamp, the OoO PDES scheduler relaxes the event and simulation time updates, allowing more threads (at

different simulation cycles!) to run in parallel and ahead of time. This results in a higher degree of parallelism and thus higher simulation speed.

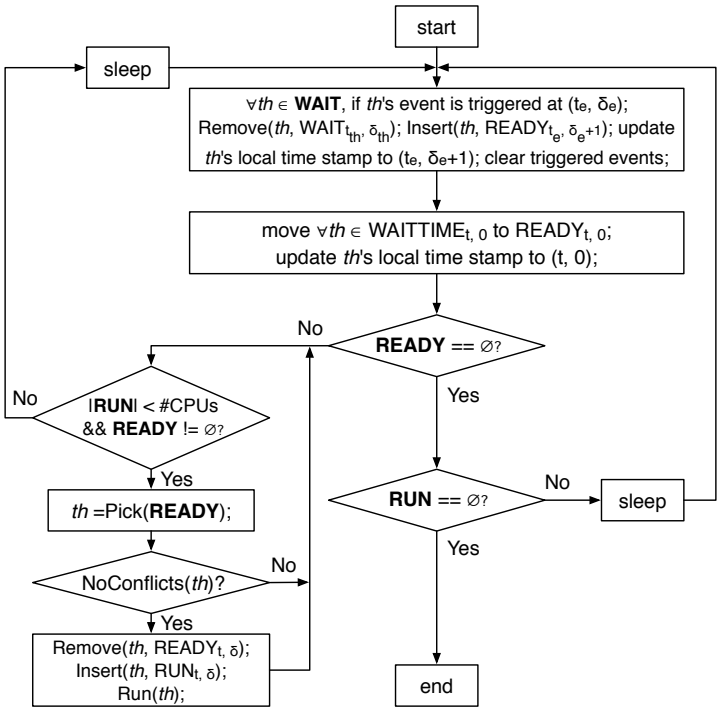

Figure 3: Out-of-Order Parallel Discrete Event Simulation (OoO PDES) scheduler for SystemC.

In comparision to the synchronous PDES in Figure 2, Figure 3 moves threads from the *WAIT* and *WAITTIME* queues into the *READY* queue *as soon as possible*. Also, there is no specific point in the scheduling flow any more for the classic delta and time cycles. Both delta and time updates are performed locally for each thread, provided that there are no possible conflicts in the way (the $NoConflicts(th)$ condition is explained below).

In contrast to Figure 2 which performs requested update methods in primitive channels in each delta cycle, Figure 3 does not contain this step any more. Due to the out-of-order scheduling and the eliminated central scheduling point for delta cycles, it is difficult to determine an efficient and safe point in the OoO PDES scheduler when primitive channel update requests can be served. However, it is always possible to safely fall back to synchronous PDES when primitive channel updates are requested.

Note the $NoConflicts(th)$ condition shown in Figure 3. As already mentioned above for the synchronous PDES, detailed dependency analysis is needed to avoid data or event conflicts for any shared variables among the parallel threads. Only if $NoConflicts(th)$ is true, a new thread is issued for parallel execution (moved from the *READY* to the *RUN* queue).

We will be using advanced static compile-time analysis (and optionally dynamic run-time analysis, see Section 3.2.2) to identify all such potential conflicts. Based on this information (a simple table lookup is sufficient), the OoO PDES scheduler can then at run-time quickly decide whether or not a set of threads has any conflicts with each other.

3 RISC Compiler and Simulator

To realize the OoO PDES approach for the SystemC language, we present now our Recoding Infrastructure for SystemC (RISC) and describe the overall RISC Compiler and Simulator proof-of-concept prototype (Beta Release V0.3.0 as of 2016-09-30). The RISC software is available as open source and can be downloaded freely from the following web site [22]: http://www.cecs.uci.edu/~doemer/risc.html

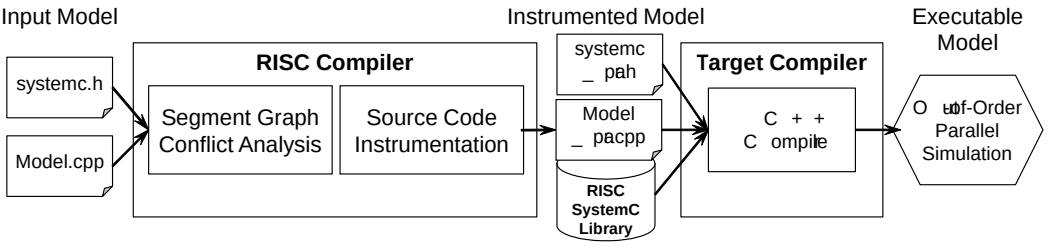

Figure 4: RISC Compiler and Simulator for Out-of-Order PDES of SystemC.

To perform semantics-compliant SystemC simulation with maximum parallelism, we introduce a *dedicated SystemC compiler*. This is in contrast to the traditional SystemC simulation where a regular SystemC-agnostic C++ compiler includes the SystemC headers and links the input model directly against the SystemC library.

As shown in Figure 4, our RISC compiler acts as a frontend that processes the input SystemC model and generates an intermediate model with special instrumention for OoO PDES. The instrumented parallel model is then linked against the extended RISC SystemC library by the target compiler (a regular C++ compiler) to produce the final executable output model. OoO PDES is then performed simply by running the generated executable model.

From the user perspective, we essentially replace the regular SystemC-agnostic C++ compiler with the SystemC-aware RISC compiler (which in turn calls the underlying C++ compiler). Otherwise, the overall SystemC validation flow remains the same as before. It is just faster due to the parallel simulation.

For reference, the detailed Linux manual page of the RISC compiler `risc` and simulator is included in Appendix A.1 of this report.

Internally, the RISC compiler performs three major tasks, namely Segment Graph construction, conflict analysis, and source code instrumentation.

3.1 Segment Graph

The first task of the RISC compiler is to parse the SystemC input model into an abstract syntax tree (AST) and then create a SystemC structural representation from the AST which reflects the SystemC module and channel hierarchy, connectivity, and other SystemC-specific relations, similar to the SystemC-clang representation [25, 26]. For details on this part of the RISC application programming interface (API), please refer to the Doxygen-generated documentation [27].

On top of this, the RISC compiler then builds a Segment Graph data structure for the model. A Segment Graph (SG) [12] is a directed graph that represents the code segments executed during the simulation between scheduling steps. That is, every segment is associated with a scheduler entry point, i.e. a `wait` statement in SystemC.

At run time, threads switch back and forth between the states of *running* (threads in *READY* and *RUN* queues) and *waiting* (threads in *WAIT* and *WAITTIME* queues). When *running*, they execute specific segments of their code. These code segments make up the nodes in the Segment Graph, whereas edges in the graph indicate the possible transitions from one segment to another (an abstraction of the model's control flow).

For a formal description of the Segment Graph and it's construction algorithm, the interested reader may refer to [15]. For details on the RISC API, please refer to the Doxygen-generated documentation [27].

3.2 Conflict Analysis

The Segment Graph data structure serves as the foundation for segment *conflict analysis*. As outlined earlier, the OoO PDES scheduler must ensure that every parallel thread to be issued has no conflicts with any other threads currently in the *READY* and *RUN* queues. Here, we utilize the RISC compiler to detect any possible conflicts between these threads already at compile time.

Potential conflicts in SystemC include data hazards, event hazards, and timing hazards, all of which may exist among the segments executed by the threads considered for parallel execution. Please refer to [15] for a detailed discussion of these hazards which, if ignored, would become race conditions at run time.

Both possible hazard detection approaches, namely *static* analysis at compile time and *dynamic* analysis at run time, are supported by RISC Compiler and Simulator Beta ReleaseV0.3.0.

3.2.1 Static Analysis

Static analysis relies purely on the available information in the SystemC source code of the design model at hand. In this case, the RISC compiler performs very conservative identification of the potential hazards in the model.

Identifying all possible hazards is a complex analysis task that requires the full "understanding" of the module hierarchy. One option is to statically extract the module hierarchy and analyze the individual threads. Here, the RISC compiler follows the approach outlined in [15].

However, in most cases not all of the needed information can be gathered statically. For instance, design parameters may be passed via the command line, for example, to define the number of modules, certain channels characteristics, or other configuration information. In such SystemC models, the instantiated modules, channels, and ports are typically created through loops in a dynamic fashion. However, these exact parameters are only available at run time, so they cannot be statically analyzed. In these cases, dynamic analysis is needed.

3.2.2 Dynamic Analysis

Dynamic analysis takes run-time information into account and then augments the classic static analysis. The combination of static and dynamic analysis is often called *hybrid* analysis [28].

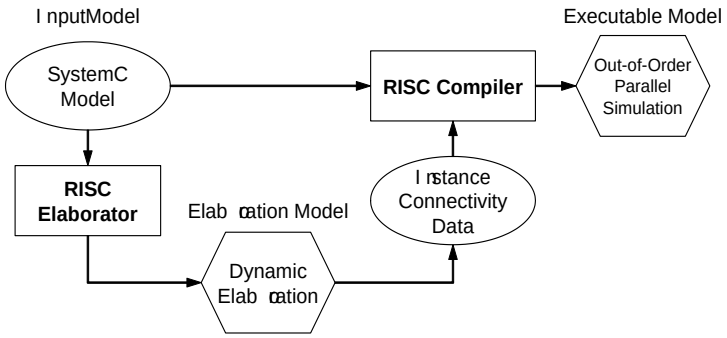

Figure 5: RISC Elaborator feeds dynamic elaboration information to RISC Compiler for precise conflict analysis.

Figure 5 shows the extended RISC design flow with support of dynamic analysis. As in the regular compilation flow discussed above in Figure 4, the input SystemC model is processed by the RISC Compiler to generate an executable model for out-of-order parallel simulation, as shown on the top half of Figure 5 from left to right.

The dynamic analysis step, shown on the bottom half of Figure 5, extends the compilation flow by a preprocessing step. The input SystemC model is fed into the RISC Elaborator `elab` which produces an executable model that only performs the SystemC elaboration phase when run. At the end of the elaboration, the executable model automatically traverses the created module hierarchy via the SystemC introspection API and dumps this detailed structural design information, shown as Instance Connectivity Data in Figure 5, into a file (*`model_name`*`.elab`). This file is in turn provided as an input to the RISC compiler, so that the dynamically created design hierarchy and specific instance connectivity can be used for precise conflict analysis. The instance connectivity data file includes the actual module hierarchy, the specific port mapping, and the actual target variable mapping of references.

Note that the use of the RISC Elaborator is optional. Design models, that can fully be analyzed statically, can be fed directly into the RISC Compiler without any preprocessing by the RISC Elaborator.

For reference, the detailed Linux manual pages of the RISC Compiler `risc` and RISC Elaborator `elab` are included in Appendix A.1 and Appendix A.2, respectively.

3.3 Source Code Instrumentation

As a result of the conflict analysis (static, dynamic, or hybrid [28]), the RISC compiler generates several conflict tables that describe all possible conflicts between threads in any two segments. Using this conservative information, the simulator can then at run-time quickly determine by a simple table look-up whether or not it is safe to issue any given thread in parallel or ahead of time.

As shown above in Figure 4, the RISC compiler and simulator work closely together. The compiler performs conservative conflict analysis and passes the analysis results to the simulator which then can make safe scheduling decisions quickly.

To pass information from the compiler to the simulator, we use automatic model instrumentation. That is, the intermediate model generated by the compiler contains instrumented (automatically generated) source code which the simulator then can rely on. At the same time, the RISC compiler also instruments user-defined SystemC channels with automatic protection against race conditions among communicating threads.

In total, the RISC source code instrumentation includes four major components:

1. Segment and instance IDs: Individual threads are uniquely identified by a creator instance ID and their current code location (segment ID). Both IDs are passed into the simulator kernel as additional arguments to scheduler entry functions, including `wait` and thread creation.

2. Data and event conflict tables: Segment concurrency hazards due to potential data conflicts, event conflicts, or timing conflicts are provided to the simulator as two-dimensional tables indexed by a segment ID and instance ID pair. For efficiency, these table entries are filtered for scope, instance path, and reference and port mappings.

3. Current and next time advance tables: The simulator can make better scheduling decisions by looking ahead in time if it can predict the possible future thread states. This optimization is discussed in detail in [14]. However, it is not supported by the current RISC Compiler and Simulator V0.3.0 at this time.

4. User-defined channel protection: SystemC allows the user to design channels for custom inter-thread communication. To ensure such communication is safe also in the OoO PDES situation where threads execute truly in parallel, the RISC compiler automatically inserts locks (binary semaphores) into these channels so that mutually-exclusive execution of the channel methods is guaranteed. Otherwise, race conditions could exist when communicating threads exchange data.

Note that the source code instrumentation is performed automatically by the RISC Compiler and no user-interaction is necessary. However, the interested user may inspect the instrumented source code. It is stored in a file named `risc_model_name.cpp` which serves as the input file to the compiler backend which in turn then generates the final executable.

3.4 Library Support

There exists a significant limitation for the described conflict analysis and source code instrumentation. It only works if the compiler has access to the entire source code of the design model. This is typically fine for smaller SystemC benchmark examples, but does not hold true for more complex SystemC models where multiple translation units and/or library files are used. In these cases, the compiler has access only to the function signatures (function declarations in header files), but not to their implementation (function bodies which are precompiled in the library file). Thus, the compiler cannot analyze the function bodies for potential conflicts, neither can it instrument any segment boundaries (i.e. `wait` calls) in the library code with segment and instance IDs.

In its previous Alpha version [23], the RISC Compiler and Simulator operated under the assumption that all library code is thread-safe without any conflicts, and does not contain any segment boundaries (no `wait`). This is reasonable for the standard C/C++ libraries used in a modern Linux environment, as well as for the specially prepared RISC SystemC simulator library. However, this assumption posed a significant limitation for more complex SystemC models built around custom application libraries.

Now, RISC Compiler and Simulator Beta Release V0.3.0 offers support for library code by use of *function annotations*. The RISC annotation scheme for library functions provides abstract information for both conflict analysis and segment boundaries.

Specifically, the user can annotate function declarations with `pragma` statements which specify whether or not the function poses any potential conflicts, and what type of `wait` calls the function body contains. For example, the standard math function `sqrt` and the blocking `read` function of the SystemC `sc_fifo` channel are annotated as follows:

```
// standard math square-root function
#pragma RISC sqrt conflict-free no-wait
double sqrt(double x);

// sc_fifo blocking read function
#pragma RISC read conflict-free looped-wait event
virtual T read();
```

Here, the `sqrt` function is declared `conflict-free` because it is thread-safe and has no dangerous side effects. Since this is true for many functions (e.g. most functions in the C standard library), the RISC Compiler assumes this by default. Thus, this `pragma` statement is not explicitly needed.

The `sc_fifo::read` function is also declared `conflict-free` because it operates in a standard SystemC channel that is safely protected by a lock in the RISC simulator library. However, this blocking `sc_fifo::read` function is annotated as `looped-wait` because it contains a `wait` statement in the body of a loop that is waiting for available data, which is indicated by some `event`. Thus, the RISC Compiler can take this segment boundary into account when analyzing a call to this function.

In general, a function is considered `conflict-free` if the corresponding function body contains no potential read/write access conflicts to any shared state with the other threads in the simulation model. Otherwise, it must be annotated as `not-conflict-free`.

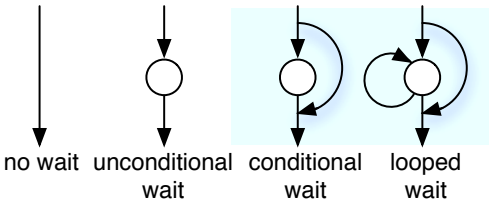

Figure 6: Control-flow abstractions for `wait` in library functions.

For the annotation of segment boundaries contained in library functions, Figure 6 shows the different control-flow abstractions with regards to `wait` function calls in the corresponding function body. In the first case, the function contains no `wait` statement and is a non-blocking function in SystemC simulation. The next two cases apply to functions with a conditional or a non-conditional `wait`. The last case covers the possible encounter of multiple `wait` statements, such as the blocking `read` call to a `sc_fifo` discussed above.

Finally, the last parameter to the `pragma` statement specifies the type of the `wait` statement in the function body, either `event` for waiting for any notified event, or the minimum time increment that the simulator will incur when executing the corresponding function, such as `sc-zero-time` or `(42,SC_MS)`.

3.5 Compiler Backend

After the automatic source code instrumentation, the RISC compiler passes the generated intermediate model in file `risc_`*`model_name`*`.cpp` to the underlying regular C++ compiler. That target compiler then produces the final simulation executable by linking the instrumented code against the RISC extended SystemC library.

By default, the RISC Compiler and Simulator rely on the GNU C++ compiler `g++` for the backend code generation. Alternatively, the Intel C++ compiler `icpc` may be used to generate a simulation executable that is optimized for Intel processors with Single-Instruction-Multiple-Data (SIMD) capabilities or the Intel Many-Integrated-Core (MIC) architecture. Please refer to the command-line options `-risc:icpc` and `-risc:mic`, respectively, which are documented in the manual pages for `risc` (see Appendix A.1) and `elab` (see Appendix A.2).

3.6 Simulator

Same as the classic Accellera proof-of-concept implementation [4], the RISC simulator is not an explicit tool, but a run-time library [29] that the generated executable SystemC model is linked against. Thus, simulation is performed by execution of the compiled model, the same way as before (just faster).

The RISC simulator identifies itself by its log message at the beginning of the simulation run, announcing `OoOPARALLEL` execution after the SystemC language version number (`SystemC 2.3.1`). It also adds the Center for Embedded and Cyber-physical Systems (CECS) as a contributor to the RISC-extended SystemC library.

A simple *HelloWorld* model is shown running in the following example:

```
sh % ./HelloWorld

        SystemC 2.3.1-OoOPARALLEL --- Sep 21 2016 11:07:46
        Copyright (c) 1996-2016 by CECS and all Contributors,
        ALL RIGHTS RESERVED
```

```
Hello World!
```

There are two environment variables that the out-of-order parallel SystemC library is sensitive to. First, the variable `SYSC_PAR_SIM_CPUS` specifies the maximum number of parallel threads allowed in out-of-order parallel simulation (namely `#CPUs` in Figure 3). For efficient simulation, this variable should be set to a value suitable for the simulation host, e.g. the number of available CPU cores. If unset, `SYSC_PAR_SIM_CPUS` defaults to 64.

Second, the environment variable `SYSC_SYNC_PAR_SIM` can be used to force the default out-of-order parallel scheduler to fall-back to synchronous parallel execution. By default (when undefined), `SYSC_SYNC_PAR_SIM` is assumed to be `false`, so out-of-order parallel simulation (OoO PDES) is performed. On the other hand, if `SYSC_SYNC_PAR_SIM` is set to true, the simulator will execute in synchronous PDES fashion.

As indicated above in Section 2.4, the RISC simulator Beta Release V0.3.0 (2016-09-30) also falls back to synchronous execution as soon as primitive SystemC channels are used with requests to update functions. Thus, such models will execute in safe synchronous manner.

4 Out-of-Order Parallel Simulatable SystemC Subset

Over more than a decade, the SystemC language [21], which technically is a C++ application programming interface (API) with a corresponding simulation library, has evolved from basic constructs for modeling parallel modules connected by signals and channels to a highly complex set of macros, types, classes, templates, and functions for very advanced modeling (i.e. Transaction Level Modeling (TLM) 2.0 [30, 31]) and highly optimized simulation of SystemC models. Usually these optimizations have aimed at higher simulation speed, i.e. by minimizing context switches in the simulator, or at higher levels of abstraction due to purposely relaxed timing. Often, the uninterrupted (sequential) execution semantics on a single processor host have been assumed or are explicitly required.

Along these lines, it has been recognized that there is considerable need to study and adjust or *evolve* the SystemC language towards better support of parallel execution (following some form of suitable PDES semantics). One example of the ongoing discussion within the SystemC community is a presentation at the SystemC Evolution Day 2016 where significant obstacles in the current language standard have been identified [32].

In contrast to the current SystemC standard [1], RISC now aims for truly parallel execution on multi- or many-core hosts. Changing the fundamental assumptions about SystemC simulator execution consequently may affect some constructs and APIs which need to be revisited and evaluated anew. The goal of this section is to start this process and enable fruitful discussions.

Below, we describe and list the out-of-order parallel simulatable SystemC subset supported by the current RISC Compiler and Simulator, Beta Release V0.3.0. In particular, Table 1 through Table 8 list for each SystemC construct whether or not it is supported at this time. If applicable, an explanation note is provided that briefly outlines the status and/or the plans for the given feature.

Overall, our current RISC proof-of-concept prototype supports the classic SystemC constructs for hierarchical modeling and multi-threaded execution, but many advanced features are not supported yet or left undecided at this stage. The status "undecided" in particular indicates that further study is needed to decide whether or not the given construct can be supported in efficient and reasonable manner by RISC and its OoO PDES approach.

4.1 SystemC Hierarchical Structure of Modules and Channels

RISC supports the regular hierarchical and structural composition of the SystemC design model. This includes the SystemC program start (`sc_main`, `sc_start`) and the general composition (`SC_CTOR`) of modules (`sc_module`, `SC_MODULE`, `sc_behavior`) and channels (`sc_channel`, `sc_prim_channel`).

Table 1: RISC V0.3.0 Out-of-Order Parallel Simulatable SystemC Subset

Name	Type	Supported or not	Notes
sc_abs	function	Undecided	This function may not work with some arithmetic SystemC datatypes.
sc_actions	typedef	Supported	typedef unsigned sc_actions
sc_argc	function	Supported	
sc_argv	function	Supported	
sc_assemble_vector	function	Undecided	Work on this function in the future
sc_assert	macro	Undecided	Work on this macro in the future
sc_attr_base	class	Undecided	Work on this class in the future
sc_attr_cltn	class	Undecided	Work on this class in the future
sc_attribute	class	Undecided	Work on this class in the future
sc_behavior	typedef	Supported	typedef sc_module sc_behavior
sc_bigint	class template	Supported	
sc_biguint	class template	Supported	
sc_bind_proxy	class	Supported	
sc_bind	macro	Undecided	Work on this macro in the future
sc_bit	type (deprecated)	Undecided	Work on this type in the future
sc_bitref_r	class template	Undecided	Work on this class template in the future
sc_bitref	class template	Undecided	Work on this class template in the future
sc_buffer	class	Supported	
sc_bv_base	class	Undecided	Work on this class in the future
sc_bv	class template	Undecided	Work on this class template in the future
sc_channel	class	Supported	
sc_clock	class	Not Supported Now	sc_clock::before_end_of_elaboration() calls sc_spawn().
sc_close_vcd_trace_file	function	Undecided	Work on this function in the future
sc_concatref	class	Undecided	Work on this class in the future
sc_concref_r	class template	Undecided	Work on this class template in the future
sc_context_begin	enumeration	Supported	
sc_copyright	function	Supported	
sc_cor	class	Supported	
sc_cor_pkg	class	Supported	
sc_cor_pthread	class	Supported	
sc_cor_pkg_pthread	class	Supported	
sc_create_vcd_trace_file	function	Undecided	Work on this function in the future
sc_cref	macro	Undecided	Work on this macro in the future
sc_cthread_process	class	Supported	
SC_CTHREAD	macro	Supported	The risc compiler can generate the segment graph for SC_CTHREAD, however, it cannot handle the clock.
SC_CTOR	macro	Supported	

Table 2: RISC V0.3.0 Out-of-Order Parallel Simulatable SystemC Subset (continued)

Name	Type	Supported or not	Notes
sc_cycle	function (deprecated)	Not Supported Now	sc_cycle() calls sc_simcontext::cycle(), which is not supported in the out-of-order simulation in the current release.
sc_delta_count	function	Supported	This function returns the local delta count of the running process.
sc_elab_and_sim	function	Supported	
sc_end_of_simulation_invoked	function	Undecided	Work on this function in the future
sc_event_and_expr	class	Supported	Initial support as of v0.3.0
sc_event_and_list	class	Supported	Initial support as of v0.3.0
sc_event_finder_t	class template	Undecided	Work on this class template in the future
sc_event_finder	class	Undecided	Work on this class in the future
sc_event_or_expr	class	Supported	Initial support as of v0.3.0
sc_event_or_list	class	Supported	Initial support as of v0.3.0
sc_event_queue_if	class	Supported	
sc_event_queue	class	Not Supported Now	The constructor function is not supported by the out-of-order simulation in the current release.
sc_event	class	Supported	The immediate notification is not supported by the out-of-order simulation in the current release.
sc_exception	typedef	Undecided	Work on this typedef in the future
sc_export_base	class	Not Supported Now	No port following in compiler analysis
sc_export	class	Not Supported Now	No port following in compiler analysis
sc_fifo_blocking_in_if	class	Supported	
sc_fifo_in_if	class	Supported	
sc_fifo_in	class	Supported	
sc_fifo_nonblocking_in_if	class	Supported	
sc_fifo_out_if	class	Supported	
sc_fifo_out	class	Supported	
sc_fifo	class	Limited Support	sc_fifo::trace() and sc_fifo::operator = are not supported in this release; execution falls back to synchronous PDES
sc_find_event	function	Undecided	Work on this function in the future
sc_find_object	function	Undecided	Work on this function in the future
sc_fix_fast	class	Undecided	Work on this class in the future
sc_fix	class	Supported	
sc_fixed_fast	class template	Undecided	Work on this class template in the future
sc_fixed	class template	Supported	

Table 3: RISC V0.3.0 Out-of-Order Parallel Simulatable SystemC Subset (continued)

Name	Type	Supported or not	Notes
SC_FORK	macro	Undecided	Work on this macro in the future
sc_fxcast_context	class	Undecided	Work on this class in the future
sc_fxcast_switch	class	Undecided	Work on this class in the future
sc_fxnum_bitref	class	Undecided	Work on this class in the future
sc_fxnum_fast_bitref	class	Undecided	Work on this class in the future
sc_fxnum_fast_subref	class	Undecided	Work on this class in the future
sc_fxnum_fast	class	Undecided	Work on this class in the future
sc_fxnum_subref	class	Undecided	Work on this class in the future
sc_fxnum	class	Supported	
sc_fxtype_context	class	Undecided	Work on this class in the future
sc_fxtype_params	class	Undecided	Work on this class in the future
sc_fxval_fast	class	Undecided	Work on this class in the future
sc_fxval	class	Undecided	Work on this class in the future
sc_gen_unique_name	function	Undecided	Work on this function in the future
sc_generic_base	class	Undecided	Work on this class in the future
sc_get_curr_process_handle	function (deprecated)	Supported	
sc_get_current_process_handle	function	Supported	
sc_get_default_time_unit	function (deprecated)	Supported	
sc_get_status	function	Supported	
sc_get_stop_mode	function	Supported	
sc_get_time_resolution	function	Supported	
sc_get_top_level_events	function	Undecided	Work on this function in the future
sc_get_top_level_objects	function	Undecided	Work on this function in the future
SC_HAS_PROCESS	macro	Supported	
sc_hierarchical_name_exists	function	Undecided	Work on this function in the future
sc_in_clk	typedef	Supported	
sc_in_resolved	class	Supported	
sc_in_rv	class	Supported	
sc_in	class	Supported	sc_in::add_trace() and other tracing functions are not supported by the out-of-order simulation in the current release.
sc_in<bool>	class	Supported	sc_in<bool>::add_trace() and other tracing functions are not supported by the out-of-order simulation in the current release.
sc_in<sc_dt::sc_logic>	class	Supported	sc_in<sc_dt::sc_logic>::add_trace() and other tracing functions are not supported by the out-of-order simulation in the current release.

Table 4: RISC V0.3.0 Out-of-Order Parallel Simulatable SystemC Subset (continued)

Name	Type	Supported or not	Notes
sc_initialize	function (deprecated)	Supported	
sc_inout_clk	type (deprecated)	Supported	
sc_inout_resolved	class	Supported	
sc_inout_rv	class	Supported	
sc_inout	class	Supported	
sc_int_base	class	Supported	
sc_int_bitref_r	class	Undecided	Work on this class in the future
sc_int_bitref	class	Undecided	Work on this class in the future
sc_int	class template	Supported	
sc_interface	class	Supported	
sc_interrupt_here	function	Undecided	Work on this function in the future
sc_is_prerelease	function	Undecided	Work on this function in the future
SC_IS_PRERELEASE	macro	Supported	
sc_is_running	function	Supported	
sc_is_unwinding	function	Supported	
SC_JOIN	macro	Undecided	Work on this macro in the future
sc_length_context	class	Undecided	Work on this class in the future
sc_length_param	class	Undecided	Work on this class in the future
sc_logic	class	Undecided	Work on this class in the future
sc_lv_base	class	Undecided	Work on this class in the future
sc_lv	class template	Undecided	Work on this class template in the future
sc_main	function	Supported	
sc_max_time	function	Not Supported Now	This function is not supported by the out-of-order simulation in the current release.
sc_max	function	Supported	
sc_method_process	class	Supported	
SC_METHOD	macro	Supported	
sc_min	function	Supported	
sc_module_name	class	Supported	
sc_module	class	Supported	
SC_MODULE	macro	Supported	
sc_mutex_if	class	Not Supported Now	This class is not supported by the risc compiler in the current release.
sc_mutex	class	Not Supported Now	This class is not supported by the risc compiler in the current release.
sc_object	class	Supported	
sc_out_clk	type (deprecated)	Supported	

Table 5: RISC V0.3.0 Out-of-Order Parallel Simulatable SystemC Subset (continued)

Name	Type	Supported or not	Notes
sc_out_resolved	class	Supported	
sc_out_rv	class	Supported	
sc_out	class	Supported	
sc_pause	function	Undecided	Work on this function in the future
sc_pending_activity_at_current_time	function	Undecided	Work on this function in the future
sc_pending_activity_at_future_time	function	Undecided	Work on this function in the future
sc_pending_activity	function	Undecided	Work on this function in the future
sc_phash	class (deprecated)	Undecided	Work on this class in the future
sc_plist	class (deprecated)	Undecided	Work on this class in the future
sc_port	class	Supported	
sc_port_base	class	Supported	
sc_ppq	class (deprecated)	Undecided	Work on this class in the future
sc_prim_channel	class	Supported	sc_prim_channel::update() is performed in synchronous manner; execution falls back to synchronous PDES
sc_process_b	type (deprecated)	Supported	
sc_process_handle	class	Supported	
sc_pvector	class (deprecated)	Undecided	Work on this class in the future
sc_ref	macro	Undecided	Work on this macro in the future
sc_release	function	Supported	
sc_report_handler_proc	typedef	Undecided	Work on this typedef in the future
sc_report_handler	class	Undecided	Work on this class in the future
sc_report	class	Undecided	Work on this class in the future
sc_semaphore_if	class	Not Supported Now	This class is not supported by the risc compiler in the current release.
sc_semaphore	class	Not Supported Now	This class is not supported by the risc compiler in the current release.
sc_sensitive_neg	class (deprecated)	Not Supported Now	This class is not supported by the risc compiler in the current release.
sc_sensitive_pos	class (deprecated)	Not Supported Now	This class is not supported by the risc compiler in the current release.
sc_sensitive	class	Not Supported Now	This class is not supported by the risc compiler in the current release.
sc_set_default_time_unit	function (deprecated)	Supported	
sc_set_stop_mode	function	Undecided	Work on this function in the future

Table 6: RISC V0.3.0 Out-of-Order Parallel Simulatable SystemC Subset (continued)

Name	Type	Supported or not	Notes
sc_set_time_resolution	function	Supported	
sc_set_vcd_time_unit	member function (deprecated)	Undecided	Work on this function in the future
sc_signal_in_if	class	Supported	
sc_signal_in_if<bool>	class	Supported	
sc_signal_in_if<sc_logic>	class	Supported	
sc_signal_inout_if	class	Supported	
sc_signal_out_if	type (deprecated)	Supported	
sc_signal_resolved	class	Supported	
sc_signal_rv	class	Supported	
sc_signal_write_if	class	Supported	
sc_signal	class	Supported	sc_signal::trace() is not supported by the out-of-order simulation in the current release.
sc_signal<bool>	class	Supported	sc_signal<bool>::trace() is not supported by the out-of-order simulation in the current release.
sc_signal<sc_logic>	class	Supported	sc_signal<sc_logic>::trace() is not supported by the out-of-order simulation in the current release.
sc_signed_bitref_r	class	Undecided	Work on this class in the future
sc_signed_bitref	class	Undecided	Work on this class in the future
sc_signed_subref_r	class	Undecided	Work on this class in the future
sc_signed_subref	class	Undecided	Work on this class in the future
sc_signed	class	Supported	
sc_simcontext	class (deprecated)	Supported	sc_simcontext::initial_crunch(), cycle() and other functions are partially supported by the out-of-order simulation in the current release.
sc_simulation_time	function (deprecated)	Supported	
sc_spawn_options	class	Supported	
sc_spawn	function	Not Supported Now	sc_spawn() is not supported by the out-of-order simulation in the current release.
sc_start_of_simulation_invoked	function	Undecided	Work on this function in the future
sc_start	function	Supported	
sc_start(double)	function	Not Supported Now	This function is not supported by the out-of-order simulation in the current release.
sc_status	enumeration	Supported	

Table 7: RISC V0.3.0 Out-of-Order Parallel Simulatable SystemC Subset (continued)

Name	Type	Supported or not	Notes
sc_stop_here	function	Undecided	Work on this function in the future
sc_stop	function	Supported	new support as of V0.3.0
sc_string	class (deprecated)	Undecided	Work on this class in the future
sc_subref_r	class template	Undecided	Work on this class template in the future
sc_subref	class	Undecided	Work on this class in the future
sc_switch	enumeration	Supported	
sc_thread_process	class	Supported	
SC_THREAD	macro	Supported	
sc_time	class	Supported	
sc_time_stamp	function	Supported	
sc_time_to_pending_activity	function	Undecided	Work on this function in the future
sc_trace_delta_cycles	function (deprecated)	Undecided	Work on this function in the future
sc_trace_file	class	Undecided	Work on this class in the future
sc_trace	function	Undecided	Work on this function in the future
sc_ufix_fast	class	Undecided	Work on this class in the future
sc_ufix	class	Supported	
sc_ufixed_fast	class template	Undecided	Work on this class template in the future
sc_ufixed	class template	Supported	
sc_uint_base	class	Supported	
sc_uint_bitref_r	class	Undecided	Work on this class in the future
sc_uint_bitref	class	Undecided	Work on this class in the future
sc_uint_subref_r	class	Undecided	Work on this class in the future
sc_uint_subref	class	Undecided	Work on this class in the future
sc_uint	class template	Supported	
sc_unsigned_bitref_r	class	Undecided	Work on this class in the future
sc_unsigned_bitref	class	Undecided	Work on this class in the future
sc_unsigned_subref_r	class	Undecided	Work on this class in the future
sc_unsigned_subref	class	Undecided	Work on this class in the future
sc_unsigned	class	Supported	
sc_unwind_exception	class	Undecided	Work on this class in the future
sc_value_base	class	Undecided	Work on this class in the future
sc_vector_assembly	class	Undecided	Work on this class in the future
sc_vector_base	class	Undecided	Work on this class in the future
sc_vector	class	Undecided	Work on this class in the future
sc_version_major	function	Supported	
sc_version_minor	function	Supported	
sc_version_originator	function	Supported	
sc_version_patch	function	Supported	

Table 8: RISC V0.3.0 Out-of-Order Parallel Simulatable SystemC Subset (continued)

Name	Type	Supported or not	Notes
sc_version_prerelease	function	Supported	
sc_version_release_date	function	Supported	
sc_version_string	function	Supported	
sc_version	function	Supported	
wait	function	Limited Support	wait(void) is not supported
next_trigger	function	Not Supported Now	This function is not supported by the risc compiler in the current release.
halt	function	Not Supported Now	This function is not supported by the risc compiler in the current release.

Connectivity and communication of the instantiated components is supported through ports (`sc_port`, `sc_in`, `sc_inout`, `sc_out`) and interfaces (`sc_interface`).

In contrast to the traditional Accellera library, which only provides a type definition `sc_channel` to `sc_module`, the RISC header files clearly distinguish channels from modules. Here, a separate `sc_channel` class is inherited from `sc_module`, providing the same functionality, but making the two classes explicit.

Most of the SystemC predefined primitive channels[2] (such as `sc_signal` and `sc_fifo`) are supported for OoO PDES, except `sc_fifo::trace` and `sc_fifo::operator=` which are not supported in the current release. For more details, please refer to the Doxygen-generated documentation [29].

4.2 SystemC Threads

The explicit and statically or dynamically [28] analyzable multi-threading of a SystemC design model is naturally supported in RISC OoO PDES. This includes SystemC processes (`SC_HAS_PROCESS`, `sc_process_handle`, `sc_cthread_process`, `sc_method_process`, `sc_thread_process`) and the corresponding threads and methods (`SC_CTHREAD`, `SC_METHOD`, `SC_THREAD`). For basic inter-thread synchronization, SystemC event notifications (`sc_event.notify`) and waiting for events or simulation time advance (`sc_wait`) are supported.

However, dynamic SystemC thread creation and deletion (`sc_spawn`, `SC_FORK`, `SC_JOIN`) is not supported at this time.

While the application programming interface (API) for these constructs remains unmodified from the SystemC user perspective, the RISC SystemC kernel internally supports extra parameters or arguments for these constructs which are utilized after the automatic source code instrumentation by the RISC compiler (see Section 3.3 above). In particular, segment and instance identifiers are supplied with each of these function calls so that the simulator kernel is aware of the exact thread state upon every scheduler entry. This includes in particular the thread creation constructs (`SC_CTHREAD`, `SC_METHOD`, `SC_THREAD`) and wait (`sc_wait`) statements, as well as standard communication interface methods (e.g. `sc_fifo_in_if::read`).

[2] As described in Section 2.4 and Section 3.6, the RISC Compiler and Simulator Beta Release V0.3.0 falls back to synchronous PDES execution when primitive channels with update requests are used in the design model.

4.3 SystemC Transaction Level Modeling (TLM)

While transaction level modeling in general is a natural feature supported by OoO PDES [15], the modeling and implementation choices made by SystemC TLM 2.0 [31] create significant problems for supporting it efficiently in RISC. The root problem here lies in the elimination of explicit channels, which were a key contribution in the early days of research on system-level design [16, 17, 18]. As most researchers agreed, the concept of separation of concerns was of highest importance, and for system-level design in particular, this meant the clear separation of computation (in behaviors or modules) and communication (in channels).

Regrettably, SystemC TLM 2.0 chose to implement communication interfaces directly as sockets in modules [33] and this indifference between channels and modules thus breaks the assumption of communication being safely encapsulated in channels. Without such channels, there is very little opportunity for safe parallel execution.

While a discussion at the SystemC Language Working Group [3] has started [32], at this point, it is unclear how this situation can be worked around or corrected. Thus, SystemC TLM 2.0 can currently not be supported by RISC.

4.4 SystemC Datatypes

A large part of the SystemC language covers special data types designed for bit-accurate hardware modeling, simulation time representation, and other ESL specifics. These SystemC data types include `sc_bigint`, `sc_biguint`, `sc_bit`, `sc_bv`, `sc_fix`, `sc_ufix`, `sc_fixed`, `sc_ufixed`, `sc_int`, `sc_uint`, `sc_logic`, and `sc_lv`.

While all these SystemC data types are available in RISC, only a few of them have been validated and tested for being safe in a truly parallel multi-threading context. At this point, RISC supports `sc_int`, `sc_uint`, `sc_fixed`, and `sc_ufixed` (which are MT-safe). All other data types are so far untested and may or may not be safely used in OoO PDES.

4.5 SystemC Utilities and Other Constructs

As listed in Table 1 through Table 8, there is a plethora of other SystemC APIs available. Some of these are easily supported in RISC (such as `sc_copyright`, `sc_version_major`, `sc_version_minor`, `sc_version_patch`, `sc_version`), others are not supported at this time, such as the SystemC built-in tracing features (`sc_trace`, `sc_trace_file`).

At this point, there is also a large number of special SystemC constructs for which it is unclear whether or not these can be supported in an OoO PDES context with reasonable effort and efficiency. An example of such constructs are those functions which involve or allow to inspect the simulator state at run-time, such as `sc_find_event`, `sc_find_object`, `sc_get_current_process_handle`, `sc_get_status`, `sc_get_time_resolution`, `sc_get_top_level_events`, `sc_get_top_level_objects`, `sc_hierarchical_name_exists`, `sc_is_running`, `sc_is_unwinding`, `sc_simcontext`, and `sc_status`.

On the other hand, access to the current simulated time (`sc_time`, `sc_simulation_time`, `sc_delta_count`), an essential part of every SystemC model evaluation, is fully supported by RISC OoO PDES.

5 Conclusion

While SystemC is the de-facto and official standard language for ESL design, SystemC simulation largely is still performed sequentially following classic DES semantics. Thus, SystemC simulation cannot utilize the parallel

processing capabilities available on today's multi- and many-core host computers.

In this report, we have described the Recoding Infrastructure for SystemC (RISC), an agressive simulation approach beyond traditional parallel DES, where a dedicated SystemC compiler and advanced parallel simulator implement Out-of-Order Parallel Discrete Event Simulation (OoO PDES) for SystemC. This approach promises to exploit parallel computing resources to the maximum extend and thus fastest simulation speed. At the same time, OoO PDES maintains the traditional SystemC modeling semantics.

At this time, this technical report documents the RISC Compiler and Simulator and details the SystemC subset supported by the RISC Beta Release V0.3.0. In contrast to the previous Alpha Release V0.2.1, the RISC Compiler and Simulator Beta Release V0.3.0 is more robust and easier to install, features new support for dynamic conflict analysis (see Section 3.2.2), safely supports primitive channels with update methods, offers new support of library functions by use of `#pragma` annotations (see Section 3.4), and provides new support for the Intel compiler and special target processors in the backend (see Section 3.5).

As we move on in the project, we will update this report and in particular the supported subset tables accordingly.

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A Appendix

A.1 Manual Page of the RISC Compiler and Simulator

NAME

risc – Recoding Infrastructure for SystemC (RISC) Compiler and Simulator

SYNOPSIS

risc [*options*] *design* [*options*]

DESCRIPTION

risc is a dedicated compiler for the SystemC language. The purpose of **risc** is to parse, analyze, instrument, and compile a SystemC source program into an executable program for out-of-order parallel simulation. **risc** is a frontend source-to-source compiler for SystemC built on top of the ROSE compiler infrastructure with GNU or Intel C++ as backend target compiler. As such, **risc** relies on and supports also most of the ROSE and GNU compiler options.

Using the command syntax shown in the synopsis above, the specified *design* is compiled. By default, **risc** reads the SystemC source file, performs preprocessing and builds an internal representation (abstract syntax tree) and a Segment Graph (SG) of the model. Next, segment conflict analysis is performed and the design model is instrumented for Out-of-Order Parallel Discrete Event Simulation (OoO PDES). Finally, instrumented C++ code is generated, compiled, and linked into an executable file that can be run for fast parallel simulation.

On successful completion, the exit value 0 is returned. In case of errors during processing, an error code with a brief diagnostic message is written to the standard error stream and the compilation is aborted with an exit value greater than zero.

For preprocessing and C++ compilation into an executable file, **risc** relies on the availability of an external C++ compiler which is used automatically in the background. By default, the GNU C++ compiler **g++** is used. Alternatively (see options *–risc:icpc* and *–risc:mic* below), the Intel C++ compiler **icpc** may be used to generate an executable optimized for Intel processors with SIMD capabilities or the Intel Many-Integrated-Core (MIC) architecture.

ARGUMENTS

design specifies the file name of the input SystemC design model; by default, the base name of *design* is used as base name for the intermediate and output files;

OPTIONS

–h | *—-help* print the **risc** compiler version and a brief usage information message to standard output and quit;

–v | *—-verbose* increment the verbosity level so that all tasks performed are logged to standard error (default: be silent); at level 1, high-level messages about the tasks performed are displayed; at level 2, additional details such as input and output file names are listed; at level 3, very detailed information about each executed task is printed;

–vv increment the verbosity level by two counts (same as *–v –v*);

–vvv increment the verbosity level by three counts (same as *–v –v –v*);

–w | —-warnings increment the warning level so that compiler warning messages are enabled (default: warnings are disabled); four levels are supported ranging from only important warnings (level 1) to pedantic warnings (level 4); for most cases, warning level 2 is recommended (*–w –w*);

–ww increment the warning level by two counts (same as *–w –w*);

–www increment the warning level by three counts (same as *–w –w –w*);

–g add a symbol table suitable for debugging (e.g. using **gdb**) to the generated object files and simulation executable (default: no debugging symbols);

–O | –O level optimize the generated simulation executable for higher execution speed and/or less memory usage (default: no optimization);

–Idir add the specified *dir* to the include path (extend the list of directories to be searched for including source files); include directories are searched in the order of their specification; the standard include path ($SYSTEMC_LW_HOME/include or $SYSTEMC_OOP_HOME/include) is automatically appended to this list; by default, only the standard include directories are searched;

–Ldir add the specified *dir* to the library path (extend the list of directories to be searched for linker libraries); the library path is searched in the specified order; the standard library path ($SYSTEMC_OOP_HOME/lib) is automatically appended to this list; by default, only the standard library path is searched;

–llib add the specified *lib* to the list of libraries for the linker so that the executable is linked against *lib;* libraries are linked in the specified order; the standard libraries (i.e. -lsystemc) are automatically appended to this list; by default, only standard libraries are used;

–c perform only the preprocessing, analysis, instrumentation, and compilation tasks; skip the final linking stage so that only an object file is created (default: perform all tasks including linking);

–o output file specify the name of the final output file explicitly (default: a.out);

–risc:dump output the computed segment graph (SG) and conflict tables as HTML files (default: no HTML files are generated); these files may be viewed by a user in a browser in order to inspect the out-of-order execution conditions of the model and improve it accordingly;

–risc:icpc use the Intel C++ compiler **icpc** in the backend for generating the executable (default: GNU C++ compiler **g++**);

–risc:mic use the Intel C++ compiler **icpc** with option *–mic* in the backend for cross-compiling an executable for the Intel Many Integrated Core (MIC) architecture (default: generate an executable for the same processor the compiler is running on);

–*risc:elab filename* import the elaboration result produced by the RISC elaborator from file *filename* and use it for more precise segment conflict analysis (default: pure static analysis);

–<*rose:option*> pass this option through to the underlying ROSE compiler (default: none);

–<*GNU option*> pass this option through to the underlying GNU compiler (default: none);

ENVIRONMENT

RISC is used at compile-time to determine the installation directory of the RISC compiler and simulator where the RISC system components are located (default: none);

SYSTEMC_LW_HOME is used at compile-time to find the RISC light-weight SystemC header files which are expected in directory $SYSTEMC_LW_HOME/include (default: none);

SYSTEMC_OOP_HOME is used at compile-time to find the RISC out-of-order SystemC header files which are expected in directory $SYSTEMC_OOP_HOME/include, and the RISC out-of-order SystemC library which is expected in directory $SYSTEMC_OOP_HOME/lib (default: none);

SYSC_PAR_SIM_CPUS is used by the RISC simulator at run-time to set the maximum number of concurrent threads allowed in the RISC out-of-order SystemC simulation (default: 64);

SYSC_SYNC_PAR_SIM is used by the RISC simulator at run-time to force the RISC out-of-order SystemC simulation to fall back to synchronous (in-order) PDES execution; note that this mode is also automatically selected when SystemC primitive channels are used with update requests (default: false);

VERSION

The RISC compiler and simulator is beta release version 0.3.0.

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BUGS, LIMITATIONS

Possibly many, since this is a beta release of a proof-of-concept prototype implementation.

A.2 Manual Page of the RISC Elaborator

NAME

elab – Recoding Infrastructure for SystemC (RISC) Dynamic Elaborator

SYNOPSIS

elab *design* [*options*]

DESCRIPTION

elab is a special compiler for the SystemC language. The purpose of **elab** is to parse, analyze, instrument, and compile a SystemC source program into an executable program for dynamic elaboration. **elab** is a frontend source-to-source compiler for SystemC built on top of the ROSE compiler infrastructure with GNU or Intel C++ as backend target compiler. As such, **elab** relies on and supports also most of the ROSE and GNU compiler options.

Using the command syntax shown in the synopsis above, the specified *design* is compiled. By default, **elab** reads the SystemC source file, performs preprocessing and builds an internal representation (abstract syntax tree) of the SystemC structural hierarchy. **elab** then instruments the design model so that its execution stops after the end of the elaboration phase (no actual simulation will take place); the dynamically built hierarchy and instance connectivity data is then dumped into a file *design.elab* which can be passed to the RISC compiler **risc** for more precise conflict analysis.

On successful completion, the exit value 0 is returned. In case of errors during processing, an error code with a brief diagnostic message is written to the standard error stream and the compilation is aborted with an exit value greater than zero.

For preprocessing and C++ compilation into an executable file, **elab** relies on the availability of an external C++ compiler which is used automatically in the background. By default, the GNU C++ compiler **g++** is used.

ARGUMENTS

design specifies the file name of the input SystemC design model; by default, the base name of *design* is used as base name for the intermediate and output files;

OPTIONS

–h | —-help print the **elab** elaborator version and a brief usage information message to standard output and quit;

–v | —-verbose increment the verbosity level so that all tasks performed are logged to standard error (default: be silent); at level 1, high-level messages about the tasks performed are displayed; at level 2, additional details such as input and output file names are listed; at level 3, very detailed information about each executed task is printed;

–vv increment the verbosity level by two counts (same as *–v –v*);

–vvv increment the verbosity level by three counts (same as *–v –v –v*);

–w | —-warnings increment the warning level so that compiler warning messages are enabled (default: warnings are disabled); four levels are supported ranging from only important warnings (level 1) to pedantic warnings (level 4); for most cases, warning level 2 is recommended (*–w –w*);

–ww increment the warning level by two counts (same as *–w –w*);

–www increment the warning level by three counts (same as *–w –w –w*);

–g add a symbol table suitable for debugging (e.g. using **gdb**) to the generated object files and simulation executable (default: no debugging symbols);

–O | –O level optimize the generated simulation executable for higher execution speed and/or less memory usage (default: no optimization);

–Idir add the specified *dir* to the include path (extend the list of directories to be searched for including source files); include directories are searched in the order of their specification; the standard include path ($SYSTEMC_LW_HOME/include or $SYSTEMC_OOP_HOME/include) is automatically appended to this list; by default, only the standard include directories are searched;

–Ldir add the specified *dir* to the library path (extend the list of directories to be searched for linker libraries); the library path is searched in the specified order; the standard library path ($SYSTEMC_OOP_HOME/lib) is automatically appended to this list; by default, only the standard library path is searched;

–llib add the specified *lib* to the list of libraries for the linker so that the executable is linked against *lib;* libraries are linked in the specified order; the standard libraries (i.e. -lsystemc) are automatically appended to this list; by default, only standard libraries are used;

–c perform only the preprocessing, analysis, instrumentation, and compilation tasks; skip the final linking stage so that only an object file is created (default: perform all tasks including linking);

–o output file specify the name of the final output file explicitly (default: a.out);

–elab:o specify the name of the elaboration result file with instance connectivity data explicitly (default: *design.elab*); this file will be produced when the executable generated by **elab** is run (after its elaboration phase);

–<rose:option> pass this option through to the underlying ROSE compiler (default: none);

–<GNU option> pass this option through to the underlying GNU compiler (default: none);

ENVIRONMENT

RISC is used at compile-time to determine the installation directory of the RISC compiler and simulator where the RISC system components are located (default: none);

SYSTEMC_LW_HOME is used at compile-time to find the RISC light-weight SystemC header files which are expected in directory $SYSTEMC_LW_HOME/include (default: none);

SYSTEMC_OOP_HOME is used at compile-time to find the RISC out-of-order SystemC header files which are expected in directory $SYSTEMC_OOP_HOME/include, and the RISC out-of-order SystemC library which is expected in directory $SYSTEMC_OOP_HOME/lib (default: none);

VERSION

The RISC elaborator is beta release version 0.3.0.

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